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Modeling and Simulation of a Fiber Distributed Data Interface Local Area Network (FDDI LAN) Using OPNET® for Interfacing Through the Common Data Link (CDL)

by

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of the requirements for the degree of

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The Optimized Network Engineering Tool (OPNET®) is a commercially available communications network simulation package. This thesis involves the modification of OPNET®'s Fiber Distributed Data Interface Local Area Network (FDDI LAN) model in order to enhance its usefulness as an aid in the development of recommendations for the characteristics and metrics to be eventually included in the Defense Service Project Office's (DSPO) Common Data Link (CDL) project. This work includes a step-by-step guide for FDDI simulation in OPNET®, and a discussion of the changes made to the original model to enhance its performance and data display characteristics. Simple tests are provided to verify the completed model's performance and usefulness as a working tool for further development.

TABLE OF CONTENTS

I. INTRODUCTION

A. PROBLEM STATEMENT

The simulation model described in this thesis was developed in support of the Defense Support Project Office's (DSPO) Common Data Link (CDL) project. The Common Data Link is a full duplex, jam resistant, point-to-point microwave communication system for use in imaging and signals intelligence collection systems (DSPO, 1993, p. 1). Essentially, CDL is to provide a protocol for communication between two or more Fiber Distributed Data Interface Local Area Networks (FDDI LAN). These include an airborne LAN providing sensor information with high data transfer rates, and a ground based LAN providing command and control information.

This work is concerned primarily with the modification and testing of a commercially available communications network simulation program, MIL 3, Inc.'s Optimized Network Engineering Tool (OPNET®). This thesis represents the first portion of three relatively independent research tasks being performed as MS theses to provide evaluations of several Network Interfaces (NI) to the CDL and a multilink point-to-point protocol, in support of the CDL project.

B. SCOPE

The scope of this thesis includes the following:

- Introduce the CDL concept as the context in which the FDDI simulation model is to be modified and tested.
- Provide a tutorial style introduction to the OPNET® FDDI model, designed to expand upon the tutorial provided by the manufacturer. This is directed to those who will conduct further studies in the CDL project, and also to students whose class laboratory work will include simulations in OPNET®.
- Discuss in detail the modifications made to the given model. Provide analysis of the model's actual simulation performance as a validation of the model's usefulness to the CDL work at NPS through comparisons against trials published in the research literature using other simulation tools.

C. BENEFITS

The primary contribution of this thesis is the development of a functioning simulation model that will support the features typically required in a CDL deployment scenario. Typical data communication requirements include the following:

- a wide range of data rates,
- a wide range of error rates and types of error correction required,
- real-time requirements such as user-specified delivery delays and its variation (jitter),

- connection requirements (whether connection-oriented or connectionless, multicasting, broadcasting, etc.),
- retransmission requirements,
- coupling and synchronization with other data sources, and
- adjustable prioritization relative to other sources.

The second benefit is to document in detail the MIL 3, Inc.'s FDDI LAN simulation model in its operation and in its modification. The third benefit lies in the use of the developed model as an instructional tool for classroom laboratory exercises supporting the study of FDDI LAN operation.

D. ORGANIZATION

This thesis is organized as follows. Chapter II provides a tutorial on the use of the FDDI LAN model provided with OPNET®. Chapter III addresses the details of the modifications made to the given model to simulate multicasting and priority-based traffic. Where applicable, clarifications regarding the OPNET® manuals are highlighted. Chapter IV presents the results of simulation tests intended to verify the validity of the modified model. The thesis ends with conclusions and recommendations for future work in Chapter V.

II. MODELING AN FDDI LAN IN OPNET®

A. OVERVIEW

This chapter is intended to provide a tutorial on the use of OPNET® to model an FDDI LAN by providing a brief set of steps to build and execute a simulation. The current version as of this writing is Release 2.4.A, dated 02/27/93, which is the third revision. Release 2.4.A, Errata 1, dated 08/01/93 is a manual update. Some prerequisite knowledge is required of the user, including "C" programming language syntax, ability to use a UNIX workstation, and an understanding of the FDDI protocol. MIL 3, Inc. provides thorough documentation in the form of an eleven volume set of manuals, the first of which is Vol. 1.0, entitled: *Tutorial Manual*. It includes a general introduction to OPNET®, a trouble-shooting guide, and five chapters presenting different communications network models. While none of these discusses FDDI in particular, all are designed to familiarize the novice user with the mechanics of the user interface, and should be studied prior to working with OPNET®. Volumes 4.0 and 4.1, the *Tool Operations Manual*, describes the editors of the user interface, and should likewise be studied. The chapter entitled "FDDI" in Vol. 8.1.0, *Example Models Manual, Protocol Models*, discusses the FDDI simulation in detail, and provides the essential information to build, develop and execute a simulation. Most of the information presented here is available in the manuals, but a number of idiosyncrasies exist which are not readily documented. These required trial and error experimentation to discover, and in many cases required explanation

from MIL 3, Inc.'s excellent technical support organization. The new user is advised to heed every sentence regarding mechanical details; much of the advice given is hard-earned.

This thesis will not present an explanation of the FDDI protocol in detail, except as necessary to emphasize or clarify the operation of the model. Many discussions exist in the research literature and textbooks, for example, Stallings, (1991, 1993). Those interested in the physical characteristics of optical fiber systems are referred to Powers, (1993). A useful introduction to modeling FDDI in OPNET® is *Modeling and Simulation of a Fiber Distributed Data Interface Local Area Network*, a Naval Postgraduate School MSEE thesis by Aldo Schenone, which summarizes OPNET®'s FDDI model, includes a detailed description of the FDDI protocol, presents the results of several simulations, and ends with a challenge to other researchers to further develop the model. (Schenone, 1993)

The remainder of this chapter will briefly introduce the structure of OPNET®'s FDDI LAN model and its components, introduce some preliminary modifications, then lead the reader through a simple simulation.

B. PRELIMINARY THEORY

1. Theoretical vs OPNET® Model of FDDI

The setting of parameters in OPNET® simulations is based on the following equation and discussion, which is found in most literature treating FDDI LANs, including

Powers (1993, p. 340), Stallings (1993, p. 225), Tari, et al (1988, p.55), and Jain, (1991, p.20), to name a few:

$$D\_Max + F\_Max + TokenTime + \Sigma SA_i \leq TTRT \quad (1).$$

where:

SA_i = synchronous allocation for station i,

D_Max = propagation time for one complete circuit of the ring,

F_Max = time required to transmit a maximum-length frame (4500 octets), and

$TokenTime$ = time required to transmit a token.

In an actual LAN, a station management protocol handles the assignments of SA_i, which may be changed in real time. OPNET® simulations represent steady-state performance, and do not contemplate changing network conditions.

All stations negotiate a common value of TTRT. Also, these timers and variables are maintained at each station:

- Token Rotation Timer (TRT)
- Token Holding Timer (THT)
- Late counter (LC)

Each station is initialized to the same TRT, which is set to TTRT. Note that LC, TRT and THT are not global; each station maintains its own copies, which will differ from those of other stations. If a given station receives the token before its TRT has expired, then that TRT is reset to TTRT. On the other hand, should the token arrive after the expiration

of TRT, then its lateness is recorded by setting LC to "1" (at that station). Two consecutive late tokens will increment LC to "2", in which case the token is considered to be lost, and a re-initialization process will commence. OPNET® has no provision for re-initialization. On the other hand, as a computerized simulation, it never permits LC=2 to occur.

When the token arrives early, (before TRT expires), THT is set to the current value of TRT. The transmission rules are as follows:

1. A station may transmit synchronous traffic for a time SA_i, as specified for that station.
2. THT is enabled after synchronous traffic is sent, or if there was no synchronous traffic to send. The station may transmit asynchronous traffic while THT > 0.

In the TX_DATA state of the MAC process model, THT is incremented from zero to THT. This is an important point in regard to the prioritization scheme by which a value T_Pri[i] is assigned to each priority setting, and the eligibility of a given priority to transmit depends on T_Pri[i] in comparison to THT That is, for THT decrementing, priority i traffic may be transmitted as long as T_Pri[i] is less than THT. This implies lower T_Pri[i] is assigned to higher priority. In OPNET®, THT increments up from zero. Priority i traffic may be transmitted as long as T_Pri[i] is greater than THT, implying that higher T_Pri[i] is assigned to higher priority. This subtle point is important to know when setting values to T_Pri[i] in the INIT state of the MAC process model.

There is an important distinction in the timing of transmission eligibility for synchronous and asynchronous traffic:

- The time spent sending synchronous traffic may not exceed SA_i for station *i*. That is, the protocol will not allow a synchronous packet transmission to commence if it can not be completed without exceeding SA_i. OPNET® supports this criterion.
- Asynchronous transmissions may commence as long as THT has not expired. Any packet transmission in progress when THT expires is allowed to complete, but no more will commence.

The protocol allows the actual token rotation time to have a maximum value of (2)TTRT, with an average value of TTRT over time.

2. FDDI LAN Equation Parameters

Each of the terms in Equation 1 given above is addressed in this section, with reference to its representation in OPNET®'s Environment file attributes.

a. "F_Max"

The time required to transmit a maximum length packet is based on the assumption that any station is capable of transmitting at the rate of 100Mbps. Since the maximum packet length is 36,000 bits (4500 octets or bytes), the algebra yields 0.360 ms for F_Max. Powers agrees (1993, p. 340), but Tari et. al. use 0.361 ms (1988, p. 55). In OPNET®, the `fddi_mac` process model defines the transmission rate as 100 Mbps, in the Header Block. F_Max is not directly assigned as an attribute, but simply exists as a physical characteristic which must be considered in determining TTRT and SA_i assignments.

b. "D_Max"

The Maximum Ring Latency is the time required for a frame to travel around the ring. The maximum value is often assumed in textbook discussions, but it should be calculated for individual cases. The total delay may be defined as follows:

```
D_Max = (total fiber length x delay rate)+(number of
     stations x station latency)
```

Powers uses 1.73 ms, Tari et.al. uses 1.62 ms, and Dykeman and Bux use 1.62 ms. D_Max includes the time required for a frame (which is basically a number of light flashes) to travel the length of the fiber on the ring, plus the time required to cross each station interface. 5.085×10^{-06} sec./km. is the value used in the literature for the delay rate of a signal in optical fiber. The reciprocal results in 1.9665e+08 m/s, which agrees with the generally accepted value of ⅔ c for the speed of light in glass. OPNET® and Dykeman & Bux use a station delay of 60.0e-08 sec. Powers assumes 1 μs as a representative value. Ultimately the value is a physical characteristic that could be measured on a real device, and may be declared in a computer simulation. The value 1.617 ms derives from using the maximum possible dimensions: 500 dual attachment stations or 1000 single attachment stations on a 200 km ring yield the following:

$$(1000 \bullet 60.0 \times 10^{-08}) + (200 \bullet 5.085 \times 10^{-06}) = 1.617 \text{ ms.}$$

The environment file attribute prop_delay represents the actual time the packet is on the fiber between two stations, rather than the delay rate, and therefore defines the size of the ring. OPNET® has no safety feature to prevent the user from entering

attributes that would define a ring larger than 200 km. Note that the stations are assumed to be equally spaced. The user should realize that the value of 5.085e-06 given with the original example environment file implies a one kilometer length of fiber, rather than a delay rate. Powers (1993, p. 328) notes that the early proposal for FDDI limits fiber length between stations to 2 kilometers. In OPNET®, the attributes `prop_delay` and `station_latency` are used in the "C" code to postpone the reception of a packet until sufficient time has passed to allow for physical delays.

c. "SA_i"

Synchronous allotment, or synchronous bandwidth, is the time a station is granted to transmit synchronous traffic, regardless of the lateness of the token. It is a form of prioritization, providing a means by which certain types of traffic are not delayed. For example, voice traffic would be synchronous. Textbook discussions represent SA_i in units of time for each station.

Asynchronous traffic is transmitted whenever there is THT remaining after the transmission of all synchronous traffic. It is the responsibility of the Station Management Protocol (in OPNET®, the user) to ensure that the synchronous bandwidth is sufficient to handle the synchronous offered load. One nuance involves the inviolate nature of SA_i for each station. A given station's synchronous offered load may amount to relatively little in terms of bits per second, while the packet size is assigned a value too large to be transmitted in the time SA_i. In this case, synchronous traffic would never be transmitted, and outbound packets would simply accumulate in the buffers of the MAC. The Environment file attribute

"sync bandwidth" corresponds to SA_i, but is expressed as a unitless fraction of TTRT, rather than as a time.

d. "TTRT"

Equation 1 suggests that physical requirements of the fiber and the stations are used to determine a workable TTRT value. The FDDI specification allows a range of settings from 4 ms to 165 ms (Powers, 1993, p. 339). Powers also notes that synchronous voice transmission requires $\sum SA_i = 10$ ms. In OPNET®, the Environment file attribute T_Req corresponds to TTRT.

C. MODEL STRUCTURE

OPNET®'s FDDI LAN model structure is hierarchical. The LAN is a ring made of stations and the connections between them. Figure 1 shows a 50-station FDDI LAN as shown in the user interface window. Figure 2 is a ten station ring provided for greater clarity of detail. The stations are modeled as connected nodes, each of which is in turn defined by a process model. The processes are represented by state transition diagrams, which are the ultimate source of the "C" language code that describes the model's behavior. Figure 3 illustrates the FDDI station model. Figures 4-6 are the process models for the source, sink, and MAC processes, respectively. These correspond one-to-one to the nodes "llc_sink," "llc_sink," and "mac" shown in Figure 3. The packets of information that travel between stations on the ring, and between nodes within the station, are also modeled and may be

Figure 1. 50-Station FDDI LAN, "fddi_net_50"

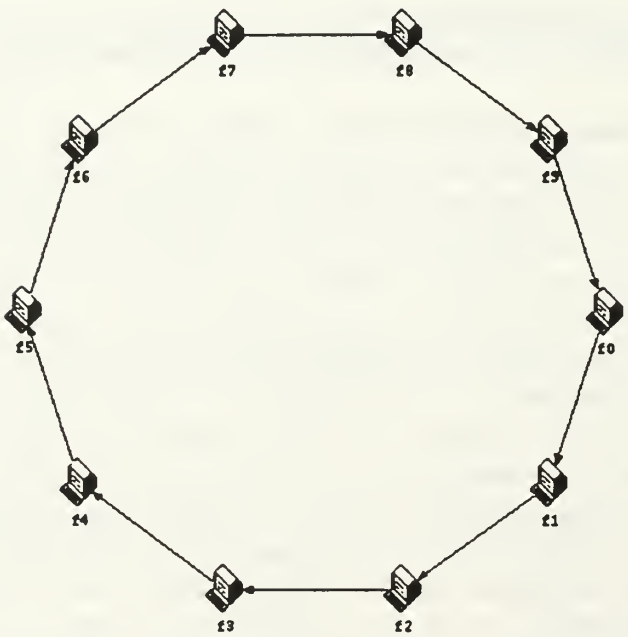

Figure 2. Ten-Station FDDI LAN, "fddi_net_10"

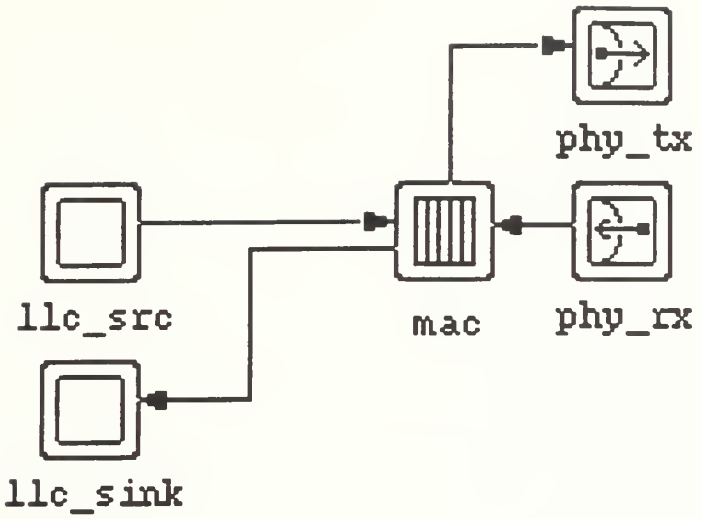

Figure 3. FDDI Station Model, "fddi_station"

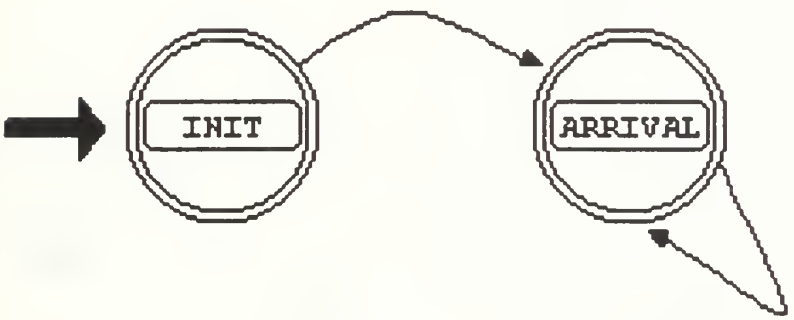

Figure 4. Source Process Model, "fddi_gen"

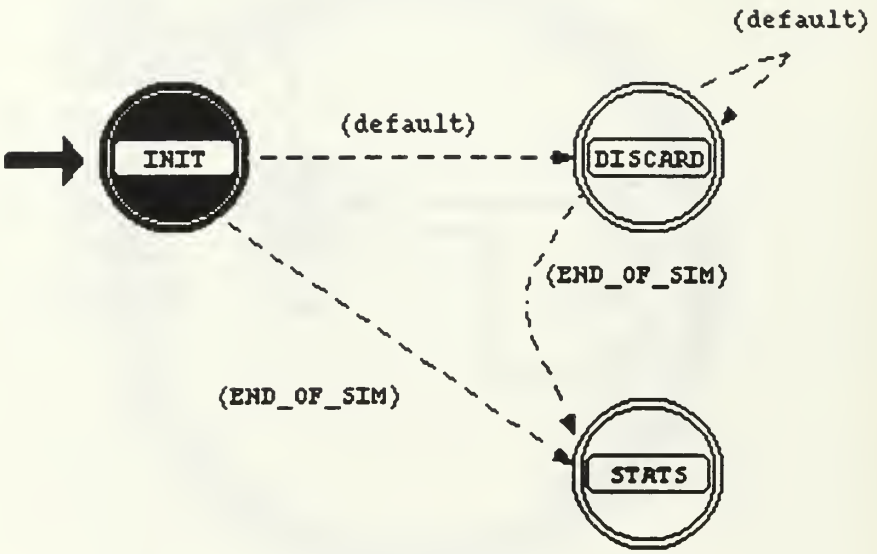

Figure 5. Sink Process Model, "fddi_sink"

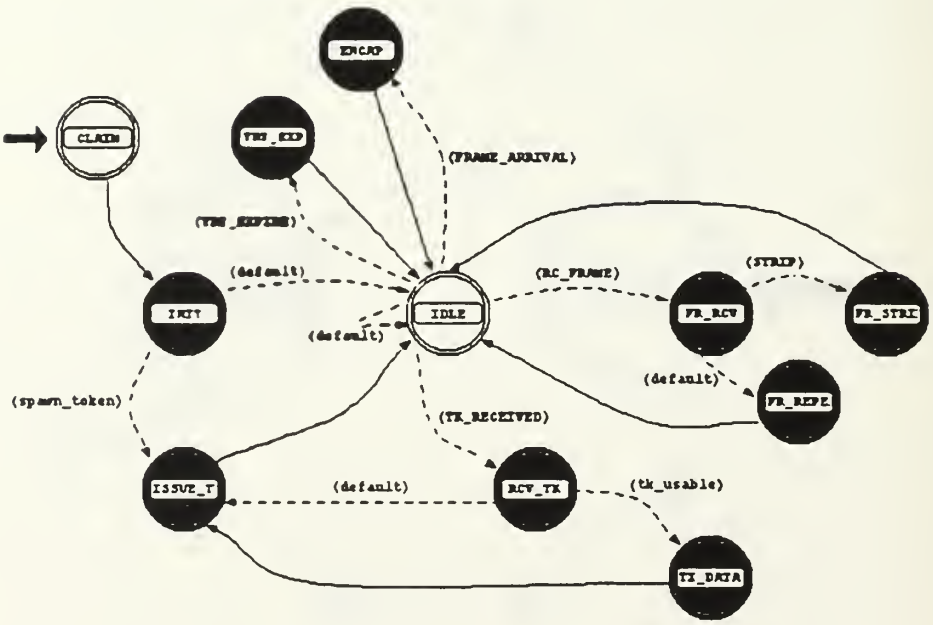

Figure 6. MAC Model, "fddi_mac"

modified. Model parameters may be entered by several methods, with the Environment file being by far the most convenient.

1. FDDI LAN

Figure 7 shows a 32-station FDDI ring in the user interface window as it would appear on a computer screen. This image is displays the Network Editor, whose icon appears toward the upper left corner of the figure. (A note on the mechanics of activating the various editors: as indicated in the tutorial manual, the center mouse button activates the desired editor. If instead either the left or right button is pressed, then the opposite button must also be pressed to "cancel" the first; only then will the center button work as expected). To the right of the ring are on-screen menus of attributes for one station (actually, three menus are shown to display simultaneously all the fields). This menu is invoked by placing the cursor over the desired station, then pressing the right mouse button. The FDDI protocol supports up to 500 dual-attachment stations on a ring, and OPNET® permits from two to 500 stations in a ring.

Actual generation of the ring is best done outside OPNET®, through a UNIX command window set to the "`~\op_models\fddi`" directory path. The command "`fddi_build.em.x <number_of_stations>`" will automatically generate an FDDI LAN with the number of stations specified. The user should verify that this function is present in the desired subdirectory. This operation is described in manual Vol. 8.1.0, "FDDI," and refers the user to Vol. 6.0, *External Interfaces Manual*, which has a more complete description of the ring building protocol. Were OPNET® active during this ring-building process, then the "Rehash" icon toward the lower right of the user interface window must be activated to update the

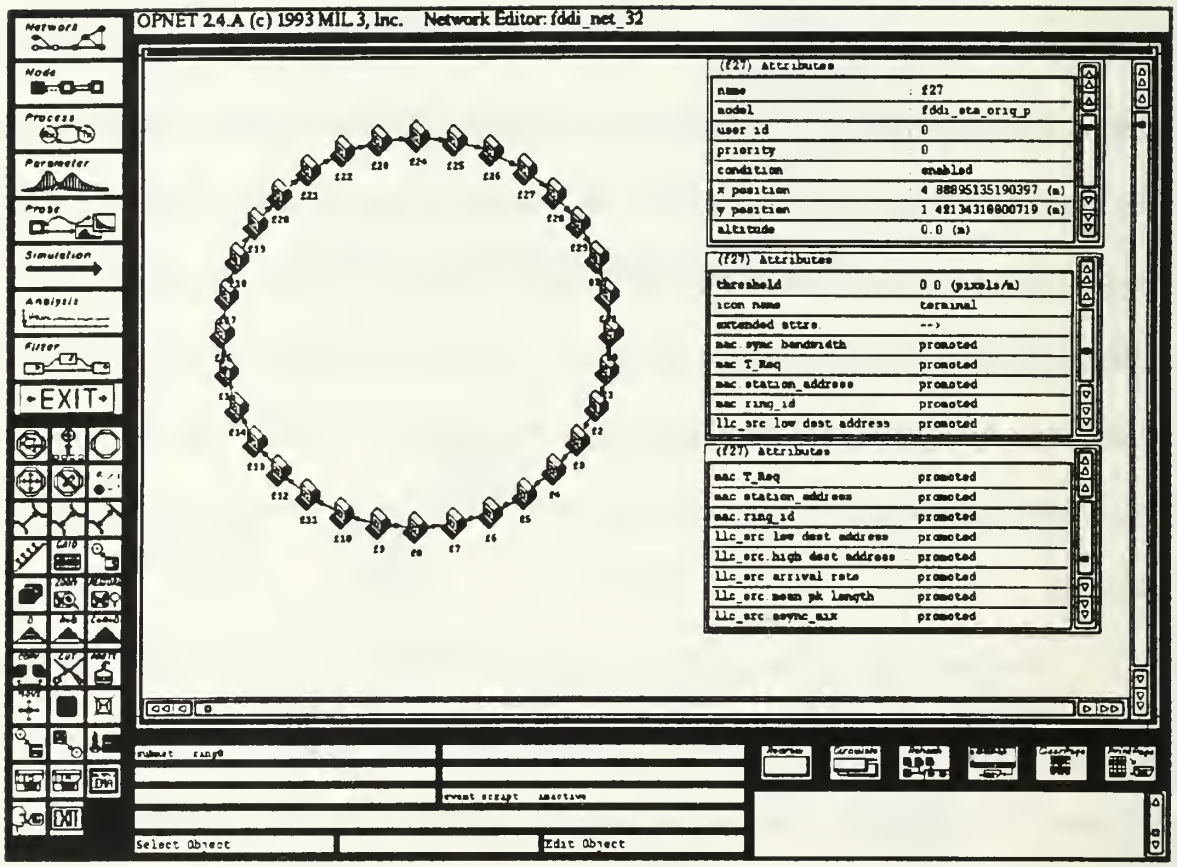

Figure 7. Network and Attributes Menus in User Interface Window

program's access to models in the subdirectory. In general, the "Rehash" command should be used frequently, particularly when new files are generated through simulation runs or through model editing.

The LAN as shown is not actually a true ring architecture, as no dedicated physical layer object exists in OPNET® for modeling ring architectures. The model is in fact a circle of point-to-point links; the ring is an abstraction whose characteristics and behavior are represented in the "C" programs that comprise the process models. (OPNET®, Vol. 8.1.0, "FDDI," p.23)

2. FDDI Station

Figure 8 illustrates the FDDI station in the user interface window, summoned and printed from the Node Editor. Also shown are the menus listing the attributes associated with each part of the station model. Message traffic in the form of packets is generated at the source, llc_src, at a rate specified by the user. The source model does not function as a true Logical Link Control (LLC) beyond correctly interfacing with the Medium Access Control (MAC) model. (OPNET®, Vol. 1.0, "FDDI," p. 21) The MAC entity is represented by mac in the model, and is responsible for encapsulating packets generated by the source, holding these packets until they can be transmitted, receiving packets from other stations, destroying packets as needed, and maintaining the locally held Token Holding Timer (THT) and Token Rotation Timer (TRT). Packets are counted and statistics gathered at llc_sink. These three nodes are modeled in detail by respective process models, which may be assigned with the attributes menus shown in Figure 8. The field "process model" may be changed

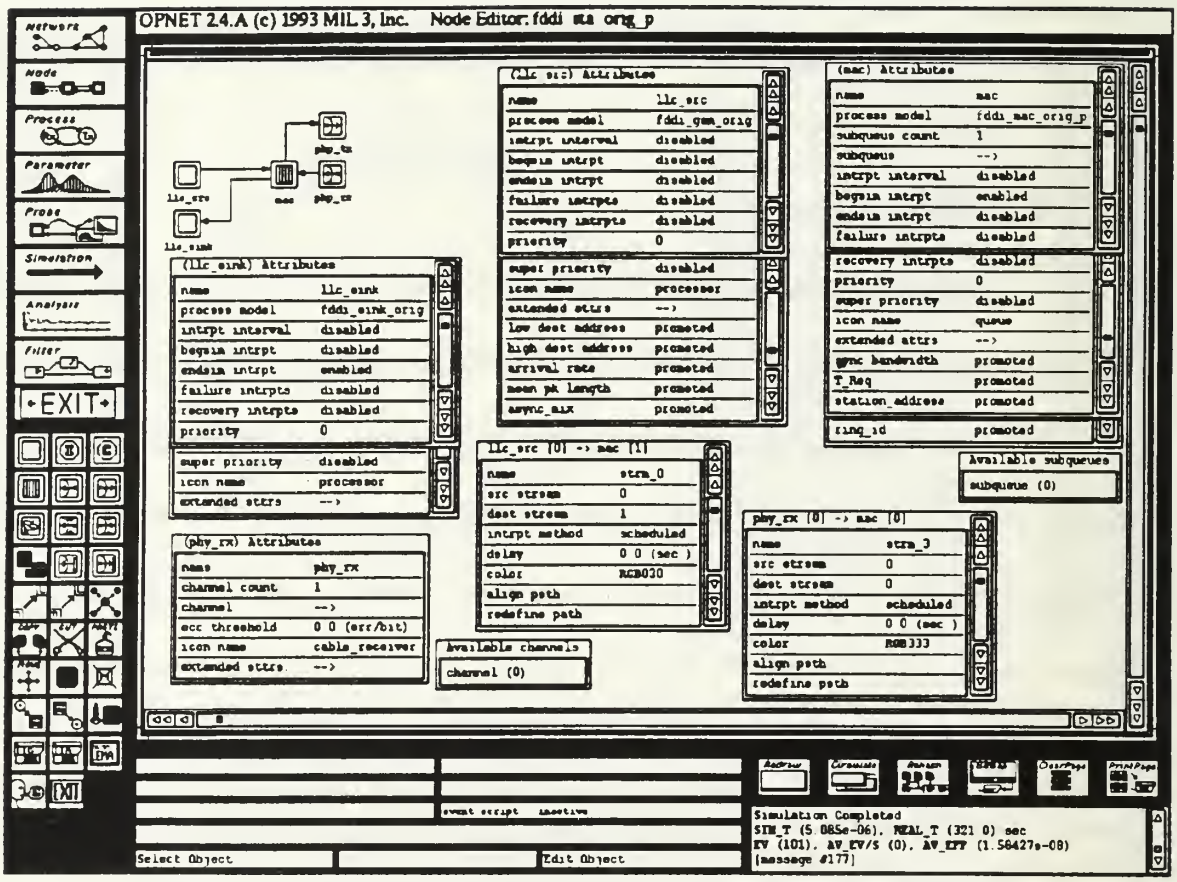

Figure 8. FDDI Station Model with On-Screen Attributes Menus

by cursor action, with selection made from the resulting submenu presenting a list of available process models. The nodes "phy_rx" and "phy_tx" represent the receiver and transmitter interfaces to the ring, and are not further defined by process models.

The user may modify the station model within the Node Editor by setting the attributes fields as desired, then saving the model by activating the "write node model" icon toward the lower left corner of the user interface window. By then exiting the Node Editor, entering the Network Editor, and calling the desired network model (e.g., "fddi_net_32"), each station acquires the new setting when the network model is archived and bound ("A+B" icon). The same modification may be effected from within the Network Editor by calling the attributes menus for each station and setting them individually. This method would be preferred only if the user desires to set differing attributes in various stations. Note that setting the attribute fields is not the same as modifying the process model itself, which is accomplished with changes to the "C" programming code accessed through the Process Editor.

3. Processes

Process models are specified by State Transition Diagrams (STD) representing the actions of the nodes within the station model. Figure 9 illustrates the process model fddi_mac as it appears in the Process Editor in the user interface window. Figure 10 shows the on-screen menu that appears when the cursor is placed over one of the states (ENCAP in this case), and the right mouse button pressed. Invoking the "enter execs" attribute calls the text editor shown in Figure 11. Here the user may inspect the programming code behind

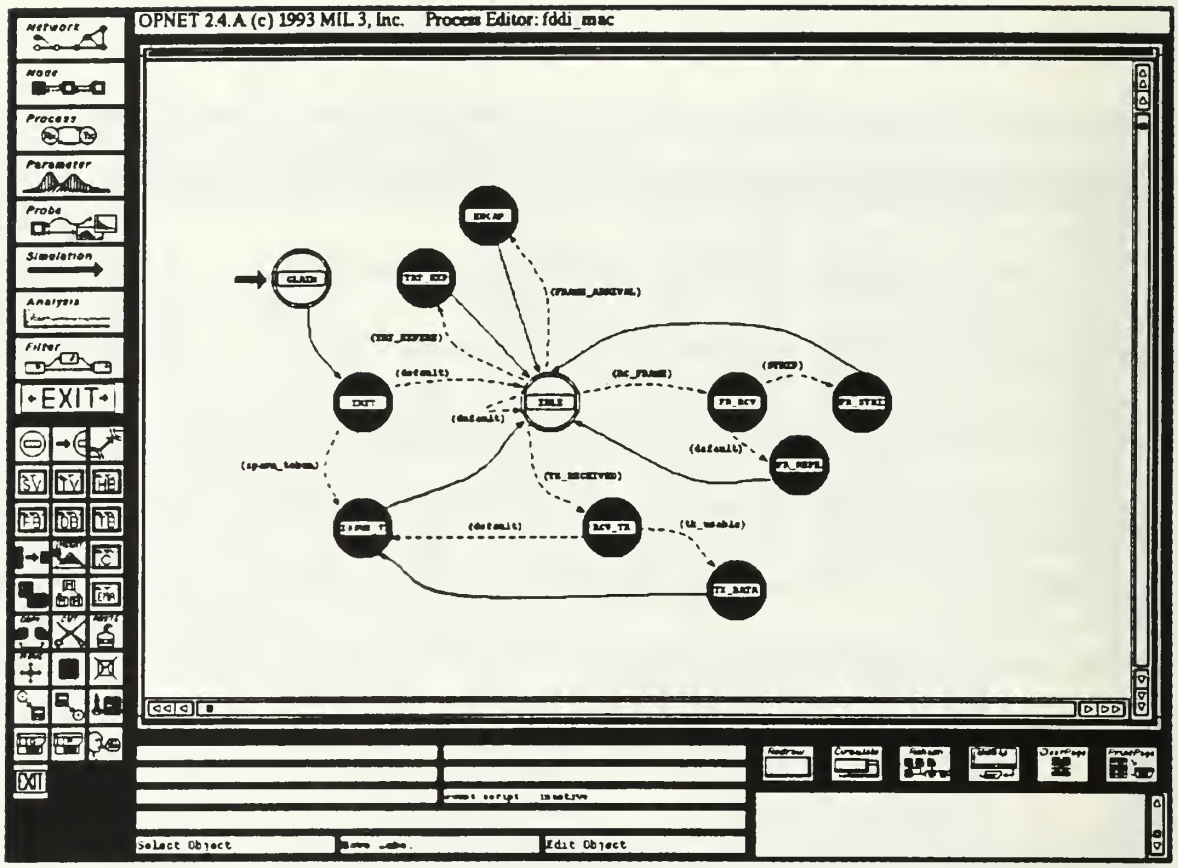

Figure 9. Process Model "fiddi_mac" in User Interface Window

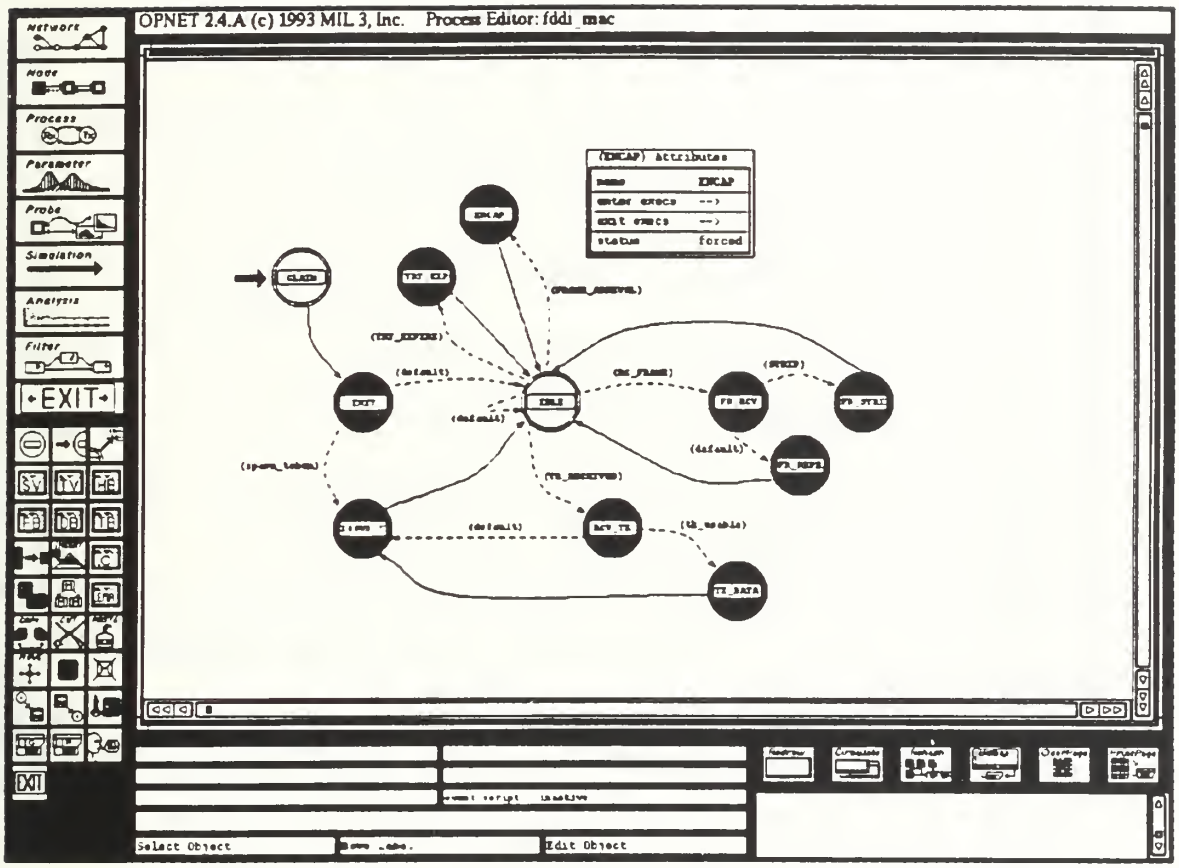

Figure 10. Process Model "fiddi_mac" with On-Screen Attributes Menu

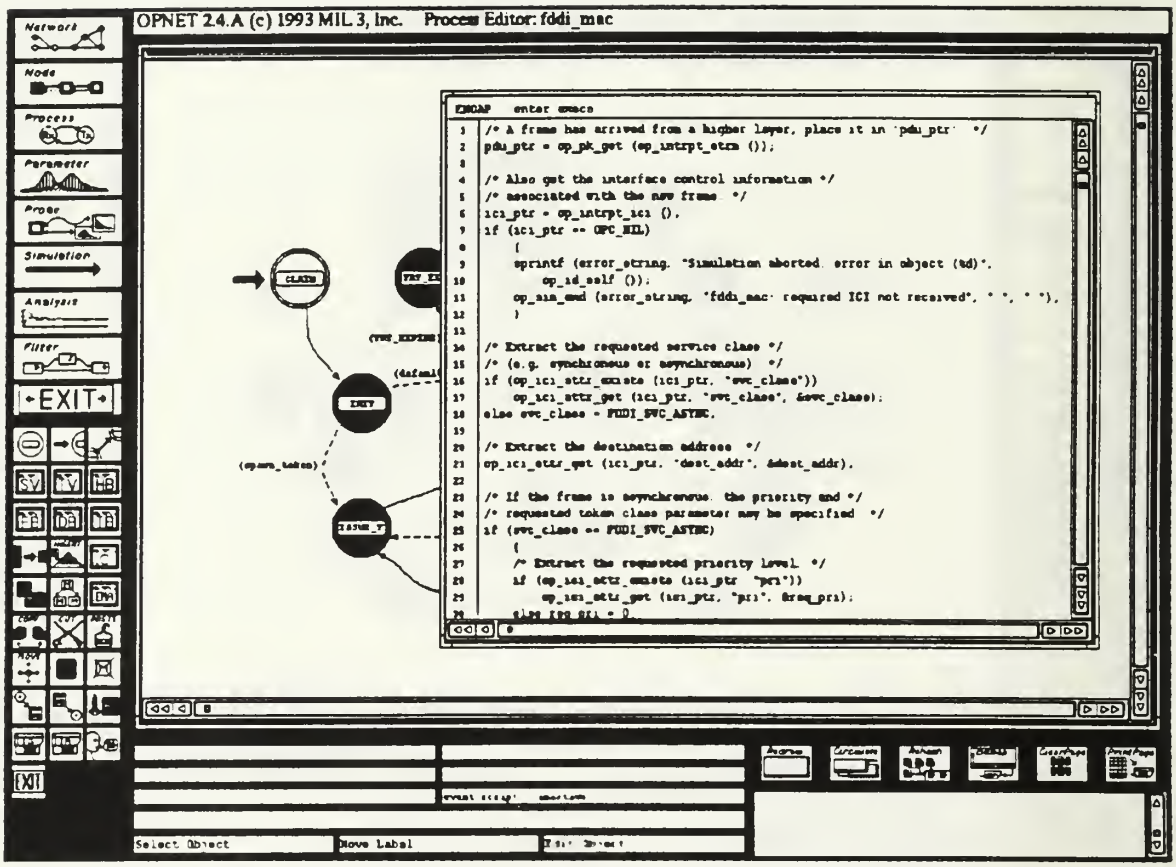

Figure 11. Process Model "fiddi_mac" with Text Editor

the model's behavior, and modify it if desired. Each state has its own section of code, and the icons to the left of the window include additional editors, all described in Vol. 4.0, *Tool Operations Manual*. These are primarily variable and function declaration sections. The entire code for the process may be called with the icon ".C," but this editor is for viewing only. Any changes made in the ".C" editor will remain when the editor is dismissed, but will disappear if the process is compiled. If calling the ".C" editor returns a "`.C file unavailable`" message, then recompiling will generate the file again (sixth icon above the lower "`EXIT`" icon in the Process Editor). Many UNIX stations include a cleanup command that deletes certain temporary files upon logging out of the system, and the "C" language codes ending in ".c" are not necessary once simulations are generated. They may always be recovered by recompiling the models. (In the file directory containing OPNET®, the model source codes have the suffix "`.pr.m`".) If changes are made (in a proper editor), the model must be recompiled. If several changes are being made within different sections, each may be saved with the keystrokes *<CTRL+S>*, deferring compiling to the end. (The set of manuals includes a summary page of OPNET®'s text editor keyboard commands.) If desired, the model may be saved without compiling by using the "`Write Process Model`" icon. When a process is changed, the station model in the Node Editor must also be called and written afresh. Then the corresponding network model must be called into the Network Editor, and be archived and bound again. If the modified process was not compiled earlier, then the "`C+A+B`" icon will compile all the process models in addition to archiving and binding them.

a. Source Node

The source node of the FDDI station model generates packets at a rate and size specified by the user. It also determines the destination address for each packet, the priority if applicable, and records the packet's creation time so that delay statistics can later be gathered. These data are passed to the MAC for encapsulation. In the Node Editor, the source is labelled "llc_src," and the process model is "fddi_gen."

In the source process' original form, as released by MIL 3, Inc. in version 2.4.A, the packet arrival (generation) rate is stochastically assigned on an exponential distribution approximating that specified by the user. If a precise, invariant arrival rate is desired, it may be assigned with the following change to the INIT state in the Process Editor, where "constant" is substituted for "exponential" in the line:

```
inter_dist_ptr = op_dist_load ("exponential", 1.0 /
     arrival_rate, 0.0);.
```

A voice traffic transmitter station, for example, would require a constant packet arrival rate from the source. Similarly, packet length is originally assigned a constant value in the given model, but may be set to a stochastic approximation of the requested value by replacing <"constant"> with <"exponential"> in the line:

```
len_dist_ptr = op_dist_load ("constant," mean_pk_len,
     0.0);.
```

If all stations on the ring are to have the same assigned attributes and the same source code (i.e., all <"constant"> or all <"exponential">), then the remaining steps

are to save and compile the process in the Process Editor, then call and save the station model in the Node Editor, and finally archive and bind the relevant LAN model in the Network Editor. If the stations on the LAN do not all have identical source code (i.e., some "constant" and others "exponential"), then the modified process models and their corresponding node models must be renamed. The following steps illustrate the creation of a station modified to allow a ring to simulate a number of voice stations amid other transmitters:

1. In the Process Editor, substitute "constant" for "exponential" in the "fddi_gen" model's INIT state editor. Save the change by keying <CTRL+S> while the cursor is inside the INIT state's editor.
2. Use the "Write Process" icon to save the modified process under a new name, for example fddi_gen_const.
3. Compile the new process model, then exit the Process Editor.
4. The Node Editor is used to create and save a new station model, by calling the original model and changing the "process model" field in the on-screen menu for the relevant node (in this example the "fddi_gen" process in the llc_src node is changed to "fddi_gen_const").
5. The new model is saved by invoking the "Write Node Model" icon. Exit the Node Editor.
6. In the Network Editor, the desired stations on the relevant LAN are reassigned using the on-screen menus: when the "model" attribute field is invoked at a

particular station, a list of available station models appears, and the desired one is chosen. If the expected model does not appear in the list, activate the User Interface Window "Rehash" icon to refresh OPNET®'s access to recently created files.

7. Desired stations are reassigned as required, and the LAN is saved, then re-archived and bound.

8. Differently named models using the same functions may cause naming conflicts at simulation run time. Should this occur, then the word "static" must be inserted in the Function Blocks ("FB" icon) of both the original and the new source model processes in the Process Editor, just prior to fddi_gen_schedule().

The above steps illustrate a change made to the source code, and do not represent the same situation in which identical stations are assigned different values in the given on-screen attributes menus.

b. Sink Node

The llc_sink node of the station model is the final destination of all message traffic. The INIT state establishes counters to hold statistical information regarding network performance (throughput and delay). The STATS state updates these counters as packets are received. The DISCARD state reports the statistical information at specified intervals, and finally destroys the packet. Because new packets are created for each transmission, they must eventually be destroyed when received, or the host computer conducting the simulation will soon fill its memory.

The "fddi_sink" model in the current version of OPNET® (Release 2.4.A, dated 02/27/93) is defective. It will cause the simulation to abort upon completion, with the error message "Program Abort: packet pointer is NIL," in the event any station did not receive traffic. Figure 12 illustrates the State Transmission Diagram as originally given, and the corrected version is shown in Figure 13. The user should correct the defective version, referring to Vol. 1.0, *Tutorial Manual*, "Bpt," pp.6-10. Saving this modification requires the same steps described for the source model, with the exception that no text has been changed. This defect and its correction are documented by Mil 3, Inc. as OPBUG 2070.

MIL 3, Inc. maintains an electronic bulletin board containing information on model corrections and upgrades between OPNET® revisions. Users may acquire these upgrades using file transfer protocol (ftp) procedures to download desired files. Appendix A includes a sample of dialogue used to acquire an upgraded file from MIL 3, Inc.

c. Medium Access Control (MAC)

The MAC process model encapsulates frames received from the source node for transmission to other stations, maintains token holding and token rotation timers, inspects received packets, decapsulates received frames, and determines token usability. Vol.8.1.0, *Example Models Manual*, "FDDI," provides a detailed description of the MAC process and of the functions of its component states.

The MAC model provided with OPNET® Version 2.4.A, (filename: "fddi_mac.pr.m") has been upgraded by MIL 3, Inc. and the newer model and

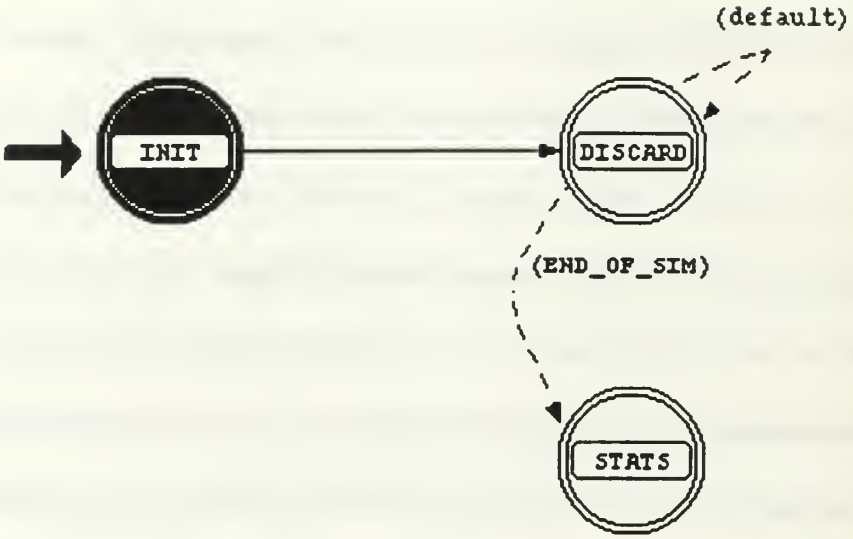

Figure 12. Original (Defective) Sink Process Model

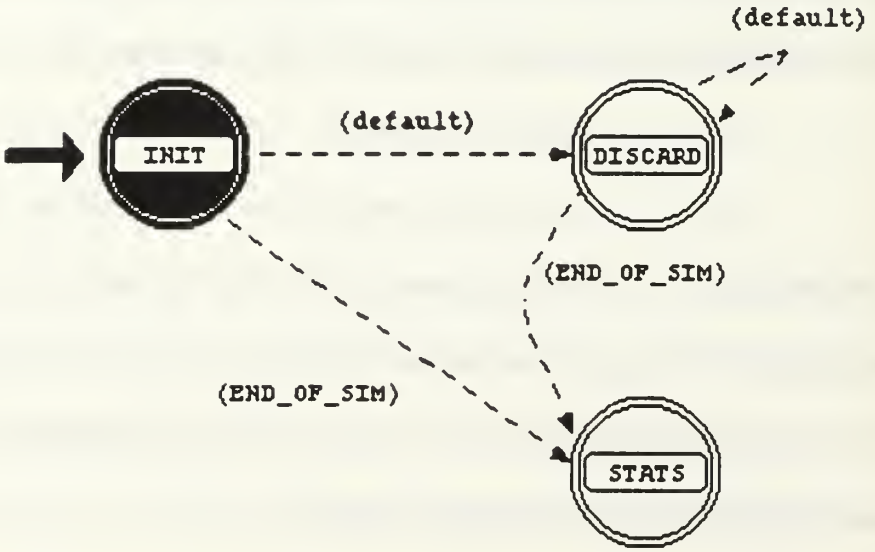

Figure 13. Corrected Sink Process Model

documentation are available via ftp on Internet, under the subdirectory "~/patches/2.4.B/opbug_2081". When retrieving a file via ftp, the user should verify the entire file is received by checking the file size listed on the bulletin board against the size of the file received. Entering <"*type image*"> at the ftp prompt should ensure a full and intact file retrieval. The original "fddi_mac.pr.m" file is then removed from the user's directory (~/op_models/fddi), stored in a safe place, and replaced with the newer version. The new file must then be compiled from the command window with the command, <"*op_mkpro -m fddi_mac*">, the procedure for which is described in Vol. 6.0, *External Interfaces Manual*, "Env". It has been observed that the "drag and drop" method of transferring files in using the File Manager in the SunOS Windows environment sometimes causes the subject file to gain or lose a byte or two, leading to "bitsum error" messages when said file is compiled in OPNET®. Standard UNIX commands are the most reliable method for moving OPNET® source files.

The patch is not necessary to operate the simulation; it is a refinement of the model, and will be included in the next revision of OPNET®. The patch OPBUG 2081 actually includes three repairs, documented as OPBUGs 2081, 2095, and 2097. OPBUG 2081 corrects existing timing and efficiency inaccuracies connected with the token acceleration feature, by which the token is destroyed and the simulation enters a "fast forward" mode in order to reduce the number of events while no station has a need to transmit. In a real FDDI LAN, TRT is reset each time the token passes, whether or not the token is used. The model in its original form allows the TRT timer to continue running when

the simulation enters "token acceleration," resulting in unexpected `Late_Count` occurrences. OPBUG 2095 is also related to the token acceleration feature, correcting the existing incorrect initialization of several variables when the simulation enters token acceleration. In particular, the variable `Fddi_Num_Stations`, the number of stations on the ring, is always reset to one, upsetting calculations predicting the proper location of the token at the end of an idle period. Finally, OPBUG 2097 addresses the fact that the original model neglects to properly account for the transmission delay associated with the token itself.

4. Packets

All communications between stations in a LAN and between the internal nodes of a station are conducted using data framed into packets. The Parameter Editor, described in Vol. 4.0, *Tool Operations Manual*, "Pm," is illustrated in Figure 14, which shows the packet structure `fddi_mac_fr`, which is used to encapsulate the frames generated in the Source node and sent to the `MAC`. Appendix B lists the five packet structures used in the FDDI LAN simulation, giving their fields and assignments. OPNET® simulation does not enforce limits on packet size required by the standard IEEE 802.5.

a. ICI Formats

Interface control information packets (ICI) are used for internal communication within a station, reporting for example service options, error conditions, and packet arrivals. Figure 15 shows the ICI Editor within the Parameter Editor, with the ICI `fddi_mac_req`. This ICI specifies the control information passed from the source to the MAC when transmission requests are generated. The ICI `fddi_mac_ind` specifies control information passed from the MAC to the LLC when a packet has been received by a station.

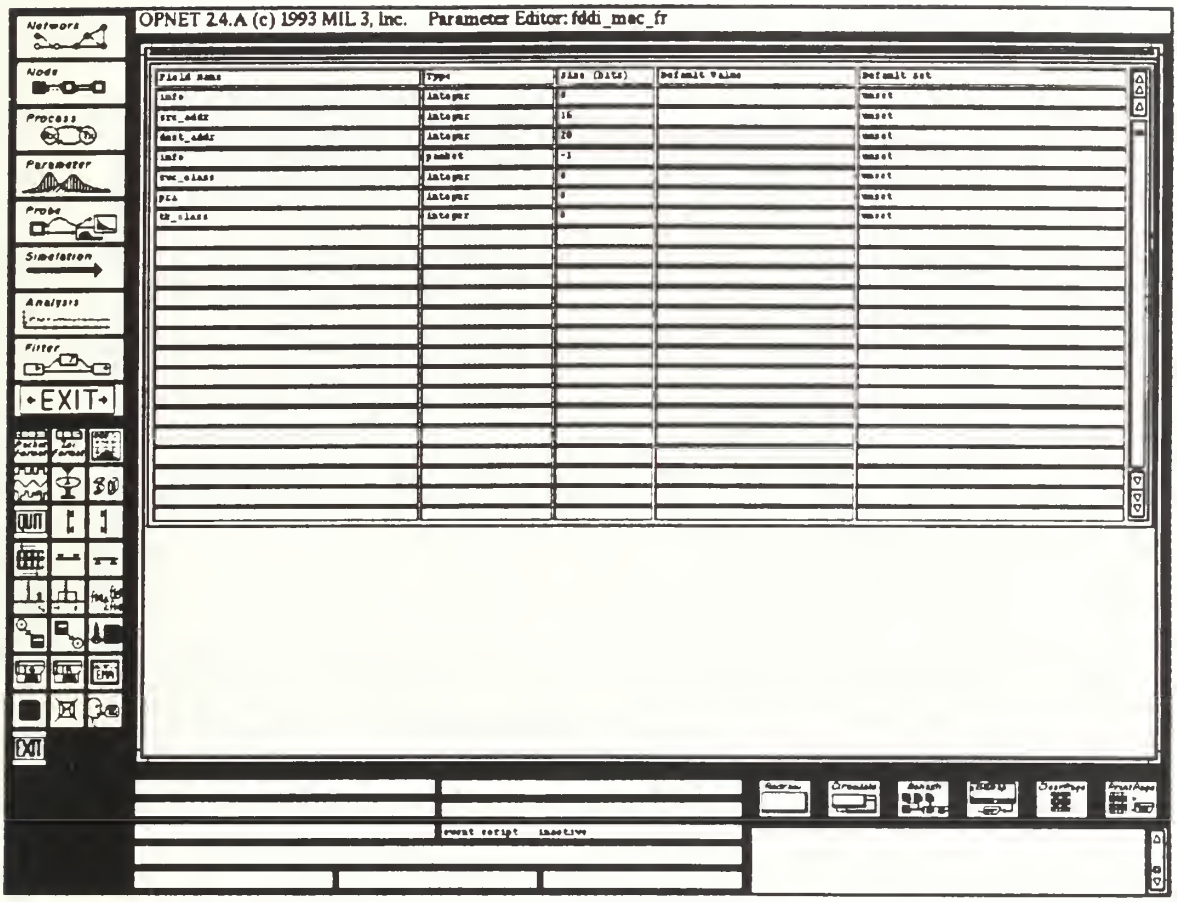

Figure 14. Packet Format "fddi_mac_fr" in the Parameter Editor

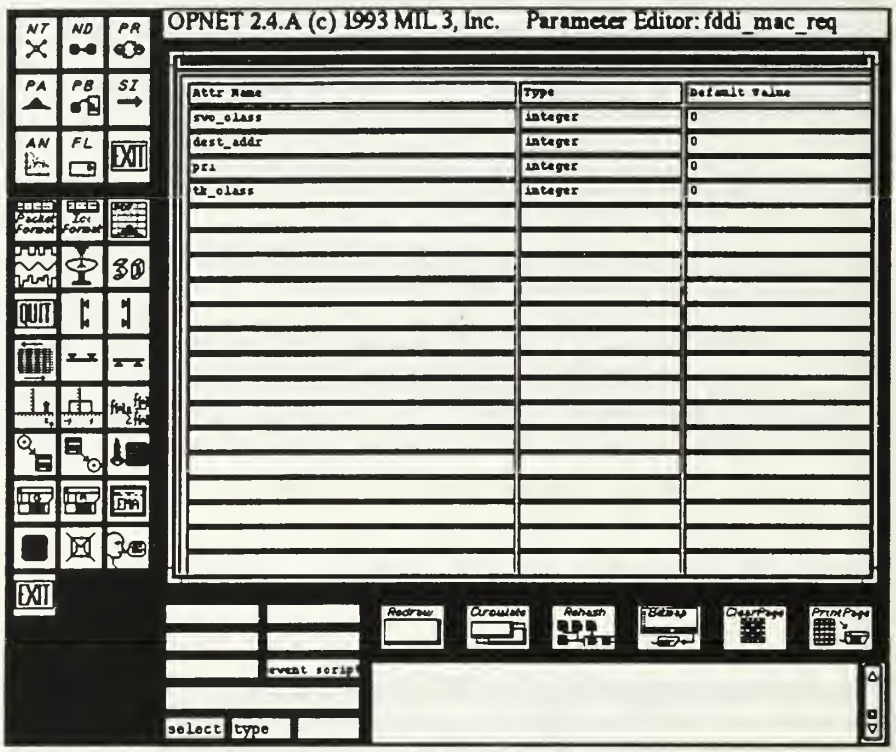

Figure 15. ICI Format "fddi_mac_rec" in Parameter Editor

For OPNET® simulation purposes, both structures are created once per station in the simulation , and reused as needed.

b. Packet Formats

Three types of packet frame formats exist in OPNET® to simulate communications between the stations. For the simulation these are created as needed and destroyed when no longer needed. Packets of format "fddi_llc_fr" are created in the LLC source as arrivals are generated. The format has only one field, containing the creation time, which is used to generate delay and throughput statistics when the packet is finally received at its destination address. The packet format "fddi_mac_fr" is used in the MAC state ENCAP to encapsulate the generated packets for transmission on the LAN. The "info" field contains the "fddi_llc_fr" structure providing the data of interest. Because OPNET® simulates only the characteristics of transmission and not the actions of stations in response to information received, the packets used are not precise replicas of real FDDI frames. The token is represented by the frame format "fddi_mac_tk". The field "fc" is inspected by each MAC process receiving a packet to determine whether it is a token or a message packet.

5. Environment File

Appendix C is an example of an environment (or configuration) file used to assign station attributes to a 32-station FDDI LAN. Inspection shows that the fields specified in the file correspond to the "promoted" fields in the on-screen attributes menus that appear in the Network and Node Editors, and to the fields in the Simulation Editor. Promoted attributes may be assigned directly within these editors, a tedious and error-prone process at best. If

a simulation is begun with none of these attributes specified, OPNET® will prompt the user for inputs at the command screen, another error-prone and tedious process. The environment file is the most efficient way to assign parameters of interest, and may be quickly modified between simulations, using the UNIX text editor. Vol. 6.0, *External Interfaces Manual*, "Env," discusses the environment file, and points out that attributes assigned in the environment file supersede those assigned in the Node and Network Editors. The file is recognized to OPNET® by its "`.ef`" suffix. The assigned attributes are described in chapter "FDDI" of manual Vol 8.1.0, *Example Models Manual*, and are presented here in their order of appearance for summary and in some cases for required elaboration. The attributes are not declared as variables in the "C" programs, but rather are generated by adding them to the on-screen menus in the Node and Network Editors. The user can create new attributes by adding them to the "`extended attributes`" field in the Node Editor, then including them in the environment file. This procedure is discussed in Chapter III. "Env," pp.33-34, in the *External Interfaces Manual* discusses the use of name wildcards in the attributes given below. Sequence of entries is not significant in the environment file.

a. "station_address"

This attribute is required for station identification; numbering of stations is from zero to `N-1`, where `N` is the number of stations on the ring. The INIT state in the MAC process calls this variable. Note that for quick changes to the file, lines may be commented with the pound key <#>.

b. "ring_id"

This attribute identifies the ring, in the event more than one may be modeled simultaneously. It is set to zero if there is only one ring.

c. "low dest address"

This attribute assigns the lowest identification address that may receive traffic from this station. The use of the wild-card asterisk, shown in Appendix C, assigns the same value to all stations. Quotation marks are used here because the attribute assigned has spaces vice underscore marks between words.

d. "high dest address"

This attribute assigns the high end of the range of addresses to which a station may send traffic. It is used in conjunction with the previous attribute by the "llc_src" process in the INIT state. It is permissible to limit the range of target addresses, down to one, but all target addresses must lie in a contiguous sector. For example, the code as given has no provision for allowing a particular station to send packets to two different stations without also possibly sending to the stations between them.

e. "arrival rate"

This attribute assigns the rate at which the source process will generate packets, and it is called by the INIT state It may be set to zero for any station intended to be idle. As originally used in the process code, arrival rate is a stochastic approximation exponentially distributed about the assigned value. To make this a precise unchanging value, as in the case of a synchronous voice transmitter, the code would have to be modified as described in the previous discussion of the source process.

f. "mean_pk_len"

Mean packet length is expressed in bits. Despite the name, this value is actually held constant by the INIT state of the source process model. The user may modify the code using the procedure described earlier to substitute <"exponential"> for <"constant"> in the line:

```
len_dist_ptr = op_dist_load ("constant," mean_pk_len,
    0.0),
```

which appears toward the bottom of the INIT state enter executives in the Process Editor. Then the station will generate non-identical packets, which may be more realistic behavior. OPNET® will permit any number of bits for the packet length; the user should know that FDDI packets have a maximum size of 36,000 bits.

g. "async_mix"

FDDI stations may generate synchronous and asynchronous traffic. This attribute sets the proportion, with 1.0 indicating all asynchronous traffic generated by the given station, 0.5 indicating half synchronous and half asynchronous. Any value between zero and one inclusive may be chosen. The INIT state in the `llc_src` model calls this attribute with the statement:

```
op_ima_obj_attr_get(my_id, "async_mix", &async_mix);
```

h. "sync bandwidth"

This attribute is used in the MAC process, INIT state, where it is expressed as a percentage of TTRT. It is analogous to SA_i in Equation 1, but is numerically a fraction of `T_Req` (TTRT), while SA_i is an amount of time. Synchronous bandwidth should not be confused with synchronous offered load. Bandwidth is expressed in time, while synchronous offered load is a bit transmission rate. Therefore, synchronous bandwidth is the time allotted for the transmission of synchronous offered load. It is entirely possible to set parameters so that these two attributes do not match well.

In describing `"sync bandwidth,"` Vol. 8.1.0, *Example Models Manual*, "Fddi", warns the user not to allow the sum of all assigned attributes `"sync bandwidth"` to exceed one, since OPNET® does not enforce FDDI protocol standards. However, this warning neglects to consider the physical delay parameters in Equation 1, which indicate that total synchronous bandwidth must be somewhat less than TTRT. The correct assignment of `"sync bandwidth"` involves some algebra, and is perhaps best explained with an example.

Given:	D_Max:	1.617 ms.
	F_Max:	0.360 ms.
	Token_Time:	0.00088 ms.
	TTRT:	8.0 ms.

Using Equation 1 yields:

$$8.0 \text{ ms.} - (1.617 \text{ ms.} + 0.360 \text{ ms.} + 0.00088 \text{ ms.}) = 6.02112 \text{ ms.} = \Sigma SA_i$$

This is 6.02112 ms. of total bandwidth to be divided among as many stations assigned. Assume there are five such stations:

$$6.02112 \text{ ms.} \div 5 \text{ stations} = 1.204224 \text{ ms} = SA_i$$

SA_i is converted to "`sync bandwidth`" with a division by TTRT:

$$1.204224 \text{ ms.} \div 8.0 \text{ ms.} = 0.150528$$

This is the value entered into the Environment file.

i. "T_Req"

This attribute is called in the MAC process, INIT state, and represents the specified station's requested value of TTRT. A real FDDI LAN has a TTRT negotiation phase; OPNET® simply chooses the smallest `T_Req` value available. The user may, but need not set different values to each. This value is in units of seconds, which is not apparent from the manuals nor from the default value that appears at the command prompt if no value is assigned.

j. "spawn station"

The spawn station is simply the starting point for the token, and may be assigned to any station on the LAN.

k. "station_latency"

This is the delay incurred by packets as they pass a station's ring interface. Powers gives 1 μsec. (1993, p.336); 60.0e-08 sec. agrees with Dykeman and Bux (1988). Station latency is a component of D_Max in Equation 1.

l. "prop_delay"

Propagation delay is the time separating stations on a ring, based on the amount of fiber between them. It is given here in seconds, and may be used to define the ring size. The INIT state of process `fddi_mac` calls this value, which is used as one of the delay

parameters applied to transmission commands. FDDI standards limit the ring size to a maximum of 200 miles of fiber (Dykeman, 1988, p. 997), and OPNET® assumes that the fiber length is divided evenly among the stations. That is, all stations are evenly spaced on the LAN in OPNET® simulations, whatever the number of stations and length of fiber. Dykeman and Bux (1988, p. 1000) define propagation delay in units of time per distance, and give a value of 5.085 μs/km. The value given in the original example environment file, 5.085×10^6 seconds, corresponds to one kilometer of fiber between stations.

m. "acceleration_token"

This attribute speeds the simulation by removing the token during idle periods when no station has packets to transmit, significantly reducing the number of events.

n. "duration"

This is the simulated run time in seconds. Most systems should reach steady state in less than a second.

o. "verbose_sim"

This feature enables on-screen reports regarding event numbers, time remaining until completion, etc.

p. "upd_int"

This specifies in seconds the intervals at which to make on-screen simulation status updates. It must be less than duration to be useful.

q. "os_file"

The output scalar file receives scalar data accumulated over several simulations. It is useful in observing the effect of varying one or more attributes, for example TTRT, over a series of experiments.

r. "ov_file"

The output vector file receives throughput and delay information relevant to one simulation run. Output vector files can be quite large, on the order of several megabits, and for this reason are often automatically deleted by a <"cleanup"> command included in a UNIX station's logoff sequence. The user should alter the filename or save desired plots as ".ac" files using the Analysis Tool, rather than log off planning to study the vector data at some future time.

s. "seed"

This is a constant used by the simulation's random number generator. It may be any positive integer, but should be left constant once chosen.

t. "debug"

This enables the Debug Tool, allowing the user to step through a simulation one event at a time. Once enabled, the command <"help"> provides a listing of the debugger's features.

D. SIMULATION

This section presents the steps involved in running a simulation and observing the resulting output data. The user must keep in mind that OPNET® is unaware of IEEE 802.5. That is, it is the user's responsibility to ensure reasonable input parameters are assigned in keeping with the established standards. The steps given will use the original model provided.

1. Build the LAN

If a 32-station LAN is not already available in the Network Editor, then one should be created using the command `< "fddi_build.em.x 32" >`, as described earlier.

2. Correct "OPBUG 2070"

The simulation will abort if the original process model `fddi_sink` is used for the `SINK` node `llc_sink`, and some station happens to have not received any packets. The correction described earlier should be applied, and the model recompiled and saved.

3. Implement "OPBUG 2081" Patch

As described earlier, this repair corrects minor timing inaccuracies in the model, related to the token and to the token acceleration feature. The simulation will work without aborting if the patch is not applied, but the user planning on implementing code changes within the model over the long term should patch the model before doing so.

4. Update Environment File

Refer to the configuration file in Appendix C for input parameters. As mentioned before, use of this file will save the user the effort involved in setting parameters by hand through the Node, Network, and Simulation Editors. Note that attributes assigned in the

Environment file will supersede any that are assigned through these editors. The file should be given some distinguishing name, for example "fddi32.ef".

5. Generate Probe File

Use of a Probe file is optional. The process code as written will generate vector file outputs only for overall throughput, delay, and mean delay. Additional outputs may be monitored through the use of a Probe file, illustrated in Figure 16. Vol. 4.1, *Tool Operations Manual*, "Pb," describes the Probe Editor. With it the user may monitor, for example, each station's arrivals (packet generation) and throughput at any physical interface point on the LAN, measured in packets and/or in bits per second. For simplicity, only packet arrivals at station fl1 will be assigned a probe in this simulation.

6. Simulation Editor

Figure 17 illustrates the Simulation Tool in the user interface window, with the settings necessary to run this simulation. Use of this tool is discussed in Vol. 4.1, *Tool Operations Manual*, "Sm" Fields are assigned by use of the cursor, and are filled by choosing from on-screen menus or keyboard entry. The "Simulation" field should be assigned the LAN filename, <*fddi_32_net*> (note that filename suffixes are not visible to the user in the various Editors) The fields "Probe File", "Vector File", "Scalar File", "Seed", "Duration", and "Upd Intvl" may all be left blank if they are assigned in the configuration file; "Probe File" is optional in any case. The "Arg Name" field should be assigned <*environment file*>, and the "Arg Value" field should be assigned the <*filename*> given to the environment file. The user may then save the work area using the

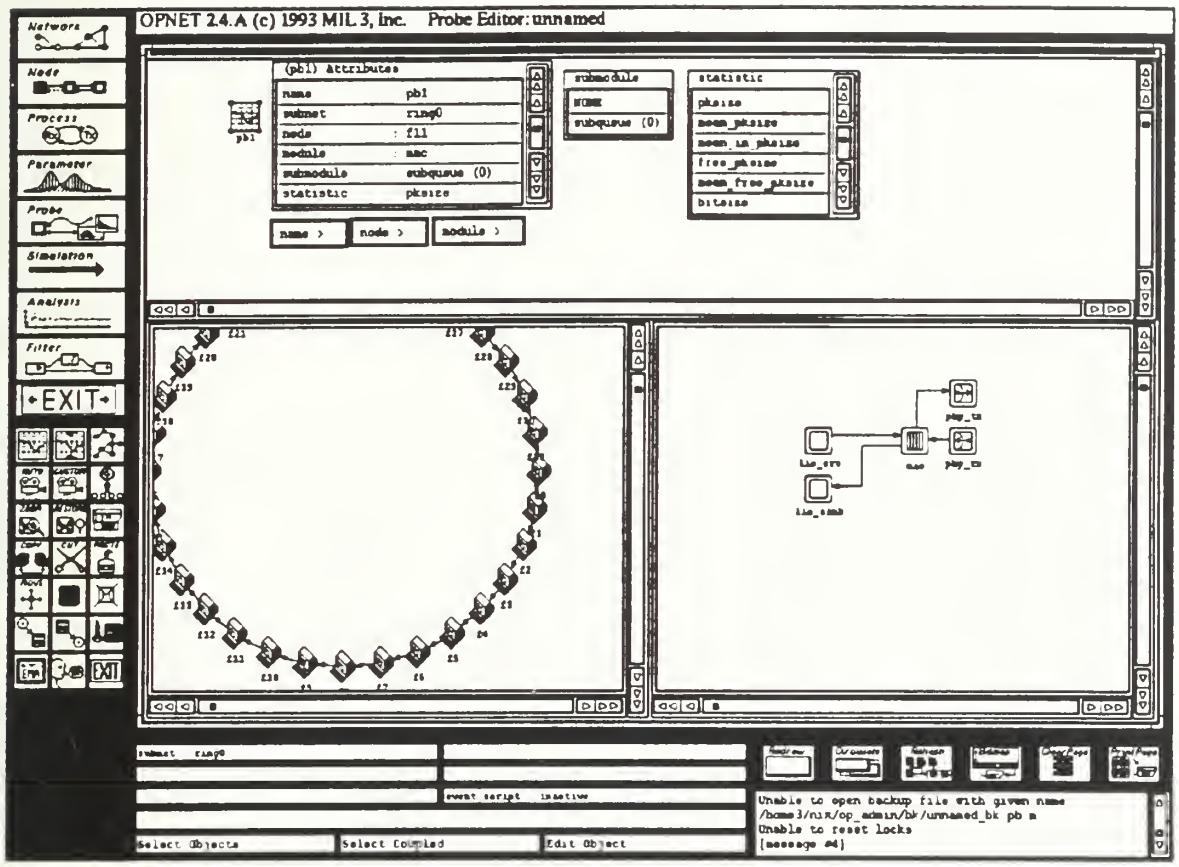

Figure 16. Probe Assignment

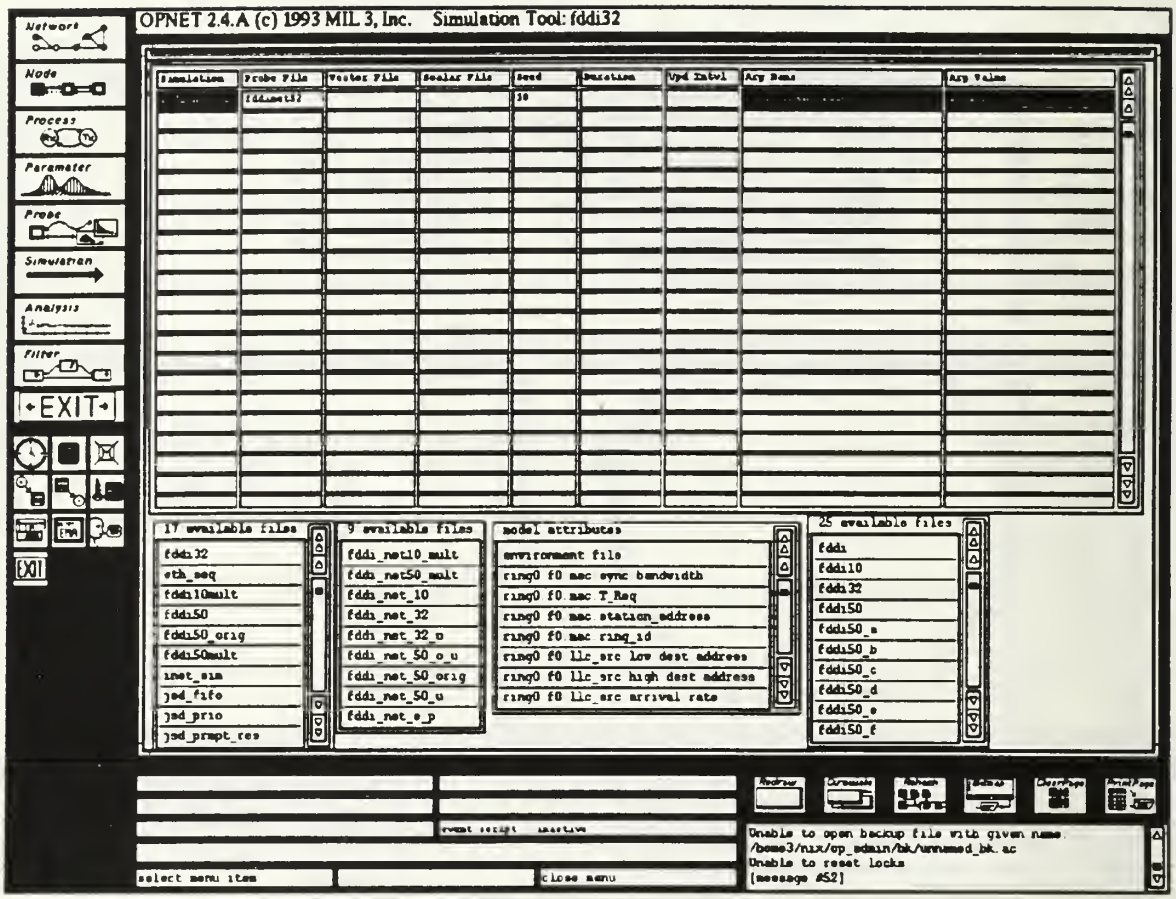

Figure 17. Simulation Tool

"Write Simulation" icon, to spare the effort of filling in these fields again on future simulations.

7. Start the Simulation

Once the fields are set, the simulation is started with the "Execute Simulation Sequence" icon. Had the user neglected to assign some parameter, the simulation will wait until the command line prompt has been answered; the user should keep the command screen in view. Upon completion of the simulation, a vector file (suffix ".ov") will be generated, along with a scalar file (suffix ".os"). The "Rehash" icon must be invoked to refresh OPNET®'s access to the files. Then the user may exit the Simulation Tool and enter the Analysis Tool.

8. Analysis

Figure 18 shows the Analysis tool in the user interface window, whose operation is explained in Vol. 4.1, *Tool Operations Manual*. The first action upon entering this tool is to call the available vector outputs, using the "Open Output Vector File" icon, then selecting from the choices presented. If more than one are present, choose the one that was assigned in the environment file. The on-screen menu will then disappear, leaving the user to select the "Create Single Vector Panel" icon, which presents the on-screen menu shown in Figure 18. The entries "end-to-end delay (sec.)", "throughput (bits/sec)", and "mean delay (sec.)" are generated directly from the SINK process model. The remaining field, "ring0.fl1.mac[0].pksize," comes from the Probe Editor. Each may be plotted by selection with the cursor, then dragging the box corners to the desired

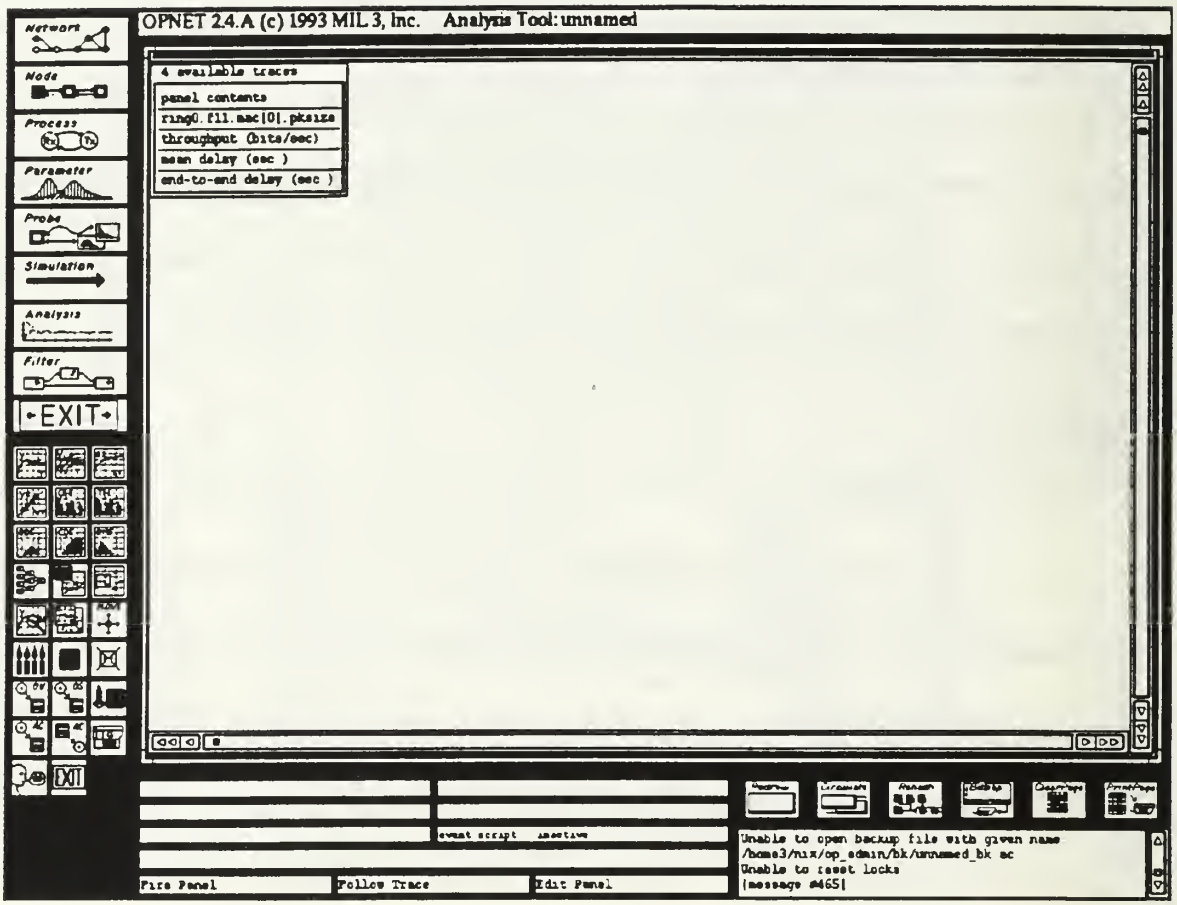

Figure 18. Vector Trace Selection in the Analysis Tool

size. Having placed the panel, the plotted points are fired by clicking the left mouse button, or by placing the cursor over the "Fire All Panels" icon and clicking the same button. Figure 19 shows all four plots generated, placed together on the screen. The "Create Multi-Vector Panel" icon is used to place several plots in the same panel, an operation that is meaningful when the Probe Editor is used to generate comparable outputs. Once the desired plots are on the screen, they may be saved with the "Write Analysis Configuration" icon, which will store the plots in a ".ac" file for later recovery. This is important because the vector file will be written over the next time a simulation is run using the same output vector filename. In addition, the UNIX station's logout sequence may include a <"*remove *.ov*"> command to prevent the accumulation of large vector files in memory.

The output scalar file, on the other hand, accumulates steady-state data over several simulations, allowing the generation of plots showing, for example, the effect of various TTRT values on total throughput. The user wishing to create su a plot should ensure the file is empty of previous data before commencing a series of simulations.

9. Debug Tool

The debug tool may be activated from the environment file. The command <"*help*"> will list the available commands. The user may step through a simulation one event at a time, or specify stopping points. The <"*fulltrace*"> command causes every variable to be reported at each event, allowing the user to follow the sequence of events in a

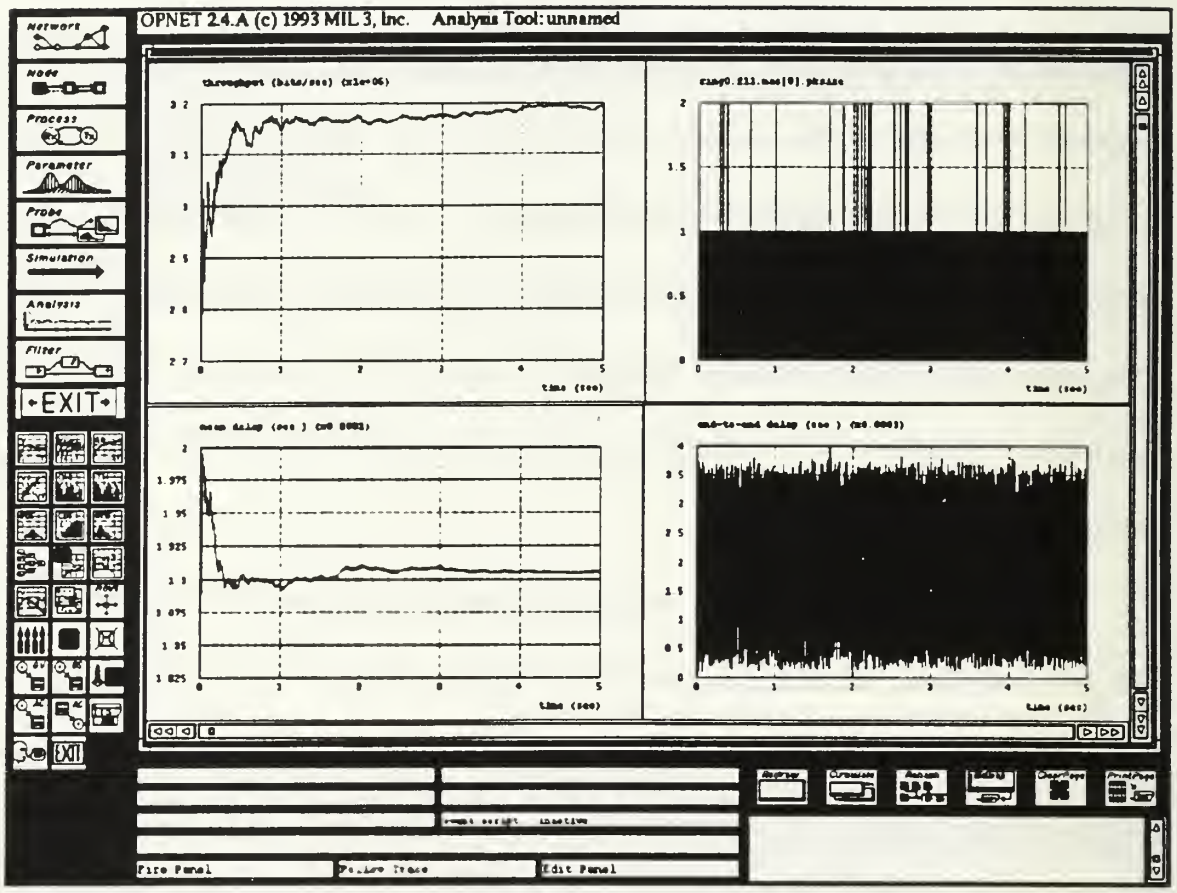

Figure 19 Analysis Tool Display, Four Panels

simulation, and to search for logic errors should failures occur. Appendix D is a short section of the debugger's output when the <"fulltrace"> command is active.

III. MODEL MODIFICATIONS

A. OVERVIEW

The FDDI station model provided with OPNET® is shown by Schenone (1993) to perform as expected by performance equations provided by research literature, for example, Dykeman and Bux (1988). However, the model as given lacks the flexibility to adequately demonstrate the characteristics and metrics required to formulate recommendations for the development of a CDL network interface. To begin, no way exists in the original model to monitor the throughput and delay statistics of synchronous traffic separately from asynchronous traffic. In addition, the code must be modified to allow the implementation of different asynchronous priority levels, and further altered to allow the generation of statistics of these subcategories. Also desirable is a method to hold traffic in the sink process of one station on the LAN, rather than destroying all packets, so that a bridging protocol may eventually be developed for communication between LANs, which is the ultimate goal of the Common Data Link Project. Finally, a multicast/broadcast facility, by which a packet may be addressed to multiple stations, is needed. Modifications implementing these features were generated for this thesis, and are described in this chapter. Appendices E, F, and G contain the final form of the "C" programs representing the process models "`fddi_mac_mult`", "`fddi_gen_mult`", and "`fddi_sink_mult`", respectively. All contain inserted comments to indicate where changes have been made. In some cases, modifications are extensive

enough that the original structure is not apparent. For these cases the reader who is interested in comparisons is referred to the original code available in the Process Editor.

The modifications described here were suggested in large part from readings in the research literature. In particular, Tari, Schaffer, Poon and Mick (1991) published results generated from another commercially produced network simulation tool, the Block Oriented Network Simulator (BONeS®), demonstrating that increased asynchronous offered traffic load has minimal effect on the throughput of synchronous voice data traffic. At the same time, the throughput of the various asynchronous levels was shown to degrade in order of priority with increasing asynchronous offered load. These findings emphasize the fact that OPNET® has no facility in place providing for the monitoring and plotting of throughput or delay data in respect to class or priority levels.

Closely interrelated are the setting of asynchronous priority levels, a system of subqueues to segregate traffic by priority, and the gathering and display of performance statistics according to priority. In the following discussion, reference is sometimes made from one to the others before all are complete.

B. PRIORITIZATION

1. Activating Prioritization in OPNET®'s FDDI Model

As given in the original released model, code exists to support a prioritization scheme, but it is not implemented. The station model includes a priority field that may be set in the Node Editor, but no setting will take effect until the INIT state in the MAC process is

modified. As given, T_Pri[i] is simply assigned the value <*Fddi_T_Opr*> (which is the negotiated value of TTRT, with the negotiation consisting of selection of the lowest requested T_Req from all stations) for all priority settings i, resulting in no distinction made between priorities. The state TX_DATA in the MAC process model contains the code that determines transmission eligibility, then transmits packets if timing conditions are satisfied. The user should notice here that unlike the real FDDI protocol, the Token Holding Timer is incremented from zero to tht_value (THT), not decremented from THT to zero. This results in a reorientation of T_Pri[i] settings, wherein higher settings allow a larger transmission window, and therefore, higher priority. To re-emphasize: in the OPNET® model, T_Pri[i] is larger for higher priority stations. Actual settings of T_Pri[i] are a matter of user's choice and real-world physical characteristics. One quick approach is to alter the code in INIT from the original:

```
for (i = 0; i < 8; i++)
{
     T_Pri [i] = Fddi_T_Opr;
}
```

by substituting the text:

```
T_Pri [i] = (double)Fddi_T_Opr/(8-i);
```

to impart some weight to the priority settings (Appendix E, line 250). Note that priority settings in OPNET® are counted from zero to seven, in keeping with the "C" programming

language convention of numbering elements of an N-element array from zero to N-1. As mentioned earlier, actual FDDI convention numbers the priority settings from one to eight.

2. Changing the Scheme and the Code

For the purposes of CDL, an ability in a station to generate traffic of differing priorities is a desirable characteristic. The modifications discussed here allow this behavior, though with a certain amount of abstraction included. Essentially, each packet generated in the source process is assigned a priority setting, in a manner that is functionally identical to the determination of the destination address. However, the priority of one packet has no influence on the priority of the next one generated at the same station. Of course in real-world transmissions, packets are grouped into messages, meaning that thousands of consecutive packet arrivals should have the same priority settings to reflect real behavior. In fact, the subqueue structure imparted to the `MAC` causes outgoing packets to be sent in decreasing order of priority, thereby modeling expected behavior to some small degree. More significantly, the user should keep in mind that the model's purpose is to model a LAN's handling of the traffic it does receive. The fact that packets are transmitted with random priorities in a scattershot fashion is not significant to the LAN's overall performance.

3. Subqueues

Because subqueues are essential to the development of the prioritization scheme, their construction is addressed first. As is seen in the FDDI station model in the Node Editor, the MAC node is represented as a queue. Therefore, only a change to "`subqueue count`" field in the attributes menu is necessary to change the MAC's structure into a bank of subqueues. Code is already in place that treats the MAC as a set of subqueues, although by

default, only one is available at first (this is labelled <subqueue (0)>, as is seen when the "subqueues" field is selected from the on-screen attributes menu for the MAC node. Nine subqueues are desired here: one for each asynchronous priority setting, plus another to handle synchronous traffic. Subqueue indexing corresponds to priority settings, so that subqueue (0) receives and releases the lowest priority asynchronous traffic, while subqueue (7) is assigned the highest priority traffic. Subqueue (8) is designated for synchronous traffic. This segregation of traffic into subqueues is necessary to support the recording of performance statistics and plots of traffic generation through the use of the Probe Editor. Also, while the Kernel Procedures (KP) available to the user include one that allows packets to be removed from the transmission queue in order of priority, rather than in the usual first-in-first-out (FIFO) order, subqueues allow simpler logic (Vols. 5.0 and 5.1 are directories of the commands and functions used by OPNET®). Vol. 2.0, *Modeling Manual*, "Nddef," pp. 27-29, describes the procedure to adjust code so that references to queues may be replaced with references to subqueues, particularly in relation to prioritization schemes. In summary, the subqueues represent a way point for packets. They are created and assigned a priority setting in the source, encapsulated for transmission in the MAC (ENCAP state), and placed in the appropriate subqueue while the station awaits the next token arrival to transmit them.

a. "RCV_TK"

In the RCV_TK (receive token) state, the first test of token usability is a determination of the presence of outbound traffic. The statement:

```
if (op_q_stat (OPC_QSTAT_PKSIZE) > 0.0)
    ... etc ...
```

calls for an inspection of the queue. This statement is replaced with a loop structure that searches all subqueues:

```
for (i = NUM_PRIOS - 1; i > -1; i--)
{
     if (op_subq_stat (i, OPC_QSTAT_PKSIZE) > 0.0)
     ... etc ...
}
```

(Appendix E, line 453)

Of course the above requires a declaration of the variable `NUM_PRIOS` and the loop counter `i`.

b. "TX_DATA"

The TX_DATA (transmit data) state contains the code that transmits packets while the token remains available, and monitors THT. The THT (`tht_value`) is checked inside a loop whose condition is, "while packets remain in the queue, transmit." This loop contains most of the code in TX_DATA. This condition must be set inside another loop which counts through each desired subqueue, and which must include an additional number of "break loop" points. Accordingly, the code:

```
while (op_q_stat (OPC_QSTAT_PKSIZE) > 0.0)
 {
     /* Remove the next frame for transmission.*/
     pkptr = op_subq_pk_remove (0,OPC_QPOS_HEAD);
          ... etc ...
}
```

is rendered into the following:

```
for (i = NUM_PRIOS - 1; i > -1; i--)
{
     while (op_subq_stat (i, OPC_QSTAT_PKSIZE) > 0.0)
     {
          /* Remove the next frame for transmission.*/
          pkptr = op_subq_pk_remove (0,OPC_QPOS_HEAD);
          ... etc ...
     }
}
```

(Appendix E, line 920)

The transmission loop is broken when any of the following occur:

- No more packets are enqueued.
- In the case of synchronous transmission, insufficient bandwidth remains to complete a transmission. (Note that synchronous traffic is allocated by the user an inviolate amount of time in which to transmit, regardless of the lateness of the token. However, the model checks the bandwidth remaining to ensure a transmission can be completed within the allotted time, and will not commence a transmission that would delay the token. This is in agreement with the actual protocol, and in contrast to the asynchronous case, in which packet transmission may commence while the THT is active, even if the transmission will keep the token past THT expiration.)
- The remaining packets are of too low a priority to be transmitted in the time remaining to THT (`T_Pri[i] < THT`).

After closure of the outer loop, the station deregisters its interest in the token by indicating it has no more traffic to send. This information is used by the token

acceleration mechanism, which will destroy the token for the time period no station has traffic to transmit, then recreate it when needed, thereby significantly reducing the number of events and the amount of time spent in a simulation. The original code appears at the bottom of the TX_DATA state:

```
if (tk_ .gistered && op_q_stat (OPC_QSTAT_PKSIZE) == 0.0)
{
     tk_ egistered = 0:
     fddi_tk_deregister ();
}
```

As before, this must be altered to search through a set of subqueues first before deciding no traffic remains to be sent:

```
q_check = 1;
for (i = NUM_PRIOS - 1; i < -1; i--)
{
     if (op_subq_stat (i, OPC_QSTAT_PKSIZE) == 0.0)
     {
          q_check = 0;
     }
     else
     {
          q_check = 1;
          break;
     }
}
if (tk_registered && q_check == 0)
{
     tk_registered = 0;
     fddi_tk_deregister ();
}
```

(Appendix E, line 1084)

4. Modifications to Prioritization

In order to enact the priority scheme described above, changes are needed in the station model (Node Editor), all three process models (Process Editor), packet format (Parameter Editor), and to the Environment file (UNIX text editor).

a. Station Model Changes

For the modified priority scheme, new attributes are needed in the on-screen menus that appear for the `mac` node of the `fddi_station` model in the Node Editor. The original field "`priority`" is left in place, but not used. Note that the "`super priority`" field (described in Vol. 6.0, *External Interfaces Manual*, "Rel," p. 15) is not related to the FDDI protocol, but is a tool for scheduling of events in the simulation; it is not used, and should be left disabled.

In order to support the priority setting protocol that occurs in the source model (described later), two new attributes are created: "high pkt priority" and "low pkt priority." The new fields are created as listed in the following steps:

1. Call the on-screen attributes menu for the "mac" node.
2. Place the cursor over the "`extended attrs`" field, and press the left mouse button to call the submenu (the right mouse button dismisses the attributes menu; try again).
3. Assign the fields as shown in Figure 20: names will be "`low pkt priority`" and "`high pkt priority`", units are <*none*>, type is <*integer*> (selected from another on-screen menu that will appear,

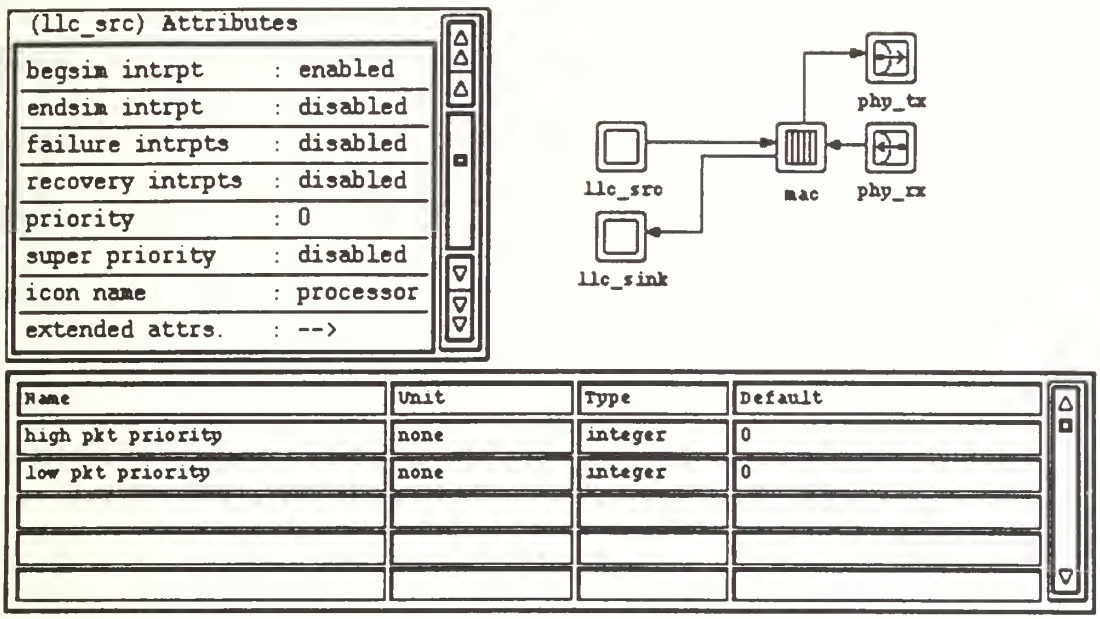

Figure 20. Adding Extended Attributes to the Station Model

rather than typed in), and defaults are at the user's discretion. Zero for both are reasonable. These assignments are preserved with the keyboard combination <CTRL+S>. Further, the model must be saved using the "Write Node Model" icon, as described in Vol. 4.0, *Tool Operations Manual*, "Nd," pp.13-15.

4. Once the model is saved, then called again to the Node Editor, unexpected values will probably appear in the fields corresponding to the new attributes (which will now appear in the primary on-screen menu as well as in the "extended attrs" submenu). These values should be set to <promoted> so that they may be assigned in the Environment file. To set the newly created field to <promoted>, place the cursor over the field, then type <CNTL+O> at the keyboard, invoking the literal "promoted."

This last item is not described in the manuals; it was obtained from MIL 3, Inc.'s technical support via electronic mail.

b. Environment File Changes

Corresponding to the attributes created and then assigned <promoted> fields above, the following code is added to the Environment file:

```
"*.*.llc_src.high pkt priority" : 7
"*.*.llc_src.low pkt priority" :  0
```

The quotation marks are required here because of the spaces in the attribute names. Had `low_pkt_priority` been used instead, then the quotation marks would have been omitted. Examples of both styles appear in Appendix H, an example Environment file that includes all

attributes added to the model (some of which are still to be described). The Environment file may be modified in the UNIX text editor. Pound signs (#) indicate comments. Order of attributes is not significant, and any attributes not used by a model are simply ignored.

c. Source Modifications

The approach to assigning a random priority to each packet as its arrival is generated is functionally identical to the procedure by which the destination address is generated. The OPNET® kernel function `op_ima_obj_attr_get()` is used to call attributes from the node model or from the Environment file. The KP `op_dist_load()` is used to load a distribution to be used in generating a stream of stochastic values. These are used in the source process model "`fddi_gen`" INIT state in the following manner:

```
op_ima_attr_get(my_id,"high pkt priority,
     &high_pkt_priority);
op_ima_attr_get(my_id,"low pkt priority"
     &low_pkt_priority);
pkt_priority_ptr = op_dist_load ("uniform_int",
     low_pkt_priority, &high_pkt_priority);
```

(Appendix F, line 108)

In the preceding, the first two lines call the desired attributes from the Environment file to the calling station ("`my_id`"). The second field in the procedure call is taken verbatim from the Environment file, while the third field is the address of the attribute. The address may have any name; <*&high_pkt_priority*> is purely a memory aid for the user, and is not required by syntax to resemble the field name to which it is assigned. The value returned by `op_dist_load()` is used in the ARRIVAL state to finally generate the priority setting with the statement:

```
pkt_prio = op_dist_outcome (pkt_priority_ptr);
```

(Appendix F, line 195)

which finally returns an integer between the values set in the Environment file. This integer is assigned to the packet with the commands:

```
op_pk_nfd_set (pkptr, "pri", pkt_prio);
op_ici_attr_set (mac_iciptr, "pri", pkt_prio);
```

(Appendix F, line 214)

Where only one priority setting is desired for a particular station, the attributes need only be both set to that desired value, with the station specified in the Environment file. A limitation here is similar to that of addressing: any section of values from zero to seven may be chosen, but it is not possible, for example, to assign a station the asynchronous priority settings two, four, and six.

In the SV (State Variable) edit window, `high_pkt_priority` and `low_pkt_priority` are declared as integers, and `pkt_priority_ptr` is declared as a pointer of type "Distribution." Because OPNET® uses a form of proto-C, the declarations made in the editor actually have the following form:

```
Distribution*  \pkt_priority_ptr;
int            \high_pkt_priority;
int            \low_pkt_priority;
```

(Appendix F, line 9)

When the code is compiled and the ".C" icon is invoked to present the entire process model code, the above will have the following appearance:

```
Distribution*  sv_pkt_priority_ptr;
int            sv_high_pkt_priority;
int            sv_low_pkt_priority;
```

OPNET® will also produce the following in the ".C" code:

```
#define pkt_priority_ptr
     pr_state_ptr->sv_pkt_priority_ptr
#define high_pkt_priority
     pr_state_ptr->sv_high_pkt_priority
#define low_pkt_priority
     pr_state_ptr->sv_low_pkt_priority
```

(Appendix F, line 44)

The above commands have the effect of choosing, on a uniform distribution, a priority setting from a given range whose endpoints are retrieved from the Environment file.

d. MAC Modifications

Because a priority scheme is already supported in the original model, little change is required in the `mac` node once a priority value is assigned in the source node. The user should note that communication between nodes is conducted with locally held variables; globals are avoided. This may result in different declared variable names for the same data,

which is acceptable. Thus, the `pri` of the source becomes the `req_pri` of the `mac` node. In the `MAC` node, priority settings are used as the indexes for the subqueues. `NUM_PRIOS` is declared for use as a loop counter.

e. Sink Node Modifications

The changes to the "`fddi_sink`" process are nearly all related to the generation of performance parameters, which are discussed in the next section. In the original model, the received packet's priority setting is not even relayed to the sink, since the only information necessary to the calculation of overall throughput and delay statistics are the packet's creation time and its time of receipt. A fundamental addition to the code in the DISCARD state is the line:

```
op_pk_nfd_get (pkptr, "pri", &pri_set);
```

(Appendix F, line 78)

which recovers the priority from the field "`pri`" in the frame structure `fddi_llc_fr`. With this information, additional modifications will bring about the ability to create throughput and delay information for each asynchronous priority setting, and also to separate synchronous traffic statistics from asynchronous. But before any of this will work, the `fddi_llc_fr` packet structure must be modified.

f. Packet Format

The frame that is created in the source and passed to the MAC, then encapsulated into a more extensive frame, then ultimately passed from the destination station's MAC to the sink process for accounting and final destruction, is of the format `fddi_llc_fr`.

To support the prioritization scheme, the format needs to include more information than only the frame's creation time. To enhance its characteristics, the Parameter Editor's "Packet Format" icon is invoked, and the format `fddi_llc_fr` is called. Another line is added, as shown in Figure 21, making "`priority`" an attribute of the packet. Type is set to "`integer,`" size can be "0," default value is "0," and default set is "unset." The changes are saved with the "`Write Model`" icon.

C. PERFORMANCE MEASURES

1. Overview

OPNET®'s original FDDI LAN model provides no ready way to monitor synchronous traffic separately from asynchronous, and no way to monitor the throughput and delay statistics for individual asynchronous levels. The inability of OPNET®'s original FDDI model to provide anything besides overall performance is a serious limitation to its usefulness in the CDL project. A major goal of this work is to augment the code in the sink process so that additional output vectors are generated, allowing the effects on individual class and priority levels to be seen. For example, in the original model, the user may assign any desired proportion of the generated traffic to be synchronous, but the original model has no facility to measure the synchronous traffic alone. The following paragraphs describe the modifications made to allow the display of statistics segregated by class and priority.

As an incidental note, the user should realize there is no particular significance to the sequence in which the states appear in the ".C" file. That is, the order DISCARD,

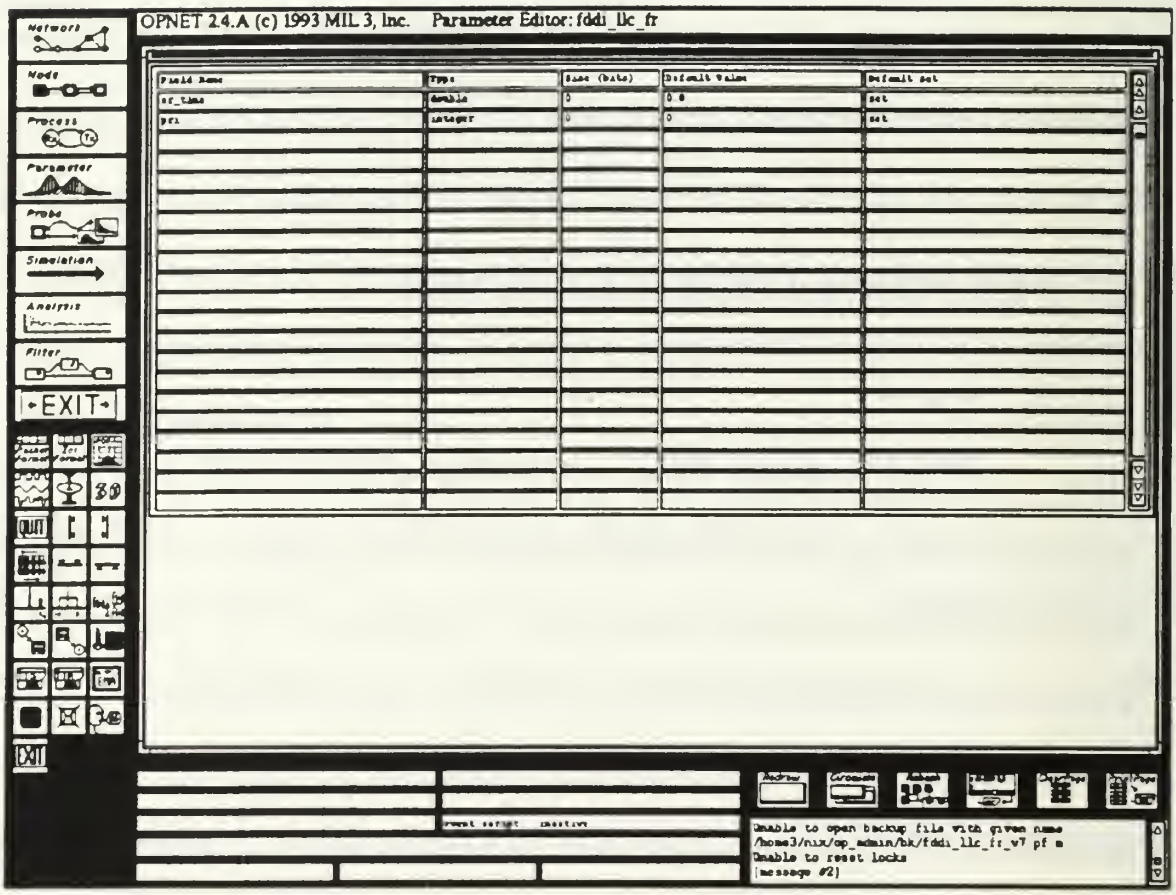

Figure 21. Adding a New Field to the "fddi_llc_fr" Frame Format

STATS, then INIT that appears in the "fddi_sink" process code is no indication of the sequence in which the simulation "visits" these states for each station. In fact, the more logical order, INIT, DISCARD, then STATS will be followed in this discussion, though preceded by a discussion of the variables needed. Appendix G is the file "fddi_sink_mult.pr.c", containing the modifications described here.

2. Variables

There are essentially four primary variables of interest in the sink process model: fddi_sink_accum_delay, fddi_sink_total_pkts, fddi_sink_total_bits, and fddi_sink_peak_delay. These exist as single integers or as floating point numbers, and are incremented or recomputed as packets are received by the station. The overall idea is to expand these into vector arrays, in which each element represents a running total for one priority setting, with the last element representing synchronous traffic totals. As mentioned earlier, this approach requires the "fddi_mac_fr" format in the Parameter Editor to be modified to include the priority as a field, since the original model needs only the packets' creation time and time of arrival in order to compute the overall throughput, mean delay, and end-to-end delay. However, as was discovered through trial and error, while the given variables can be changed to vector arrays, and the model can be modified to accommodate the new structure, and the code will compile (if the syntax is correct), any attempted simulation will abort with a segmentation violation error. This is because the "C" programs, which may be modified by the user, must interface with OPNET®'s kernel procedures, which are beyond the user's access. Ultimately, the given variables must be kept and new ones

created. Since the overall performance remains a useful statistic, the variables mentioned are left in place, while new ones are declared with the desired vector structure. These are `fddi_sink_accum_delay_a`, `fddi_sink_total_pkts_a`, `fddi_sink_total_bits_a`, and `fddi_sink_peak_delay_a`, which are declared and initialized in the Header Block.

The following declarations are added to the State Variables section:

```
Gshandle \thru_gshandle_a[10];
Gshandle \m_delay_gshandle_a[10];
Gshandle \ete_delay_gshandle_a[9];
```

Once compiled, the following appear in the ".C" file:

```
Gshandle sv_thru_gshandle_a[10];
Gshandle sv_m_delay_gshandle_a[10];
Gshandle sv_ete_delay_gshandle_a[9];
#define thru_gshandle_a  pr_state_ptr->sv_thru_gshandle_a
#define m_delay_gshandle_a
     pr_state_ptr->sv_m_delay_gshandle_a
#define ete_delay_gshandle_a pr_state_ptr-
     >sv_ete_delay_gshandle_a
```

(Appendix G, line 43)

With these declarations in place, the modifications to the rest of the code are straightforward, and generally follow the examples set by the original code. In addition, the variables `Offered_Load` and `Asynch_Offered_Load` are declared for use in generating scalar plots over a series of simulations. These are assigned values called from the Environment file by the state STATS.

3. Initialization State

The primary purpose of the Initialization State is to assign handles to the global statistics that are generated at the end of the simulation. The "KP" statement op_stat_global_reg (<*gstat_name*>) returns a handle used to reference a globally accessible statistic. This handle is needed to furnish new values (as they arrive with new packets) for the "KP" op_stat_global_write(), which appears in the DISCARD state. The field entry <*gstat_name*> is the text in the on-screen menu that appears in the Analysis Tool when the Create Single Vector Panel or "Create Multi Vector Panel" icons are invoked. The code added is very similar to what is already in place. For example:

```
thru_gshandle_a[0] = op_stat_global_reg ("pri 1 thruput
    (bps)");
```

(Appendix G, line 325)

creates a handle for the collection of data for priority 1 level throughput, and creates a field which will appear in the on-screen menu in the Analysis Tool. In another example:

```
m_delay_gshandle_a[9] = op_stat_global_reg ("async mean
    delay (sec)");
```

(Appendix G, line 364)

creates a handle for total asynchronous mean delay, and generates another field which will appear in the on-screen menu in the Analysis Tool. Each priority level has its corresponding handle assignment line. The actual statistics to accompany these handles are generated in the DISCARD state. As seen in the code itself (Appendix G), each element has a handle assignment line, for throughput, mean delay, and end- to-end delay. Figure 22 shows off the

resulting on-screen menu from the "Create Single Vector Panel" icon, reflecting the new data that may be plotted. Note that the code here reconciles "C" language vector numbering conventions with real-world priority level settings.

4. "DISCARD" State

The DISCARD state is where the received packet is "opened" and statistics generated from the packet contents. As mentioned before, "KP" op_stat_global_write (<*gstat_handle*>, <*value*>) is the statement that accumulates data. The <*value*> field may be a previously computed figure, or may be calculated within the "KP". The <*gstat_handle*> field is the same used in the INIT state. DISCARD uses the priority value found with the arriving packet as the index for the vector structure. For example:

```
op_stat_global_write (thru_gshandle_a[5],
fddi_sink_total_bits_a[5] / op_sim_time());
```

(Appendix G, line 128)

generates a current throughput figure for asynchronous priority level six (recall the necessary offset for vector element numbering convention). Also necessary is the recording of delay values for each priority, which is done with the following :

```
op_stat_global_write (ete_delay_gshandle_a[pri_set],
     delay);
```

(Appendix G, line 182)

This state also destroys the packet once its contents are recorded. This is necessary to prevent the simulation from filling the host computer's memory with dead packets.

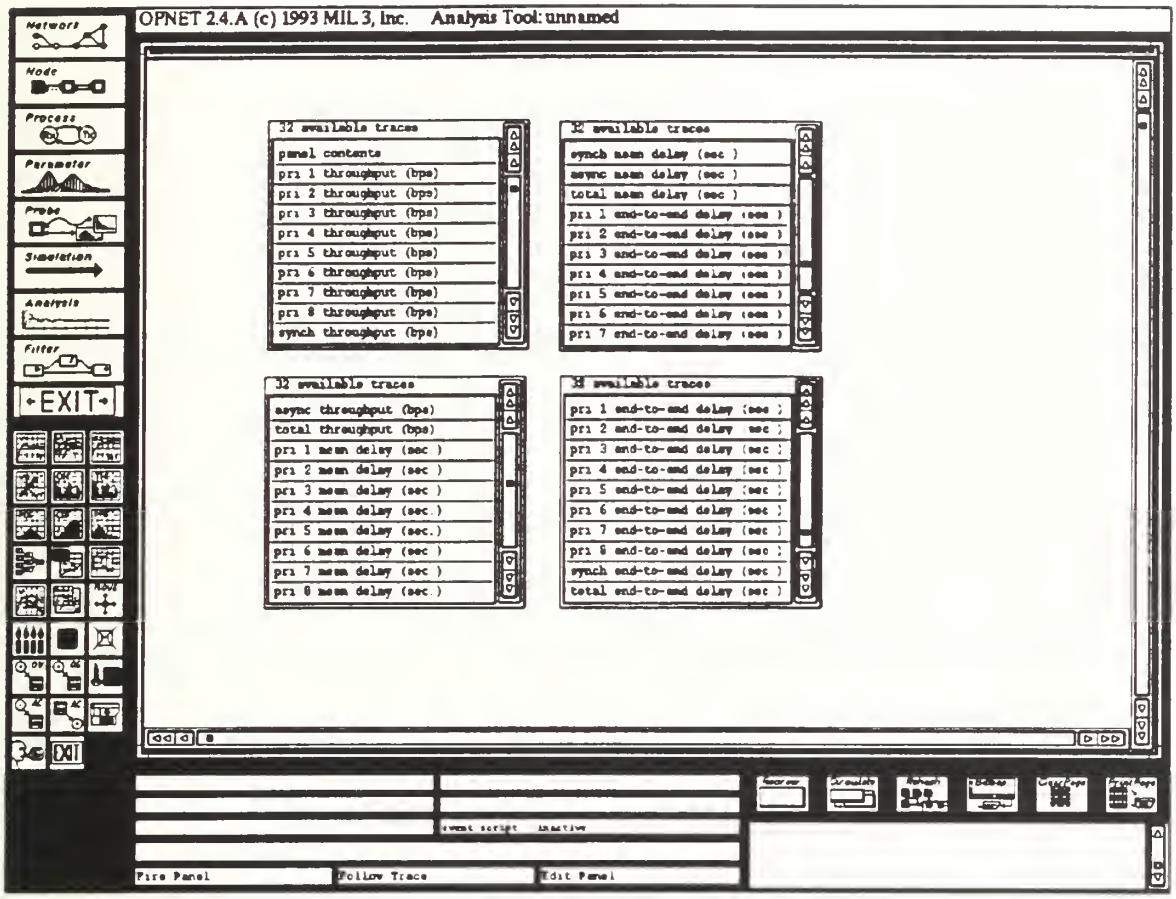

Figure 22. Newly Created Vector Traces Available

5. "STATS" State

The STATS state produces the steady-state scalar data that may plotted using the "Create Scalar Panel" icon. These are saved, rather than written over, so that the user may observe changes to output as the input is varied. For example, the throughput of synchronous traffic over several simulations as different TTRT values are used. From these, a plot of throughput vs. TTRT may be generated. The statement:

```
op_stat_scalar_write (<scstat_name>, <value>);
```

is similar to the "write" command described before, and writes a scalar steady-state statistic in this case. The field "scstat" appears in the on-screen menu called with the "Create Scalar Panel" icon. Examples of the use of this statement appear in Appendix G, lines 210-226.

Another on-screen menu line item is drawn from the Environment file. The values <*total_offered_load*> and <*asynch_offered_load*> are placed and assigned in the Environment file (see Appendix H), as described earlier. These correspond to the

variables `Asynch_Offered_Load` and `Offered_Load` declared in the header block. These are joined by the commands:

```
op_ima_sim_attr_get (OPC_IMA_DOUBLE,
     "total_offered_load", &Offered_Load);
op_ima_sim_attr_get (OPC_IMA_DOUBLE,
     "asynch_offered_load", &Asynch_Offered_Load);
```

(Appendix G, line 299)

and added to the on_screen menu with the commands:

```
op_stat_scalar_write ("Total Offered Load (Mbps)",
     Offered_Load);
op_stat_scalar_write ("Asynchronous Offered Load (Mbps)",
     Asynch_Offered_Load);
```

(Appendix G, line 305)

This code in Appendix G contains a warning to the user that the offered load settings are not automatically updated in any way. If the user desires to plot throughout or delay as a function of offered load over a series of simulations, then the user must remember to keep the offered load assignments current in the Environment file for each simulation.

D. BRIDGE LINK

The alteration described here represents a simple first step toward a network interface. Instead of destroying frames after they are received, one station on the LAN is modified to hold its packets in subqueues. Further development will bring about a protocol for removing these buffered packets from the original LAN and transferring them to another.

1. Station Model Modifications

The received frames are stored in subqueues according to their priority in a manner analogous to that already described for the `mac` node. This requires that the `sink` node, `llc_sink`, be changed from its original processor form into a queue node. To affect this modification to the station model, the following steps are followed:

1. The `sink` node is selected by placing the cursor over it and clicking the left mouse button.
2. The node is removed by invoking the "`Cut`" icon. The "`Packet Stream`" between the `sink` and the `mac` also disappears.
3. The "`Create Queue`" icon is selected, and the resulting box is dragged to the location just vacated. Clicking the left mouse button places the new node. The station now has a queue node rather than a processor node.
4. The on-screen attributes menus is called with the right mouse button, and the fields are all set to the same values that were in effect before, including the node name.
5. The "`process model`" field will be set to the newly modified sink process model. If the process has not yet been modified, then the original assignment may be used, then changed when the process has been modified and saved under a new filename.
6. The "`subqueues`" field is set to *<9>*, accommodating eight levels of asynchronous traffic and also synchronous traffic.

7. The packet stream line between the `sink` and the `mac` nodes must be replaced.

8. The new station model is saved under a new filename, for example, *<"fddi_sink_link">*.

Figure 23 shows the resulting station model, with the appropriate on-screen attributes menu.

2. Process Model Modifications

The changes to the process model code are few, and are included in Appendix G as inactive code ("commented out"). The code that destroys received packets:

```
op_pk_destroy (pkptr);
```

(Appendix G, line 98)

is replaced by code that enqueues the packets according to their priority settings:

```
op_subq_pk_insert (pri_set, pkptr, OPC_QPOS_TAIL);
```

(Appendix G, line 105)

The user should realize that for the moment, no more code exists for the disposition of these enqueued packets; a long simulation simply accumulates packets and fills computer memory. The subqueues are infinite by default, but may be limited (using another on-screen menu) to demonstrate overflows. In that case, there is no code for the disposition of packets that are lost through buffer overflow, and these will simply accumulate in the host computer's memory as well. In sum, the user must be aware of the memory demands of OPNET® simulations.

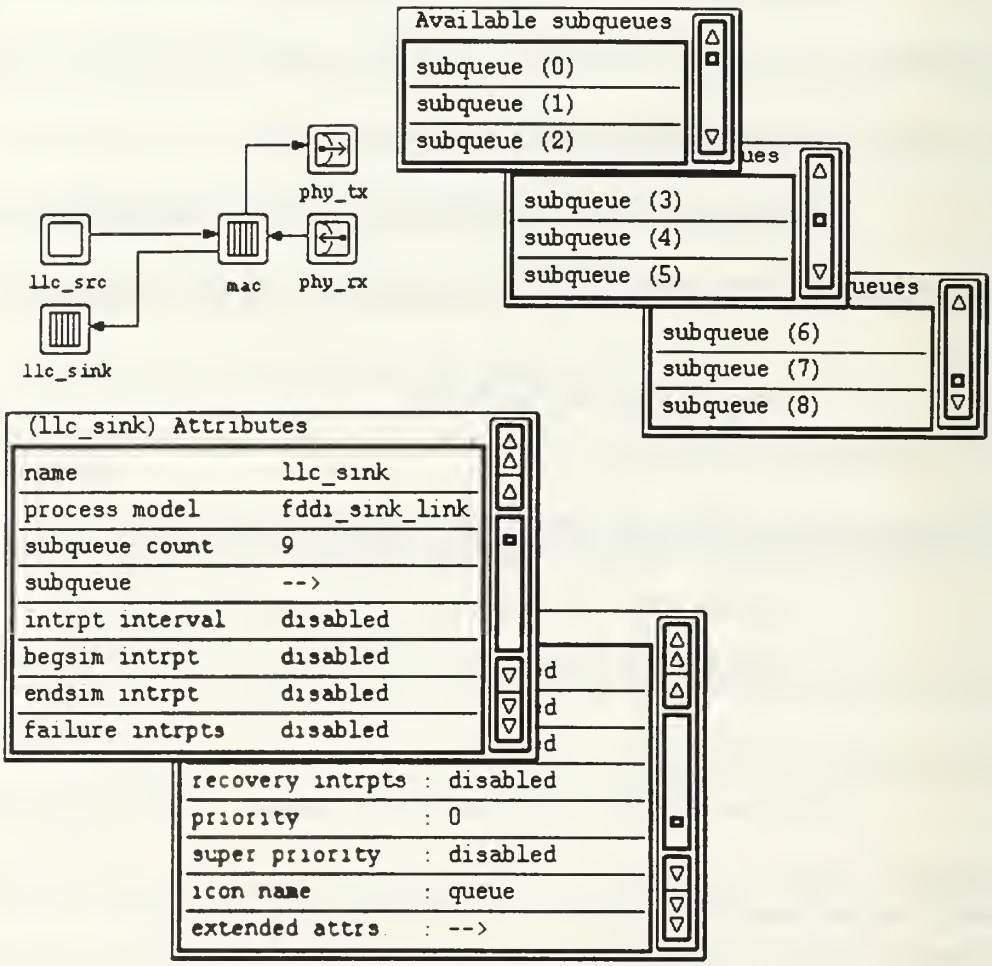

Figure 23. Modified Station Model, "fddi_sink_link"

E. MULTICAST

1. Overview

Multicast is the addressing of a packet to more than one station. Broadcast is a special case with the transmission of the packet to all stations. In the original model, a station desiring to send the same message to all stations would repeat the packet transmission for each destination station. Actually, this last is an abstraction; OPNET® simply generates packets from each station addressed to randomly generated destinations, with no indication that any particular transmission represents a copy of any previous transmission. However, the fact remains that each packet is addressed to only one station. In the actual FDDI protocol, each packet is passed from station to station until it reaches its destination, but then continues past its destination until it is finally removed from the ring by its originating station. The OPNET® simulation economizes on the number of simulation events (and therefore the simulation time) by having the destination stations remove the packets they receive. This occurs in the `mac` state `FR_REPEAT`, which also contains comments suggesting that the user may wish to overrule this economizing feature in the event that group addressing is desired. Figure 24 is the State Transition Diagram for the `mac` process, repeated from Figure 6 for the reader's reference. The state `FR_STRIP` includes the code by which an originating station removes packets that have completed a circuit of the LAN. It is not used in the original code, nor will it be used in the modifications described here.

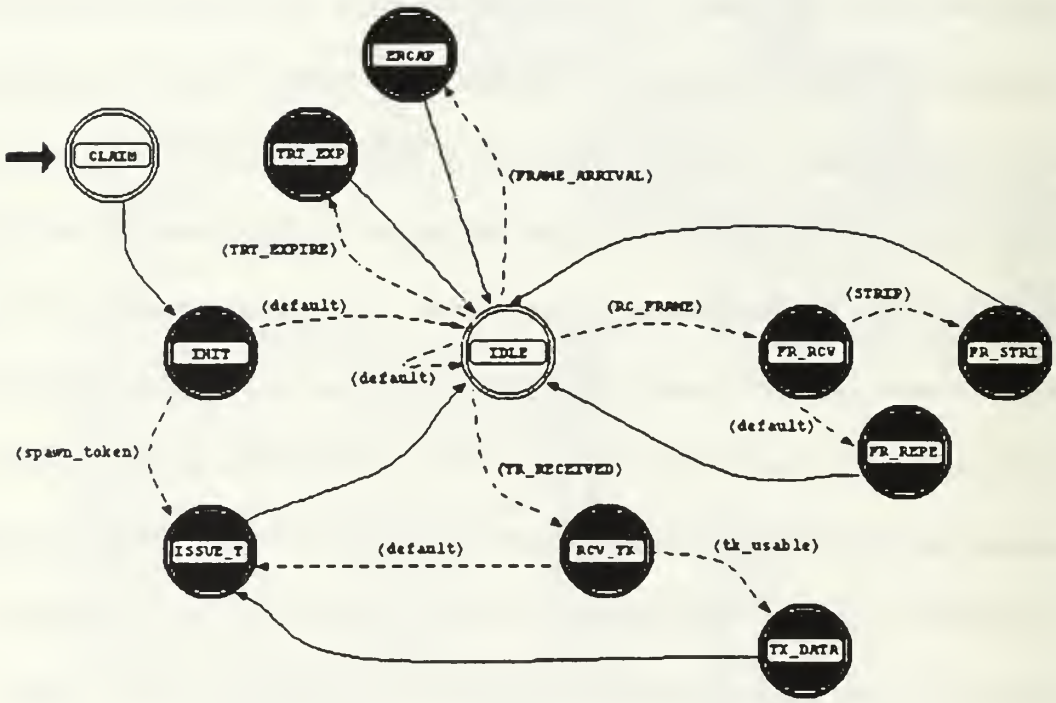

Figure 24. MAC Process Model State Transition Diagram

The basic idea is as follows: rather than carry a destination address, each packet carries an array with a number of elements equal to the number of stations on the ring. These elements are simply ones and zeros, with a one indicating by its location that the packet is addressed to a station corresponding to that location. For example, in a five station LAN, the address field `[0 1 0 1 1]` would indicate the packet is addressed to stations one, three, and four (as in the case of indexing vector array elements, the stations on an N-station LAN in OPNET® are numbered from zero to N-1). As the packet is passed around the LAN, each station inspects this array, passing the packet on if the station is not designated a destination address. Destination addresses keep a copy of the packet's information, set their place in the address field to zero, then pass the packet on to the next station. The last destination address will destroy the packet after verification that only zeros remain in the destination address array, thereby preserving some of the economy gained in minimizing the number of events in the simulation. While the destination address array would need to be transported with each packet, OPNET®'s Kernel Procedures can only accommodate a pointer to the array. The following sections describe the changes necessary to effect multiple addressing. The reader is again referred to Appendices E, F and G, containing the ".C" files for the MAC, source, and sink process models, respectively. However, the implementation of multicasting involves no changes to the sink process model.

2. Environment File

Each station is assigned a maximum and a minimum number of possible destination addresses to use in addressing each packet. The following are added to the Environment file:

```
"*.*.llc_src.min num addees":   <user assigned integer>
"*.*.llc_src.max num addees":   <user assigned integer>
```

These will be called by the source process model's INIT state. Appendix H contains an example Environment file including these new attributes. The minimum number must be at least one, and the maximum should be no greater than N-1, where N is the number of stations (the logic written in the source model's code does not allow stations to address packets to themselves). If it is desired that some station generate no traffic, then the "arrival rate" field should be set to zero. Setting "min num addees" and "max num addees" to zero will only result in packets transmitted with no destination addresses assigned. Setting "max num addees" to a number greater than N-1 will cause an endless loop in the code that generates destination address assignments in the source process (the logic in the loop is, "assign *x* different destinations, but do not repeat any.").

3. Station Model

The additions to the environment file must be added to the station model in the Node Editor, in a manner analogous to the method described in the discussion of prioritization (III.B.4.a. Station Model Changes). The steps required to add another attribute to the on-screen menu are not recounted here, but the desired final result is shown in

Figure 25. It is a good practice to save the changed station model under a new name, for example <"fddi_sta_mult">, until the user is certain that the modifications do more good than harm to the original model.

4. Source Process Model

A real packet on a real LAN would necessarily be self- contained, carrying with it all its destination addresses. However, the functions used to assign the packet address field in OPNET® will not support a vector structure, and so pointers to memory locations must be used, with these memory locations containing the destination addresses. The modifications to the source process to effect multiple addressing are summarized in the following sections. As is the case with the station node model, each of the process models should be saved under new names, for example `fddi_gen_mult`, as a matter of practice.

a. Variables

To begin, a global variable, `NUM_STATIONS`, is defined in the Header Block (it is also defined in the `mac` process model), to be used as a loop counter. This variable must be kept updated to accurately reflect the correct number of stations on the FDDI LAN. This is easily forgotten when different LANs are created, using the same station model with different numbers of stations.

As was mentioned, the destination address field, originally a single integer value, must now be made a pointer to an array of integers. Although "C" programming language treats the name of a vector array as the array's pointer, the established Kernel Procedures do not support the simple change of syntax. In short, a new variable, `*da_ptr` (destination address pointer), is declared as an integer pointer in the Temporary Variables

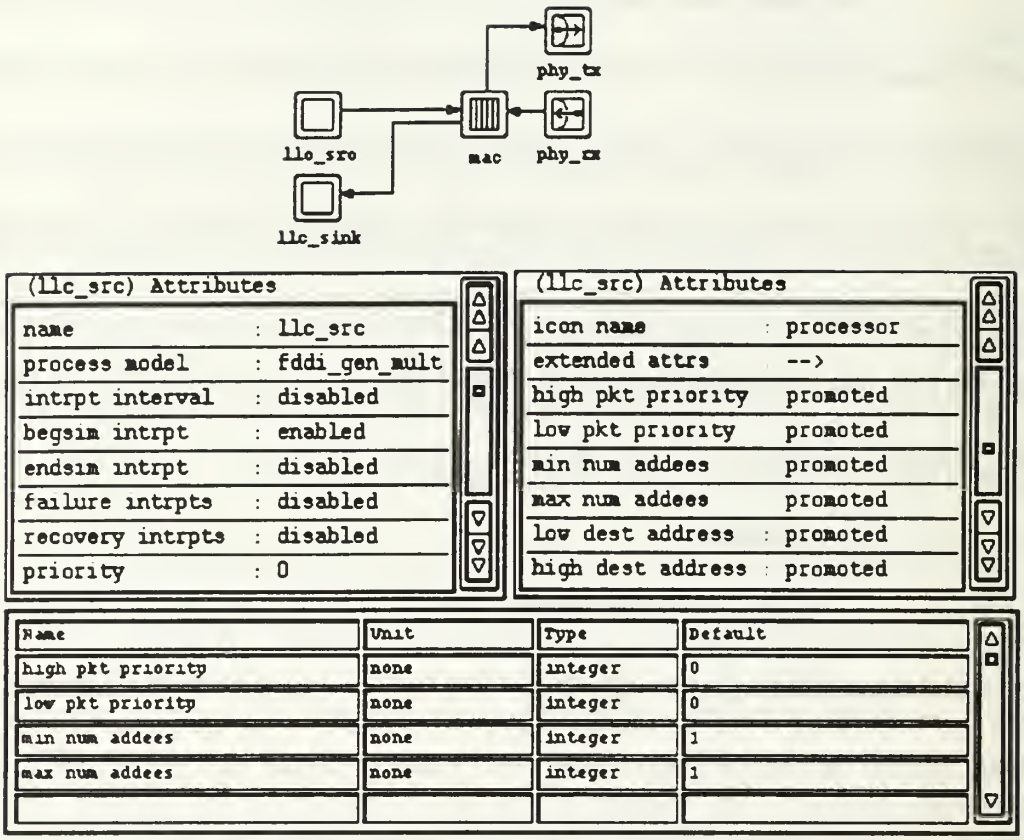

Figure 25. Multicast-Capable FDDI Station with Attributes Menus

(TV) editor. The original variable, `dest_addr`, is removed from the TV section, and declared with the State Variables (SV) as an array of integers, with dimensions `[NUM_STATIONS + 1]`. In addition, variables to accompany the Environment file attributes are declared with the State Variables: `min_num_addees` and `max_num_addees`. The resemblance between these variables and the attributes they accompany is meaningless in respect to "C" language syntax, but is of course a useful memory aid to the user. The pointer `num_addees_dist_ptr` represents the value used in determining stochastic values. Also, the integer `num_addees` is declared in the State Variables, to represent the number of stations to which a given packet will be assigned. It is used as a loop counter, and will be different for each packet.

b. Initialization State

In the source model's Initialization state, the range bounding the number of destination addresses is determined with a call to the Environment file:

```
op_ima_obj_attr_get(my_id, "min num addees",
     &min_num_addees);
op_ima_obj_attr_get(my_id, "max num addees",
     &max_num_addees);
```

(Appendix F, line 86)

These result in the assignment of the values from the Environment file to the addresses of the corresponding variables. A distribution is established with the following:

```
num_addees_dist_ptr = op_dist_load ("uniform_int",
     min_num_addees, max_num_addees);
```

(Appendix F, line 114)

This value is used to generate streams of stochastic values, and is used in the ARRIVAL state.

c. "ARRIVAL" State

In the ARRIVAL state, the actual number of stations to receive the new packet is determined, and then a loop is used to choose these stations one at a time. Each loop iteration is very similar to the original procedure that was in place when only one station was assigned to each packet. The loop contains a provision to prevent the repeated assignment of the same station, and also to prevent the assignment of the originating station as a destination address. The following statement determines the number of destination addresses for a given packet:

```
num_addees = op_dist_outcome (num_addees_dist_ptr);
```

(Appendix F, line 172)

The following loop is used to find and assign the chosen stations:

```
for (i = num_addees; i > 0; i--)
{
     gen_packet:
     nix = op_dist_outcome (dest_dist_ptr);
     if (dest_addr[nix] == 1 || nix == station_addr)
     {
          goto gen_packet;
     }
     dest_addr[nix] = 1;
}
```

(Appendix F, line 174)

This loop continues to iterate until the specified number of stations, without repetition, is assigned.

Recall that the destination address array is declared with one element more than the number of stations in the LAN (`dest_addr [NUM_STATIONS + 1]`). Here, all the elements in the destination array must be shifted one space to the right, and a simple loop is used to set `dest_addr[i+1]` equal to `dest_addr[i]`, for `i` iterations. This step is necessary because the first array element will be overwritten with the array's memory address in the course of the packet's travels, as it is transmitted from one station, received by the next, and its destination address field is opened, inspected, then closed by each station in turn. This behavior is verified by use of the debug tool set to "`fulltrace`," accompanied by strategically placed `printf` statements. The author does not pretend to know why this happens. The array element shift is a deft enough way to sidestep the problem. However, all references in other states to the destination address array must be offset to accommodate this shift.

5. MAC Process Model

The MAC process model receives each packet, inspects it, the decides whether or not the packet is addressed to the station. In the original model, the packet is removed from the ring by the destination station, and relayed by other stations. With multicasting of packets, a third case arises, in which a station receives a packet addressed to it, but must also pass the packet on to other destination stations. A number of print commands are placed in the code, but left inactive. They are of much use in the verification of the model's operation.

a. Variables

The same variable `NUM_STATIONS` used in the source model is also defined in the Header Block of the MAC. The user must remember to keep this value updated in both places when the same station model is used in a different size LAN. The destination address is changed from an integer into an open-ended array of integers, `dest_addr[]`, in the Temporary Variables editor. An integer pointer, `*da_ptr`, is declared as well.

b. Encapsulation State

The Encapsulation state receives frames from the source process, and place them inside the format `fddi_mac_fr` for transmission on the LAN. An intermediate step is to inspect the frame received from the source for its destination address, which must be written into the encapsulating frame's destination address field as well. The original statement:

```
op_ici_attr_get (ici_ptr, "dest_addr", &dest_addr);
```

is unworkable with `dest_addr` in array form, which is why the pointer `*da_ptr` is declared. Instead, the following is used:

```
op_ici_attr_get (ici_ptr, "dest_addr", &da_ptr);
```

(Appendix E, line 806)

followed by a loop assigning each element in the array a value from the corresponding element in the array found at address `da_ptr`. Vol. 5.0, *Simulation Kernel Manual*, discusses this command statement. This loop uses as a counter the value `NUM_STATIONS+1`, for the reason mentioned earlier: use of a `printf` statement would reveal that the first element in the array has been written over and replaced with a number representing the memory address of the array. Fortunately, the entire original array of zeros and ones has been shifted, so the first element is intact. Correspondingly, all references to the array from within the MAC process must be made with respect to this shift.

c. Frame Repeat State

The Frame Repeat state inspects each received packet and acts on one of three cases: the packet is addressed only to this station, or the packet is not addressed to this station at all, or the packet is addressed to this station and to other stations as well. The first two cases are already present in the original model, and require some simple modifications. The third case represents a significant change, requiring the addition of an entire block of code to the state, in which the packet information is copied first, then passed on to the next station.

The first statement in the FR_REPEAT state opens the arriving packet's address field for inspection:

```
op_pk_nfd_get (pkptr, "dest_addr", &da_ptr);
```

(Appendix E, line 604)

Here, `&da_ptr` has been substituted for the original `&dest_addr`. As is the case in the ENCAP state, the arriving pointer is used to initialize the locally held destination address array:

```
for (i= 0; i < NUM_STATIONS+1; i++)
     dest_addr[i] = da_ptr[i];
```

(Appendix E, line 610)

This destination address array is then passed through a loop to determine if it has more than one destination address, and to see if the element corresponding to this station is set to one, indicating the packet is addressed to this station. If the packet proves to be addressed to this station only, then the actions of the original code are carried out: relevant fields are copied to an ICI packet format for transmission to the sink process, then the packet is destroyed. If inspection of the destination address array shows the packet is not addressed to this station at all, the original code is again sufficient to place the packet back on the ring.

The third case represents a new situation. When the packet is addressed to this station and to others as well, the information must be saved here, and be transmitted onto the ring again. These actions are carried out with commands borrowed from the first two cases, and include some new considerations as well. The function "`op_pk_nfd_get()`"

is used to retrieve data from specified fields in the packet. That is, it is a decapsulation function. When the data happens to be in a structure format, the function has the effect of destroying the information. This is an important point because the information must be preserved for retransmission. Therefore the function:

```
op_pk_nfd_get (pkptr, "info", &data_pkptr);
```

(Appendix E, line 719)

must be followed at some point with the statement:

```
info_ptr = op_pk_copy (data_pkptr);
```

(Appendix E, line 727)

in which `info_ptr` has been declared in the Temporary Variables to have the same type as the structure in the "`info`" field. When this information is summoned for re-encapsulation with the function:

```
op_pk_nfd_set (pkptr, "info", info_ptr);
```

(Appendix E, line 747)

the field information has been preserved. The other fields, "`src_addr`", "`dest_addr`", and "`pri`", do not require this procedure, since they are not lost with decapsulation. Two packets result: one is an ICI frame carrying the received information to the sink process, and the other is a re-encapsulated packet sent to the LAN. The station's last action before

re-encapsulating the destination address array is to zero its corresponding address element. The last station to receive the packet will destroy it, upon determining that only zeros remain in the destination address array (Appendix E, lines 644 and 692).

6. Limitation

The primary limitation of the model with multicasting active is in the generation of throughput and delay statistics. As the code currently stands, each packet is counted by every station that receives it, leading to multiplication of throughput data. On the other hand, some use may possibly be made of this characteristic by comparing the tallied throughput against the actual offered load, as a measure of the effectiveness of the multicasting scheme. For example, a measured throughput of 100 Mbps versus a known offered load of 50 Mbps could indicate that multicasting effectively generates 50 Mbps without physically taxing the bandwidth capability of the FDDI LAN, since no additional packets are generated.

IV. MODEL TESTING

This chapter provides several test results intended to verify the enhanced capabilities added to the basic FDDI station model in OPNET®. In addition to improving the performance and display features, the modifications must preserve the basic behavior to be useful. The tests presented here include an illustration of the model's treatment of synchronous and priority-based asynchronous traffic, a simple demonstration that the modified sink process does indeed store the traffic it receives, a check against theoretical performance equations, and a demonstration on the monitoring of multicast.

Not all of the modifications described in this thesis are incorporated simultaneously in all station models. The State Transition Diagram correction to the original `fddi sink` process model is of course installed in all models that use the `llc_sink` node. Likewise, the patch for OPBUG is installed for all MAC model versions. The generation of randomly differing priority assignments for all packets is closely interrelated to the generation of output statistics segregated by class and priority; both involve extensive modifications to the original process models for the source, sink and MAC, which are stored as `fddi_gen`, `fddi_mac`, and `fddi_sink`. (The actual file names found in the UNIX subdirectory have the suffix "`.pr.m`," and when these are compiled, corresponding files with the suffix "`.pr.c`" are created.) The original models of those names are retired under the ending "`_orig.pr.m`." Likewise, the original node-level station is stored under the name `fddi_sta_orig`, while `fddi_station` is the upgraded

version. The multicasting capability is stored in a separate station model, `fddi_sta_mult`, which includes the processes `fddi_mac_mult`, `fddi_gen_mult`, and `fddi_sink_mult`, and contains all of the other modifications as well. The multicast capability was the last feature installed for this thesis, and has not been completely developed.

A. SYNCHRONOUS THROUGHPUT

1. Overview

This test was motivated by an earlier study, in which a software simulation tool was used to demonstrate the effect of increasing offered load on the asynchronous and synchronous throughput of an FDDI LAN. Using the Block Oriented Network Simulator (BONeS®), Tari et. al. were able to show that synchronous throughput does not degrade appreciably, even when the offered load is well in excess of 100 Mbps. They were able further to show the decay in throughput suffered in the asynchronous priority levels as offered load is increased. Figure 26 is taken from this study. In addition to providing an idea of expected performance, this figure also highlights the fact that in its original form, OPNET®'s FDA model provides no method to display any throughput data other than the overall performance parameters, total throughput, total mean delay, and total end-to-end delay.

2. Setup

The set up was intended to imitate the older experiment as closely as possible. Ten simulations were conducted using the same 50-station LAN. Ten stations (`f0, f5, f10, f15, f20, f25, f30, f35, f40, f45`) were designated "`constant`" generators.

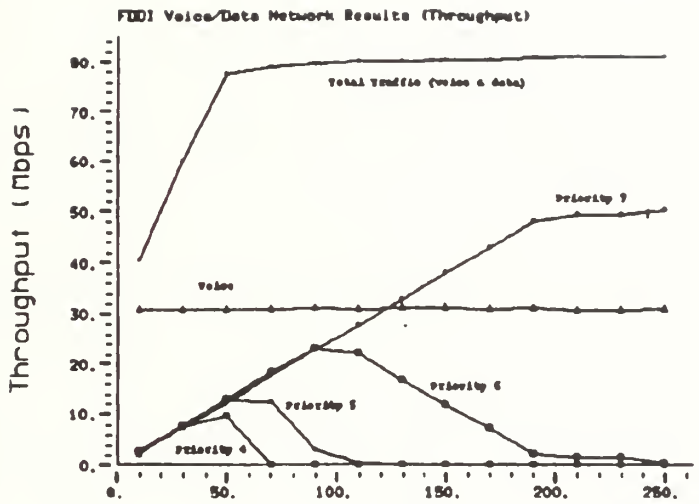

Figure 26. Throughput Measurements Using BONeS® (Tari, et. al., 1988, p. 58)

That is, they were modified inside the source process "C" code to generate packets at a constant rate, rather than at an exponentially distributed approximation of the rate specified in the Environment file. These ten stations transmitted synchronous traffic only, while the other 40 stations were allocated no synchronous bandwidth, sending asynchronous traffic only. The synchronous stations generated 512-bit packets at a constant arrival rate of 6000 packets per second. The asynchronous stations transmitted 1000-bit packets, at a rate stepped so as to increase the total offered load by 10 Mbps in each simulation. Therefore, total offered load ranged from 40.72 Mbps to 140.72 Mbps over ten tests, with steady state data recorded in a data file of scalar data. TTRT was set to 10.7 ms, in order to support the delay of 50 stations on a 50 km ring while maintaining $\sum SA_i \geq 10$ ms. This differed from the older study, for which a TTRT value of 8.0 ms was used, despite the authors' stated intention to model voice network traffic. As noted by Powers (1993, p.340), synchronous voice transmission requires that the token visit each voice transmitter every 20 milliseconds, resulting in $\sum SA_i = 10$ ms. This in turn requires TTRT to exceed 10 ms by an amount sufficient to account for physical delays inherent to the ring fiber and the stations. Asynchronous priority threshold levels. `T_Pri[i]`, in the INIT state of the MAC process code were set in increments of 0.125 TTRT, so that Priority 1 traffic transmission was cut off when THT incremented to 1.3375 ms, and Priority 8 traffic could be sent for the entire THT period. Recall that this incrementing THT is a function of OPNET, reversed from the decrementing timer described in most literature.

3. Results

Figure 27 shows a plot of total throughput over a series of ten simulations in OPNET, demonstrating a roughly linear rise until the offered load begins to exceed 90 Mbps, which is in qualitative agreement with the older experiment, and agrees with results published by Dykeman and Bux (1988, pp. 1003-1007). This plot also serves to show off one of the model's improvements in data display, allowing `"Offered Load"` to be an abscissa for scalar plots.

Figure 28 illustrates the fact that synchronous bandwidth allocation is not affected by the asynchronous offered load; therefore, synchronous throughput remains nearly constant as the offered load is increased.

Figures 29-31 illustrate the effect of increasing offered load on the throughputs of asynchronous traffic at priority levels two, three, and four, respectively. No priority settings above four suffered any degradation within the range of offered load observed in this test. The decay of Priority 3 traffic at approximately 100 Mbps of offered load suggests that (3/8)TTRT may be a guiding point for setting priority thresholds that will take effect before the LAN's capacity is reached.

B. PRELIMINARY LINKING MODEL

1. Overview

The modifications made to the sink process as a step toward a bridging node are the simplest of all described in this thesis, and the testing of the finished model is also simple.

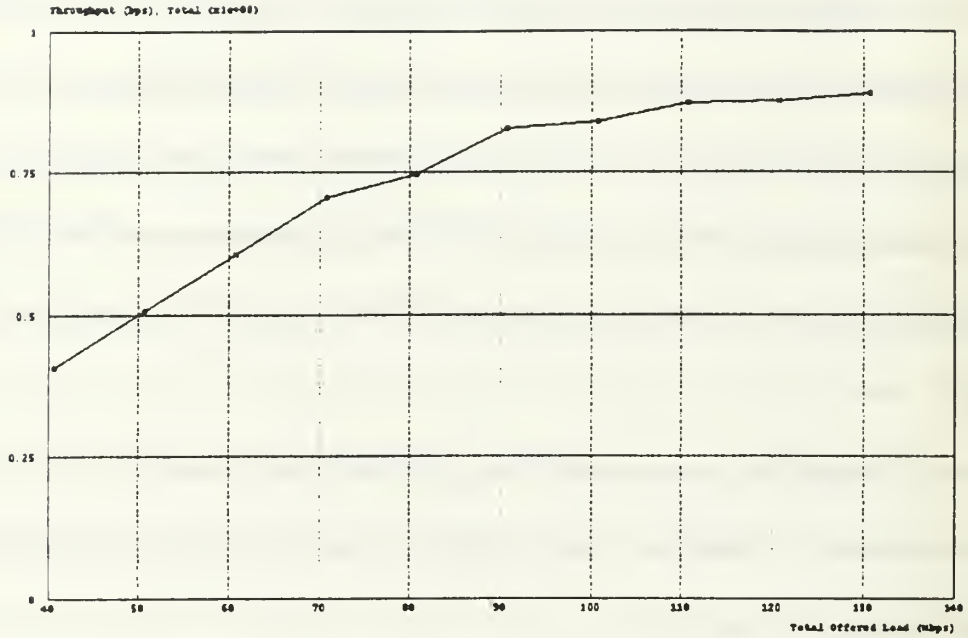

Figure 27. Total Throughput vs. Total Offered Load

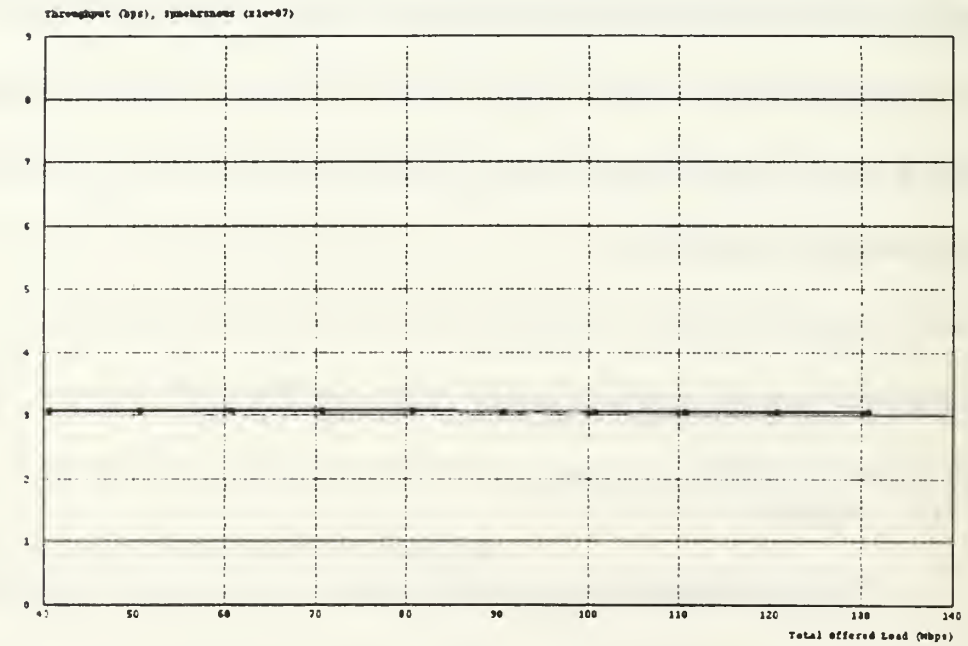

Figure 28. Synchronous Throughput vs. Total Offered Load

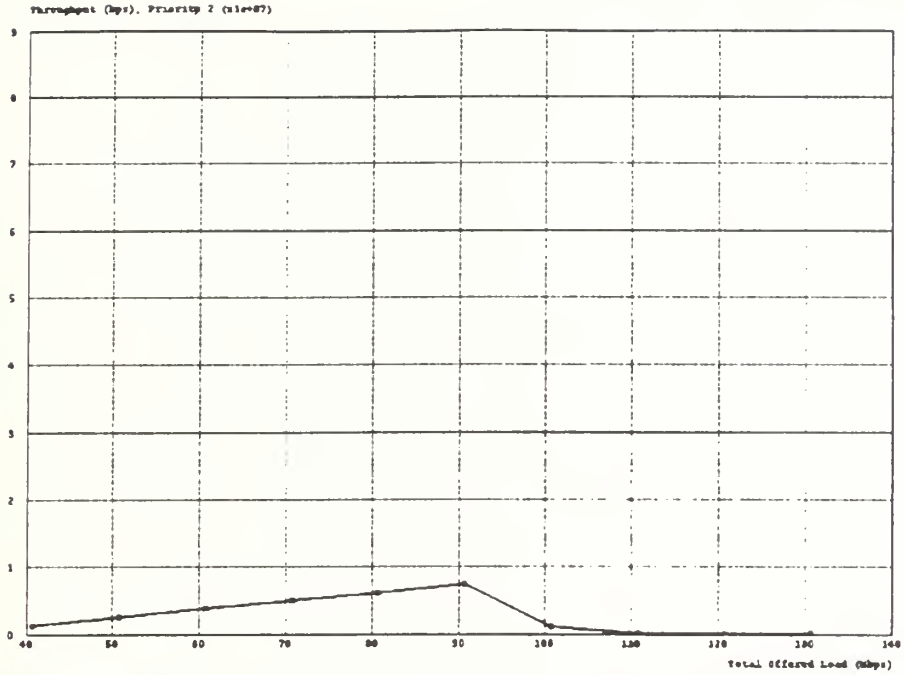

Figure 29. Priority Two Throughput vs. Total Offered Load

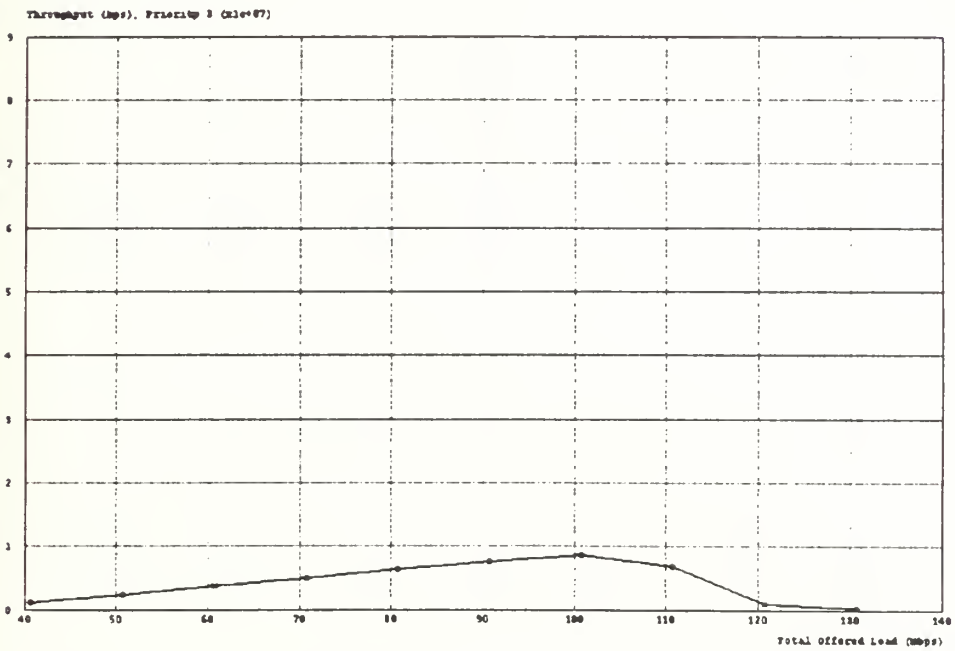

Figure 30. Priority Three Throughput vs. Total Offered Load

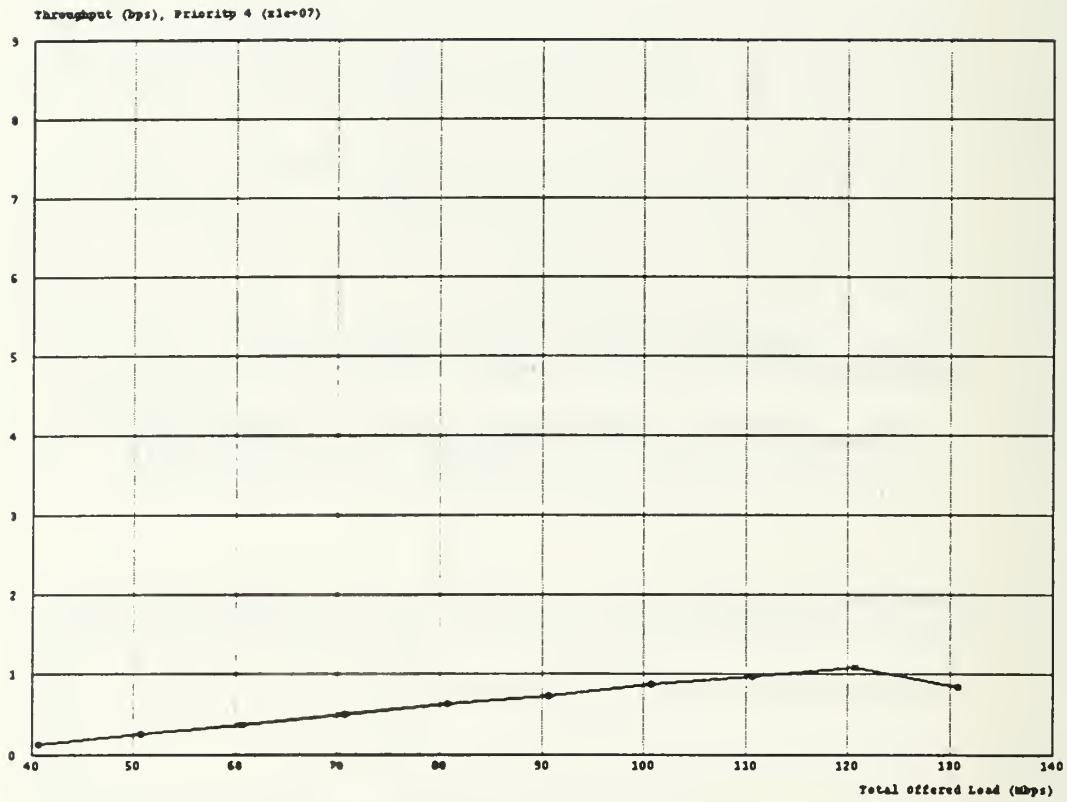

Figure 31. Priority Four Throughput vs. Total Offered Load

It is necessary only to show that all the traffic transmitted to the station is received and held. To demonstrate this action, a ten-station LAN was created, with station `f9` designated the link node. The Environment file assignments were made so that all stations directed all traffic to station `f9`, which itself was not transmitting. A modest offered load was used to prevent an overload of packets in the simulation host computer's memory. All transmitting stations were assigned the same full range of asynchronous priority settings, and were assigned equal portions of synchronous bandwidth. Each of the buffers at station `f9` should have then been seen to receive packets pre-sorted for eventual transmission. The Probe Editor was used to monitor "`pksize`," the number of packets in each subqueue of the modified sink process model.

2. Setup

The following calculations and Environment file settings in Table 1 were chosen as reasonable:

Table 1. ENVIRONMENTAL FILE SETTINGS.

packet size:	2000 bits
arrival rate:	10 pkt./sec.
offered load:	9 • 10 • 2000 = 180,000 bps.
TTRT (T_Req):	0.004 sec.
async mix:	0.9
prop delay:	5.085×10^{-06} sec./km. × 1 km.) = $5.085 \times^{-06}$ sec.
station latency:	60.0×10^{-08} sec.
D_Max:	(5.085 μsec. • 10) + (60.0×10^{-08} • 10) = 0.05685 ms.
F_Max:	0.360 ms.
Token Time:	0.00088 ms.
synchronous BW:	4-(0.05685 + 0.360 + 0.00088) = 4 - 0.41773 = 3.58227 ms.

Dividing bandwidth evenly among 9 stations gives the following:

$$\frac{3.58227\ ms.}{9\ stations} = 0.39803\ ms./station. \quad (2).$$

This result is compared with TTRT to determine the "sync bandwidth" attribute:

$$\frac{0.39803\ ms.}{4.0\ ms.} = 0.0995075\ ms. \quad (3).$$

which is a unitless fraction of TTRT.

3. RESULTS

The given parameters were applied, and the receiving buffers were inspected at the end of one second of simulation time. Figure 32 is the plot obtained illustrating the accumulation of packets in all nine subqueues. (Although all the plots may be placed in one panel, the data are divided into two panels for readability of the hardcopy). The plots show the following after one second:

Table 2. SUBQUEUE ACCUMULATION

```
subqueue(0): 13 packets
subqueue(1): 11 packets
subqueue(2): 10 packets
subqueue(3):  7 packets
subqueue(4):  7 packets
subqueue(5): 12 packets
subqueue(6): 11 packets
subqueue(7): 10 packets
subqueue(8):  9 packets
-----------------------
   total: 90 packets
```

This agrees with the given packet generation rate of 90 packets/sec., and verifies that the preliminary linking model holds all received packets on station.

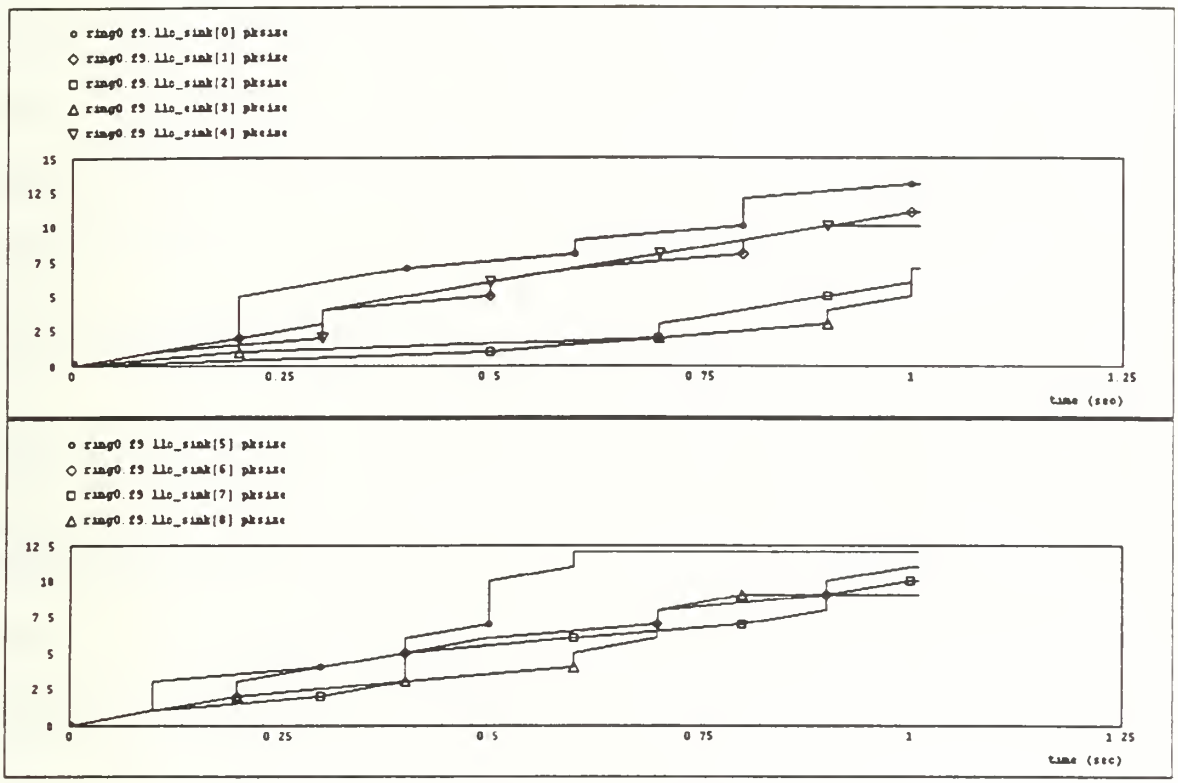

Figure 32. Packet Accumulation at Linking Model "fddi_sink_link"

C. SYNCHRONOUS TIMING

1. Overview

In order to preserve the real-time nature of synchronous traffic, the bandwidth allocated to a given station may not be exceeded. Although transmission time (bandwidth) may remain to a station, the protocol determines a priori whether the next packet transmission would cause the allotted bandwidth to be exceeded. If the bandwidth limit would be exceeded, then the station will not transmit. The code in the `TX_DATA` state of the MAC process ensures OPNET®'s FDDI LAN models adhere to this behavior, as shown in the following simple experiment.

2. Setup

A new LAN was created, using 13 stations on 91 km. of fiber, resulting in seven kilometers of fiber between stations. The odd figures were chosen in order to avoid symmetries in the arithmetic involved, thereby enhancing the instructional value of the test. All 13 stations were assigned an equal portion of the total available bandwidth for synchronous traffic, and all stations transmitted only synchronous traffic. Physical attributes were identical for each station. The following calculations apply:

Table 3. 13-STATION LAN ENVIRONMENT SETTINGS

prop_delay:	7 km. • 5.085×10^{-6} sec./km. = 0.0355950 ms.
station latency:	60.0×10^{-8} sec./station
F_Max :	0.360 ms.
D_Max :	(13 stations • 60.0×10^{-8} sec./station) + (13 links • 0.035595 ms./link) = 0.4627428 ms.
Token_Time :	0.88 μsec.
	D_Max + F_Max + Token_Time = 0.8236228 ms.
TTRT:	4.0 ms.
	$TTRT \geq \sum SA_i + 0.8236228 \Rightarrow \sum SA_i \leq 3.1763772$ ms.
Divide the synchronous allotment evenly among 13 stations:	
	3.1763772 ms./ 13 stations = 0.2443367 ms./station.

Given the standard transmission rate of 100 Mbps, each station may transmit 24,433 bits with its synchronous allotment. This bandwidth is converted to a fraction for the Environment file attribute "sync bandwidth":

$$\frac{0.2443367\ ms.}{4.0\ ms.} = 0.06108418 \qquad (4)$$

For this test, all stations were constant transmitters. That is, the code for the source process was adjusted so that packet transmission rates and packet lengths were assigned invariant values, rather than stochastic approximations of the attributes assigned. A TTRT of 4 ms indicates 250 token passes per second, which was chosen as the packet arrival rate for all stations. Packet size was 24,000 bits, resulting in a total offered load of 78.0 Mbps. A larger packet size, for example 25,000 bits, should be too large to transmit, resulting in no throughput at all.

3. Results

Figure 33 shows the resulting total throughput derived from the given setup. Invoking the "`convert to text`" attribute of the on-screen menu indicates a value of 77.7 Mbps after one second, with a slow rise still in progress. This is in close agreement with the offered load.

Figure 34 is an empty panel, accompanied by a text screen indicating that no throughput results when the packet size exceeds that allowed by the synchronous bandwidth assignment. In this case, packet size was increased to 25,000 bits. Arrival rate was also reduced to 100 packets/sec., resulting in a total offered load of 32.5 Mbps. Figure 35 illustrates the accompanying accumulation of packets in station `f3`'s subqueue(8), which is reserved for synchronous traffic. The perfectly linear shape is a result of the constant arrival rate assignment.

D. ASYNCHRONOUS EFFICIENCY

1. Overview

A network's throughput efficiency is calculated from the following equations, which assume that only asynchronous traffic is being transmitted:

$$E = \frac{N(T-D)}{(TN+D)} \quad \textbf{(5)}.$$

where:

N = number of stations,

T = TTRT and,

D = ring latency (total time required for a token to circulate the ring in the absence of data traffic).

Ring latency D is in turn defined as

$$D = L(\frac{n}{c}) + N \cdot T_s \quad \textbf{(6)}.$$

where:

L = length of ring in kilometers

$n/_c = 5.085 \times 10^{-6}$ sec./m.

T_s = token processing time (0.88 μs.)

(Powers, 1993, pp. 336-337; uses $T_s = 1.0$ μs. as a typical value)

The same 13 station, 91 km LAN described previously was used here, with the Environment file adjusted so that no synchronous bandwidth was assigned, and only asynchronous traffic was generated. Offered load was 78.0 Mbps. The given equations indicate the expected efficiency is 87.44%.

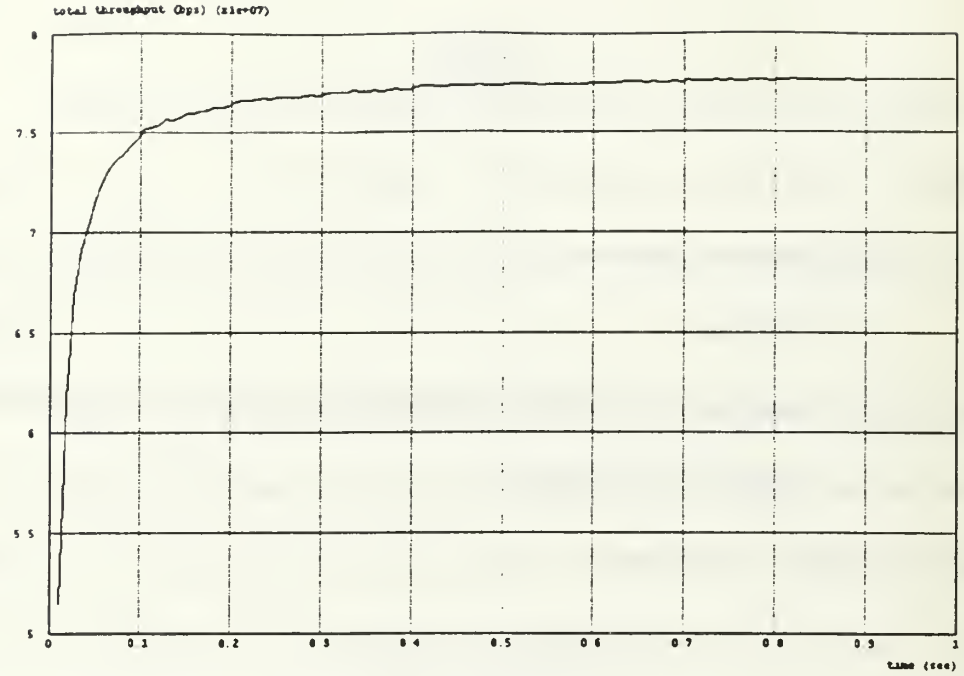

Figure 33. Synchronous Throughput: BW Exceeds Packet Transmission Time

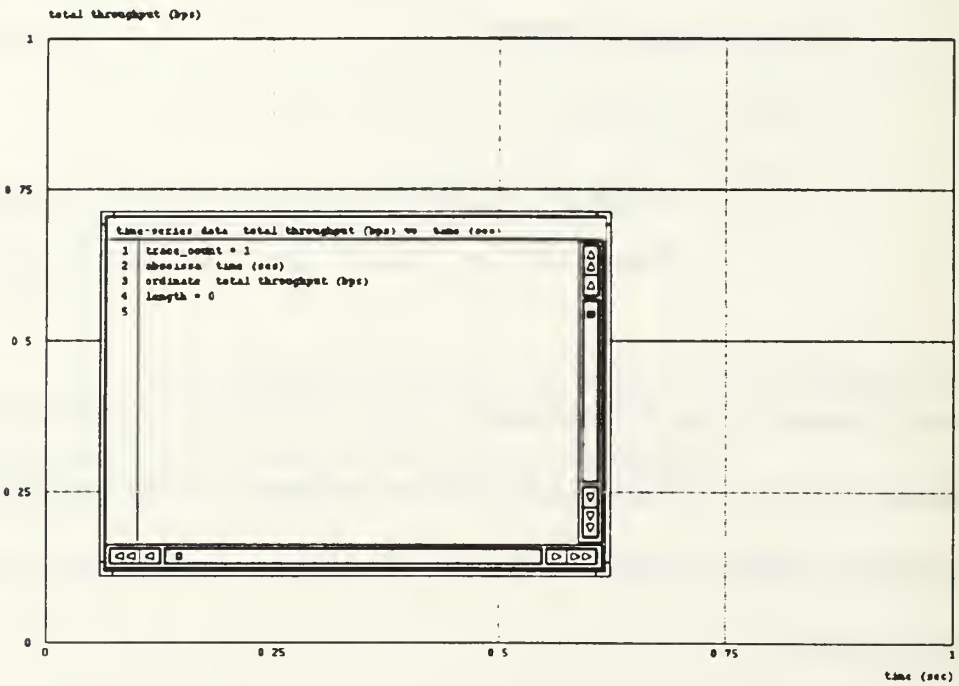

Figure 34. Synchronous Throughput: Packet Transmission Time Exceeds BW

2. Results

Figure 36 illustrates the resulting throughput, 68.48 Mbps, which is 87.79% of the offered load. This agrees quite well with the predicted throughput.

E. MULTICASTING

1. Overview

The multicasting function is designed to assign to each packet a randomly chosen number of destination addresses, then to use this selected number as an index for a loop in which a different destination address is assigned on each iteration. Upon completion of the loop, a vector array of ones and zeros has been created, in which a one in position *i* indicates the packet is addressed to station `f`$_i$. Of course, a real transmitter would send many consecutive packets to the same set of addresses, in streams that comprise messages. However, this modification is in keeping with the original model's actions, which randomly assigned destinations on a packet-by-packet basis.

Preliminary tests indicated the model is not fully developed. Because no change was made to the statistics generation mechanism in the sink process model, the new model as written should have resulted in packets being counted toward total throughput each time they were received at a destination, giving an inflated throughput computation. On the other hand, this inflated figure could possibly be of value as a measure of "virtual throughput." In any event, tests indicate that throughput statistics are not being counted correctly. However, a study of the model's behavior using the debug tool in conjunction with "`printf`" statements

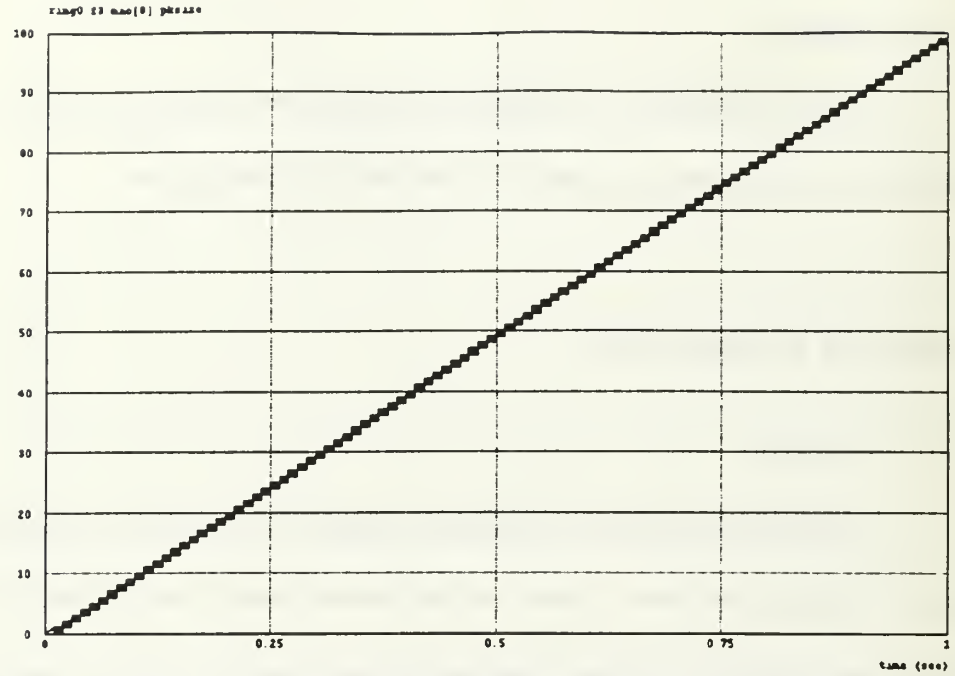

Figure 35. Accumulating Synchronous Packets in MAC

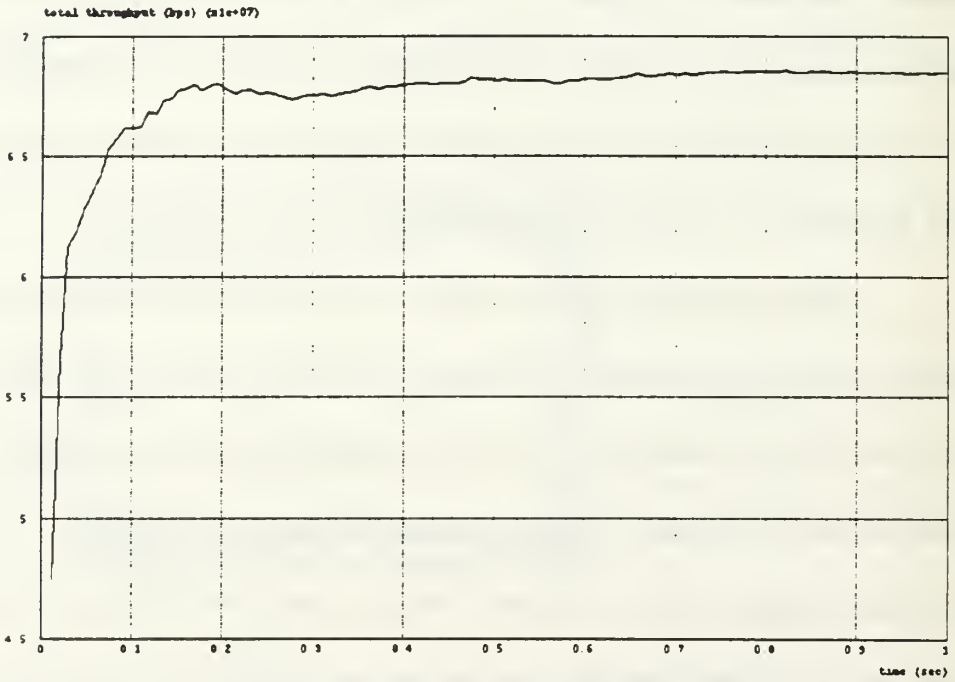

Figure 36. Asynchronous Throughput

indicated that no packets were being lost, and all stations seemed to carry out correctly the operation of receiving a multicast packet, saving the information, and passing the packet on to the next station.

Three tests are described here. The first is an attempt to reproduce the results generated for one of the simulations used in Section A, then uses the same input parameters to see if any differences exist in the handling of synchronous-only and asynchronous-only traffic. The second test compares all synchronous throughput with all asynchronous throughput. The third test is a verification that no packets are actually being lost. The fourth test is an attempt to generate a plot of expected throughput when one station broadcasts all its traffic.

The tests generally indicate that the multicast capable model, when limited to single addressing, does not behave in the same manner as the single-address-only model. In particular, throughput is lower than expected. Further development will be necessary to make the multicasting model a reliable tool.

2. First Test

a. Setup

Figure 37 displays the throughput plots resulting from the third run in the set of simulations used to construct Figure 27. This test is based on the idea that if the multicast-capable model is limited to single addressing, and given the same input parameters provided in the earlier test, then the resulting throughput ought to be identical for total asynchronous and synchronous traffic. Therefore, the following inputs were used for a 50 station, 50 km. LAN using the multicasting station model.

Ten stations (`f0, f5, f10, f15, f20, f25, f30, f35, f40, f45`) transmit only synchronous traffic in 512-bit packets with a constant arrival rate of 6000 packets/sec, resulting in a synchronous offered load of 30.72 Mbps. The remaining 40 stations transmit asynchronous traffic only, in 1000-bit packets at an arrival rate of 750 packets/sec, for an asynchronous offered load of 30 Mbps. TTRT is set to 10.7 ms. For all stations, the Environment file attributes "`min num addees`" and "`max num addees`" are both set to 1, enforcing the limit of one destination address per packet. Appendix H shows the Environment file used here.

b. Results

Figure 38 illustrates the resulting throughput. Table 1 summarizes the throughput results for the two simulations. The unpredictable throughput of the multicast-capable model's synchronous and asynchronous modes would suggest some logic error in coding. The reduction in overall throughput suggests perhaps some difficulty with the simulation's timing mechanism.

Table 4. THROUGHPUT COMPARISON, FIRST TEST

	Offered Load	*Single Address Capable*	*Multicast Capable*
Total	60.72 Mbps	60.60 Mbps	55.00 Mbps
Synchronous	30.72 Mbps	30.80 Mbps	43.00 Mbps
Asynchronous	30.00 Mbps	30.00 Mbps	12.00 Mbps

A real FDDI station would not generate destination addresses in the manner coded into the simulation model; the time required would be unreasonable. On the

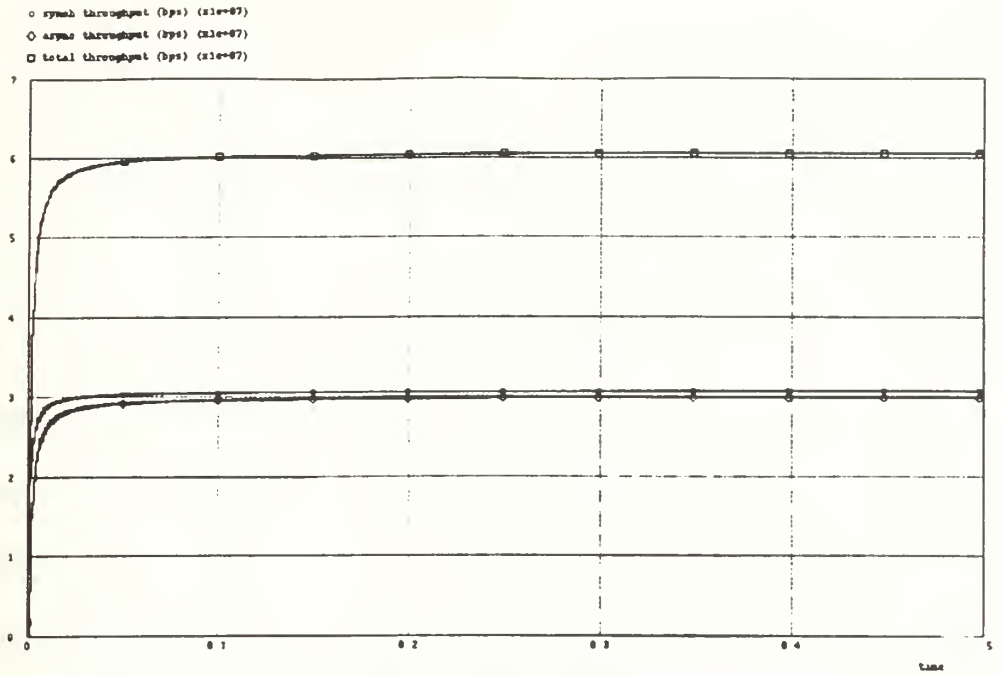

Figure 37. Throughputs: Single Addressing Only Stations

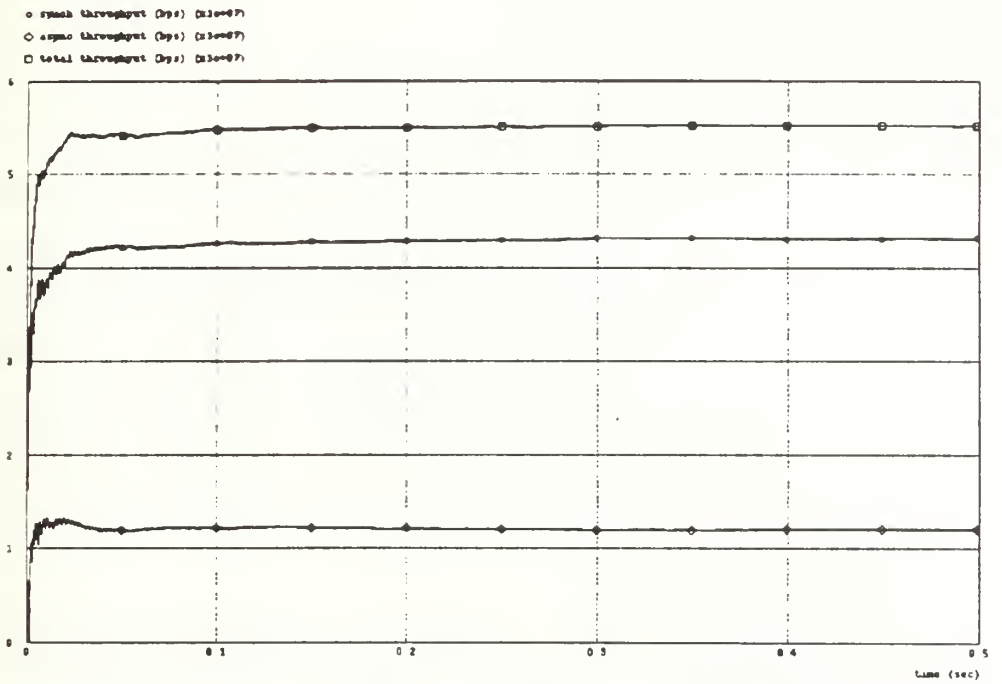

Figure 38. Throughput: Multicast Capable Stations, Single Addressing Mode

other hand, although the simulation's execution is noticeably slowed by the extra events generated in the destination address assignment loop, the throughput rate should not be affected because the simulated passage of time is controlled by the Kernel Procedures. That is, one microsecond does not pass until the simulation has completed one microsecond's worth of events, at all points on the LAN. This is how simultaneous events around the LAN are conducted, one at a time. For example, a new packet arrival at station `f7`, and a packet destruction at station `f23`, and a token release at station `f0` may all occur simultaneously. A study of the sequence of events, revealed through use of the debug facility, shows that the simulation's clock is incremented after these events are all completed, thereby modeling simultaneous events. The eight percent reduction in total throughput may indicate the simulation timer is proceeding without waiting for the completion of the loop.

3. Second test

a. Setup

This test was intended to compare the throughput resulting from two nearly identical simulations, in which the first involved only synchronous traffic, and the second involved only asynchronous traffic. A ten-station LAN was created, with all stations having the same packet arrival rate (750 packets/sec) and packet size (1000 bits). As in the previous test, destination addresses were limited to one station per packet. TTRT was set to 4.0 ms.

b. Results

Figure 39 and Figure 40 illustrate the throughputs resulting from 7.5 Mbps offered load of all-synchronous and all-asynchronous traffic, respectively. Interestingly, both plots are identical, suggesting perhaps that the disparate results of the previous test indicate

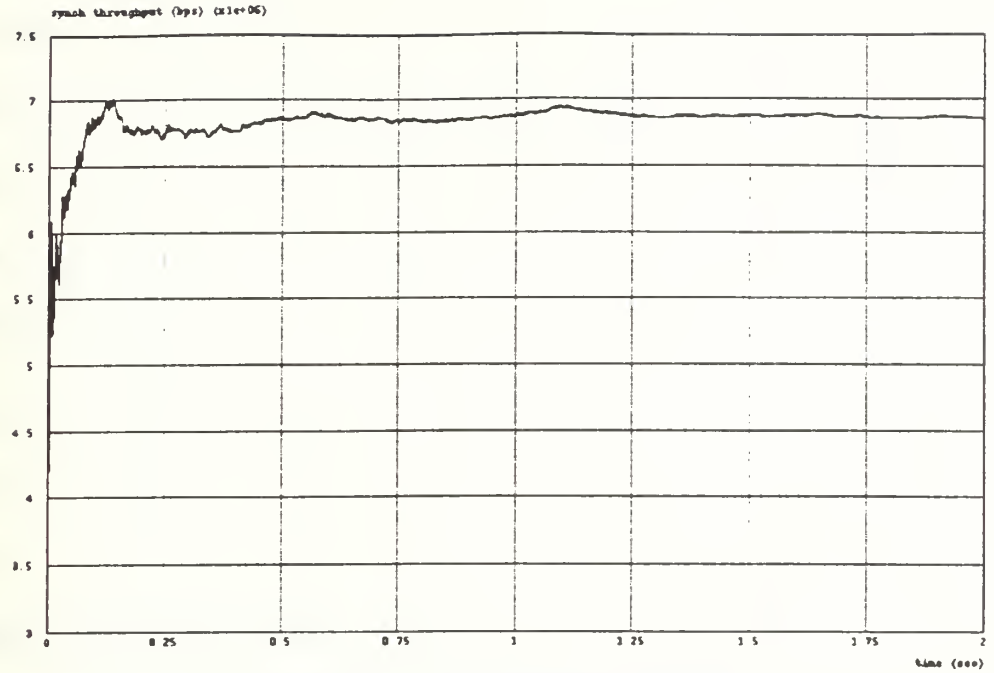

Figure 39. Synchronous Throughput

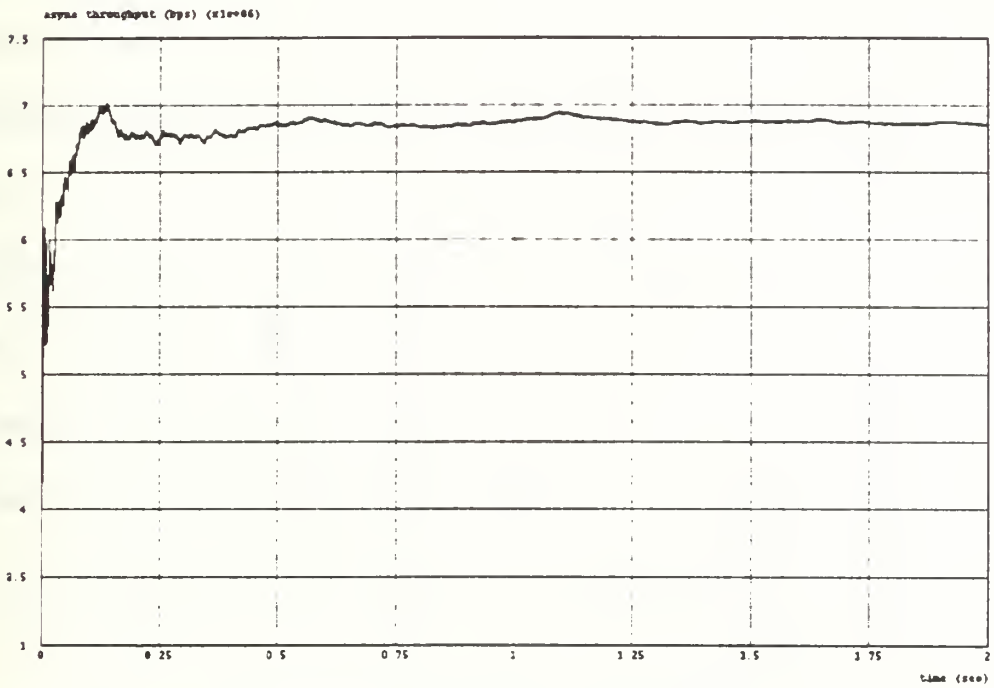

Figure 40. Asynchronous Throughput

a coding error whenever mixed loads of synchronous and asynchronous traffic must be tracked simultaneously. Also noteworthy is the fact that the resulting throughput in both cases, 6.86 Mbps, was 91.5% of the offered load. This is comparable with the previous test: the 55 Mbps throughput in Figure 38 represents 91.7% of the 60.72 Mbps offered load. Again, this suggests a coordination problem between the addressing loop and the simulation's timekeeping function.

4. Third Test

a. Setup

This test was intended to compare the throughput from two different perspectives available. One result was generated in the familiar accounting procedure conducted in the sink process model, and the other result used the Probe Editor to place a monitoring probe on one of the station transmitter nodes, `phy_tx`. If the transmissions are arranged so that all traffic passes this node, then the result should be two identical throughputs.

The same LAN of ten multicast-capable stations was used, with only station `f9` transmitting, and with all of its traffic directed to station `f8`. In between, a probe was placed on the transmitter node of station `f7`. The packet arrival rate was 7500 packets/sec., and the packet length was 1000 bits, giving an offered load of 7.5 Mbps, all of which was asynchronous, with the full range of prioritization available. TTRT was 4.0 ms.

b. Results

Figure 41 shows a roughly constant difference of approximately 0.15 Mbps between the throughput monitored from the physical transmitter and that calculated in the

receiving station's sink process. The irregular plots are unusual, since the simulations normally show a smooth steady state after less than half a second. Significantly, both curves remain close to the offered load, though they are jagged.

The 40 bits per packet overhead (created in the Parameter Editor, where the fields "fc", "src_addr," and "dest_addr" are assigned sizes of 8, 16 and 16 bits, respectively) is a possible source of disparity, although a difference of 0.3 Mbps would be expected in that case (7500 x 40). A study of the sink process model shows that overhead is not included in throughput calculations, while the probes do count the bits in the encapsulating packet structure. As before, the problem requires more study.

5. Fourth Test

a. Setup

A final test of the expected throughput of the multicast-capable model actually observed the multicasting facility, or more exactly, the broadcasting facility of the model. Again, a ten station LAN was used, with only one station transmitting, with TTRT set to 4.0. This transmitter, station f7, generated 7500 packets at 1000 bits per packet, for a total offered load of 7.5 Mbps, all asynchronous. In addition, each packet was addressed to all nine of the other stations, which would be expected to yield a throughput of 67.5 Mbps. (The throughput statistics are gathered by comparing timestamps at the receipt of a packet with the packet's creation time, which is carried as a field in the fddi_llc_fr packet format).

b. Results

Figure 42 illustrates the actual result. The probe at station `f7`'s "`phy_tx`" node reflects the offered load of 7.5 Mbps with reasonable accuracy. However, the expected throughput of 67.5 Mbps was not nearly realized. The throughput plotted, 48.9 Mbps, represents 72.44% of the predicted amount. Again, this indicates a need to further develop the model, and to better define the meanings of throughput and offered load when packets are addressed to more than one station at a time.

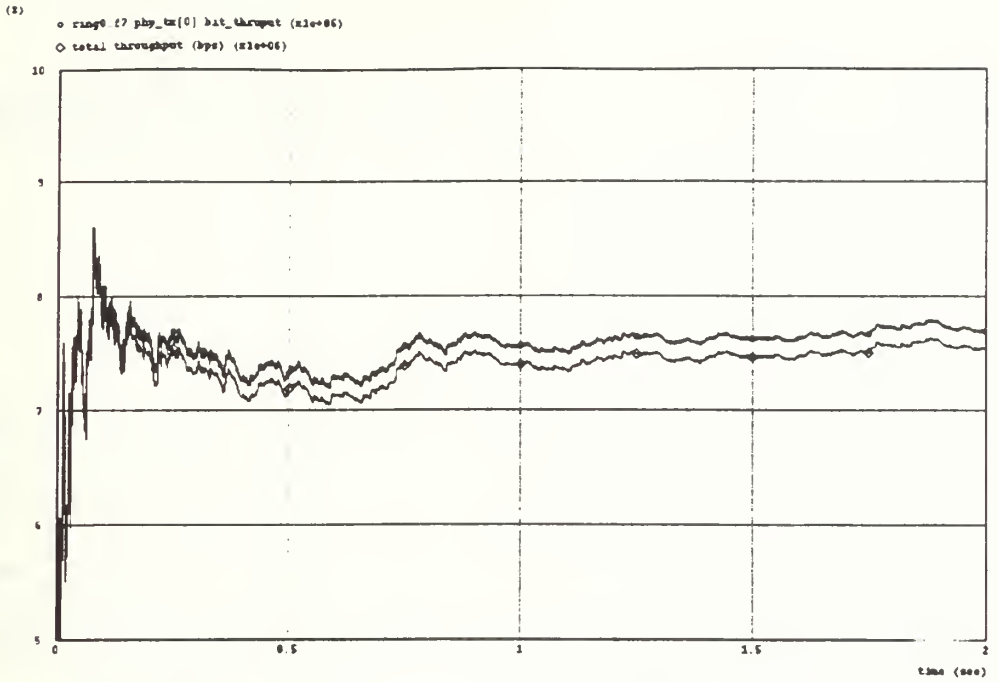

Figure 41. Throughput from Two Vantage Points

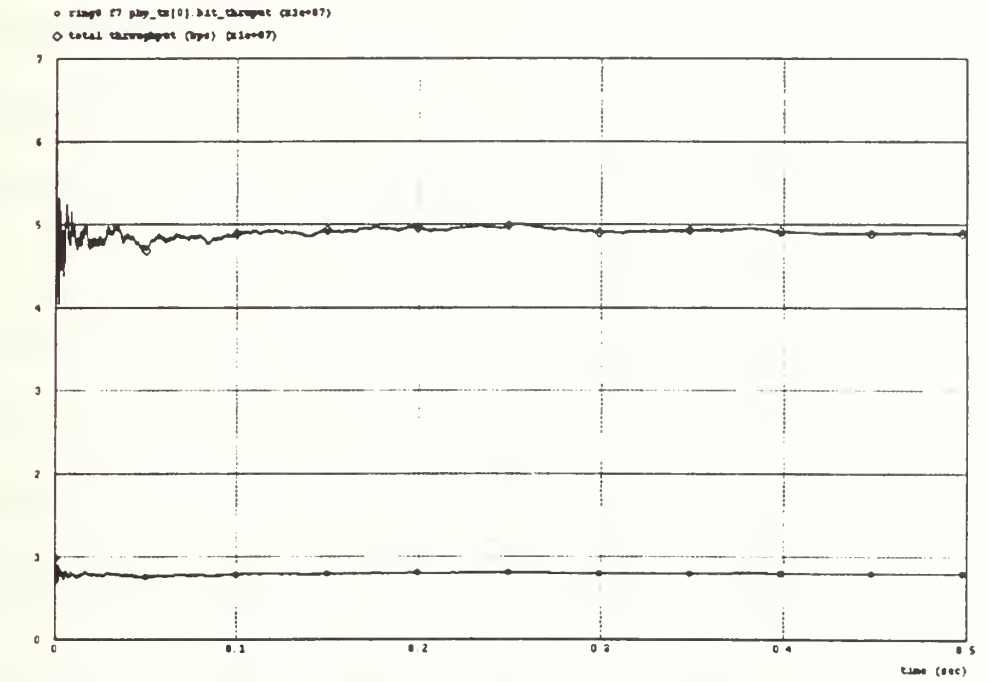

Figure 42. Broadcast from One Station

V. CONCLUSIONS AND RECOMMENDATIONS

A. CONCLUSIONS

This thesis has been directed to two main purposes: to explain the use and operation of OPNET®'s FDDI LAN simulation, and to describe the changes that were installed in the model in order to make it a more useful, accurate and versatile tool for the Common Data Link project. Both goals share another common objective, which is to develop and document the "corporate knowledge" of the CDL working group which will continue to work with the model studied here.

One immediate conclusion is the observation that OPNET is a powerful and flexible tool, but it requires much time and study to be used effectively. Where desired model attributes are lacking, the patient user may code his or her own. The following are the accomplishments documented with this thesis:

1. Throughput, mean delay, and end-to-end delay data are recorded separately and may be displayed separately for synchronous packet traffic and each priority level of asynchronous packet traffic. In addition, two new scalar plot axes, "Asynchronous Offered Load" and "Total Offered Load" are available for the display of LAN performance data in relation to usage.
2. The FDDI station model is capable of randomly choosing different priority threshold settings to assign to generated asynchronous packets, where before the

model was bound to one setting per simulation. This modification allows an additional measure of flexibility in assigning the transmission characteristics of an FDDI LAN.

3. A rudimentary linking node has been created, which accumulates received traffic in buffers for eventual transmission to another LAN.

4. A multicasting capability has been added to the FDDI station model, enabling packets to be addressed to more than one destination station. Preliminary tests and studies indicate the model correctly generates, transmits, receives, and disposes of the multicast packets, although unexpected throughput data suggest the model possibly has coding inaccuracies or improper interfacing of simulation timing and destination address generation.

5. A modest series of preliminary tests verifies the continuing valid operation of the modified models, for the most part.

6. A significant number of unexpected features of OPNET's FDDI LAN model are documented, for the benefit of those researchers continuing to work on the CDL proje.

B. RECOMMENDATIONS

Because the eventual goal of the work begun in this thesis is the development of a large scale model simulation for a Network Interface, the development of the FDDI LAN model may be expected to continue. The following are some possible areas for further development:

1. As mentioned, the multicasting feature yields unexpectedly low throughput. The complex nature of the changes made to the code to enable multicasting suggests the likelihood of some logic error or timing interface discrepancy. This ought to be found and corrected.

2. The testing presented in this thesis is only preliminary, and could well be expanded upon. Because of time constraints, delay characteristics were not addressed at all.

3. The buffered subqueues of the preliminary bridging model are infinite by default, but may be assigned limits through the use of on-screen attributes. The code currently has no provision for the proper disposal of packets lost to buffer overflows; lost frames will fill the host terminal's memory until none remains. A related but separate issue is the fact that there is no retransmission protocol in effect for the current model.

APPENDIX A. FILE RETRIEVAL VIA FTP

The following is an excerpt from a screen dialogue demonstrating access to MIL 3 Inc.'s bulletin board, movement among subdirectories, and retrieval of files.

```
sun24:/home3/nix
% ftp
ftp> open
(to) mil3.com
Connected to mil3.com.
220 rmaxwell FTP server (SunOS 4.1) ready.
Name (mil3.com:nix): anonymous
331 Guest login ok, send ident as password.
Password: nix@ece.nps.navy.mil
230 Guest login ok, access restrictions apply.
ftp> dir
200 PORT command successful.
150 ASCII data connection for /bin/ls (131.120.20.124,2926) (0 bytes).
total 18
-rw-r--r--  1 0      100   2110 Jul 21 00:32 FTP.instructions
-r--r--r--  1 0      1     5135 Jan 21  1992 README.OPSIG
dr-xr-xr-x  2 0      1      512 Jan 14  1992 bin
dr-xr-xr-x  2 0      1      512 Jan 14  1992 dev
dr-xr-xr-x  2 0      1      512 Jan 14  1992 etc
drwxr-xr-x  5 0      1      512 Oct 27  1992 examp
drwx--x--x  2 0      1      512 Dec 20 16:40 incoming
dr-xr-xr-x 10 0      1      512 Dec 23 17:34 model_depot
drwxr-xr-x 10 0      1      512 Dec 15 10:09 patches
drwxrwxrwx  3 0      1      512 Dec 29 03:05 tmp
dr-xr-xr-x  3 0      1      512 Jan 14  1992 usr
226 ASCII Transfer complete.
700 bytes received in 3.1 seconds (0.22 Kbytes/s)
ftp> get README.OPSIG
200 PORT command successful.
150 ASCII data connection for README.OPSIG (131.120.20.124,2927) (5135
   bytes).
226 ASCII Transfer complete.
local: README.OPSIG remote: README.OPSIG
5257 bytes received in 6.8 seconds (0.75 Kbytes/s)
ftp> cd patches
250 CWD command successful.
ftp> dir
200 PORT command successful.
150 ASCII data connection for /bin/ls (131.120.20.124,2928) (0 bytes).
total 8
```

```
drwxr-xr-x  6 101    1           512 Mar 10  1993 2.3.L
drwxr-xr-x  4 101    100         512 Apr  2  1993 2.3.Lhp
drwxr-xr-x  6 101    1           512 Sep 24 22:42 2.3.M
drwxr-xr-x  4 101    1           512 Dec  6 18:37 2.4.A
drwxr-xr-x  5 101    1           512 Dec 17 22:32 2.4.B
drwxr-xr-x  4 101    1           512 Jan 29  1993 2.4.beta2
drwxr-xr-x 13 101    1           512 Apr 26  1993 2.4.beta3
drwxr-xr-x  5 101    1           512 Jun 10  1993 2.4.beta4
226 ASCII Transfer complete.
511 bytes received in 0.63 seconds (0.8 Kbytes/s)
250 CWD command successful.
ftp> cd 2.4.B
250 CWD command successful.
ftp> dir
200 PORT command successful.
150 ASCII data connection for /bin/ls (131.120.20.124,2934) (0 bytes).
total 3
drwxr-xr-x  2 101 100           512 Dec 15 10:11 opbug_2081
drwxr-xr-x  2 0       1            512 Dec 17 22:33 tsup_3139
drwxr-xr-x  2 0       1            512 Dec 17 20:15 tsup_3182
226 ASCII Transfer complete.
205 bytes received in 0.28 seconds (0.7 Kbytes/s)
ftp> cd opbug_2081
250 CWD command successful.
ftp> dir
200 PORT command successful.
150 ASCII data connection for /bin/ls (131.120.20.124,2935) (0 bytes).
total 57
-rw-r--r--  1 101    100         3988 Dec 15 07:12 README
-rw-r--r--  1 101    100         53379 Dec 15 07:12 fddi_mac.pr.m
226 ASCII Transfer complete.
141 bytes received in 0.22 seconds (0.62 Kbytes/s)
ftp> get README
200 PORT command successful.
150 ASCII data connection for README (131.120.20.124,2936) (3988
   bytes).
226 ASCII Transfer complete.
local: README remote: README
4092 bytes received in 5 seconds (0.81 Kbytes/s)
ftp> bin
200 Type set to I.
ftp> get fddi_mac.pr.m
200 PORT command successful.
150 ASCII data connection for fddi_mac.pr.m (131.120.20.124,2937)
   `(53379 bytes).
226 ASCII Transfer complete.
local: fddi_mac.pr.m remote: fddi_mac.pr.m
53388 bytes received in 64 seconds (0.81 Kbytes/s)
ftp> quit
```

APPENDIX B. PACKET AND ICI FRAME STRUCTURES

The following packet structures are used in the FDDI LAN model.

A. PACKET FORMATS

1. "fddi_llc_fr"

Field Name	Type	Size (bits)	Default Value	Default Set
cr_time	double	0	0.0	set
pri[Note]	integer	0	0	unset

Note: Added to allow generation of statistics related to prioritized traffic

2. "fddi_mac_fr"

Field Name	Type	Size (bits)	Default Value	Default Set
fc	integer	8	-	unset
src_addr	integer	16	-	unset
dest_addr	integer	16	-	unset
info	packet	-1	-	unset
svc_class	integer	0	-	unset
pri	integer	0	-	unset
tk_class	integer	0	-	unset

3. "fddi_mac_tk"

Field Name	Type	Size (bits)	Default Value	Default Set
fc	integer	8	-	unset
class	integer	0	-	unset
res_station	integer	0	-	unset

B. ICI FORMATS

1. "fddi_mac_ind"

Attribute Name	Type	Default
src_addr	integer	0
dest_addr	integer	0

2. "fddi_mac_req"

Attribute Name	Type	Default
svc_class	integer	0
dest_addr	integer	0
pri	integer	0
tk_class	integer	0

APPENDIX C. EXAMPLE ENVIRONMENT FILE

FOR 32-STATION FDDI LAN

```
# fddi32.ef
# sample simulation configuration file for fddi example model
# 32 station network

#*** Attributes related to loading used by "fddi_gen" ***

# station addresses
*.*.f0.mac.station_address: 0
*.*.f1.mac.station_address: 1
*.*.f2.mac.station_address: 2
*.*.f3.mac.station_address: 3
*.*.f4.mac.station_address: 4
*.*.f5.mac.station_address: 5
*.*.f6.mac.station_address: 6
*.*.f7.mac.station_address: 7
*.*.f8.mac.station_address: 8
*.*.f9.mac.station_address: 9
*.*.f10.mac.station_address: 10
*.*.f11.mac.station_address: 11
*.*.f12.mac.station_address: 12
*.*.f13.mac.station_address: 13
*.*.f14.mac.station_address: 14
*.*.f15.mac.station_address: 15
*.*.f16.mac.station_address: 16
*.*.f17.mac.station_address: 17
*.*.f18.mac.station_address: 18
*.*.f19.mac.station_address: 19
*.*.f20.mac.station_address: 20
*.*.f21.mac.station_address: 21
*.*.f22.mac.station_address: 22
*.*.f23.mac.station_address: 23
*.*.f24.mac.station_address: 24
*.*.f25.mac.station_address: 25
*.*.f26.mac.station_address: 26
*.*.f27.mac.station_address: 27
*.*.f28.mac.station_address: 28
*.*.f29.mac.station_address: 29
*.*.f30.mac.station_address: 30
*.*.f31.mac.station_address: 31
```

```
*.*.*.mac.ring_id :0

# destination addresses for random message generation
"*.*.llc_src.low dest address" :    0
"*.*.llc_src.high dest address" :   31

#"top.ring0.f0.llc_src.low dest address"  :
#"top.ring0.f0.llc_src.high dest address" :

# arrival rate(frames/sec), and message size (bits) for random message
# generation at each station on the ring.
"*.*.*.arrival rate" :        200
"*.*.*.mean pk length" :      500

# set the proportion of asynchronous traffic
# a value of 1.0 indicates all asynchronous traffic
"*.*.*.async_mix" :        1.0

#*** Ring configuration attributes used by "fddi_mac" ***

# allocate percentage of synchronous bandwidth to each station
# this value should not exceed 1 for all stations combined; OPNET does
not
# enforce this; 01FEB94: this must be less than 1; see equation below

"*.*.mac.sync bandwidth" :    0.0
# "*.f0.mac.sync bandwidth" : .0935487

"*.*.mac.T_Req" : .010

# Index of the station which initially launches the token
"spawn station":  0

# Delay incurred by packets as they traverse a station's ring interface
# see Powers, p. 351 for a discussion of this (Powers gives 1usec,
# but 60.0e-08 agrees with Dykeman & Bux)
station_latency:  60.0e-08

# Propagation Delay separating stations on the ring.
prop_delay:       5.085e-06

#*** Simulation related attributes

# Token Acceleration Mechanism enabling flag.
# It is reccomended that this mechanism be enabled for most situations
accelerate_token: 1
seed:                10

# Run control attributes
duration:          5
```

```
verbose_sim:        TRUE
upd_int:            .1
os_file:            fddi32mod

ov_file:            fddi32mod
# Opnet Debugger (odb) enabling attribute
# debug:            TRUE
```

APPENDIX D. DEBUG TOOL EXCERPT

The following is an excerpt from the debugger in fulltrace, showing the arrival and reception of a message traffic packet at its destination address. In this case, station `f11` has sent a packet to station `f31`. The packet has been passed "hand to hand" through each station in between, with the simulator enacting every step along the way. Note that when reception is complete, the simulation returns to station `f11`, which still has the token.

```
_____________________ (ODB 2.4.A: Event)_____________________________

      * Time   :  0.995744186808 sec, [00d 00h 00m 00s . 995ms 744us
        186ns 808ps]
      * Event  :  execution ID (461704), schedule ID (#502933), type
        (self intrpt)
      * Source :  execution ID (461703), top.ring0.f30.phy_tx (pt-pt
        transmitter)
      * Data   :  code (0)
    > Module :  top.ring0.f30.phy_tx (pt-pt transmitter)
odb> next

         * Kernel Action: Pt-Pt Transmitter object
           Completing transmission of packet(s)
           channel     (0)
           packet ID   (4459)

_____________ (ODB 2.4.A: Event) ____________________________________

      * Time   :  0.995744186808 sec, [00d 00h 00m 00s . 995ms 744us
        186ns 808ps]
      * Event  :  execution ID (461705), schedule ID (#502934), type
        (remote intrpt)
      * Source :  execution ID (461703), top.ring0.f30.phy_tx (pt-pt
        transmitter)
      * Data   :  code (0)
      > Module :  top.ring0.f31.phy_rx (pt-pt receiver)

odb> next

         * Kernel Action: Pt-Pt Receiver object
           Beginning reception of packet
           channel     (0)
```

```
            packet ID   (4459)

_____________ (ODB 2.4.A: Event) ____________________________

    * Time    :  0.995744186808 sec, [00d 00h 00m 00s . 995ms 744us 186ns
   808ps]
    * Event   :  execution ID (461706), schedule ID (#502935), type
   (remote intrpt)
    * Source  :  execution ID (461703), top.ring0.f30.phy_tx (pt-pt
   transmitter)
    * Data    :  code (0)
    > Module  :  top.ring0.f31.phy_rx (pt-pt receiver)

odb> next

         * Kernel Action: Pt-Pt Receiver object
            Completing reception of packet
            channel     (0)
            packet ID   (4459)

         * Kernel Action: Pt-Pt pipeline
            Calling (error) pipeline stage
            packet ID   (4459)

         * op_td_get_int (pkptr, tda_index)
            packet ID   (4459)
            TDA attribute    (OPC_TDA_PT_LINK_OBJID)
            TDA value   (82)

         * op_ima_obj_attr_get (objid, attr_name, value_ptr)
            objid    (82)
            attr_name   (ber)
            value    (0.0)

         * op_pk_total_size_get (pkptr)
            packet ID   (4459)
            total size  (25040)

         * op_td_set_int (pkptr, tda_index, value)
            packet ID   (4459)
            TDA attribute    (OPC_TDA_PT_NUM_ERRORS)
            TDA value   (0 bit errs)

         * Kernel Action: Pt-Pt pipeline
            Calling (ecc) pipeline stage
            packet ID   (4459)

         * op_td_is_set (pkptr, tda_index)
            packet ID   (4459)
            TDA attribute    (OPC_TDA_PT_ND_FAIL)
```

```
            tda is set  (false)

         * op_td_get_int (pkptr, tda_index)
            packet ID   (4459)
            TDA attribute    (OPC_TDA_PT_RX_OBJID)
            TDA value   (734)

         * op_ima_obj_attr_get (objid, attr_name, value_ptr)
            objid    (734)
            attr_name   (ecc threshold)
            value    (0.0)

         * op_pk_total_size_get (pkptr)
            packet ID   (4459)
            total size  (25040)

         * op_td_get_int (pkptr, tda_index)
            packet ID   (4459)
            TDA attribute    (OPC_TDA_PT_NUM_ERRORS)
            TDA value   (0 bit errs)

         * op_td_set_int (pkptr, tda_index, value)
            packet ID   (4459)
            TDA attribute    (OPC_TDA_PT_PK_ACCEPT)
            TDA value   (1)

         * Kernel Action: Pt-Pt Receiver object
            Packet successfully received
            channel     (0)
            packet ID   (4459)

___________ (ODB 2.4.A: Event) ____________________________

    * Time   :  0.995744186808 sec, [00d 00h 00m 00s . 995ms 744us 186ns
   808ps]
    * Event  :  execution ID (461707), schedule ID (#502936), type
   (stream intrpt)
    * Source :  execution ID (461706), top.ring0.f31.phy_rx (pt-pt
   receiver)
    * Data   :  instrm (0), packet ID (4459)
    > Module :  top.ring0.f31.mac (queue)

odb> next

         * invoking process ("fddi_mac")

         _________state (IDLE): exit executives_________

         * op_intrpt_type ()
            intrpt type (stream intrpt)
```

```
* op_intrpt_strm ()
   active strm (0)

* op_pk_get (instrm_index)
   strm. index (0)
   packet ID   (4459)

* op_pk_nfd_get (pkptr, fd_name, value_ptr)
   packet ID   (4459)
   field name  (fc)
   value    (0)

* op_intrpt_type ()
   intrpt type (stream intrpt)

* op_intrpt_type ()
   intrpt type (stream intrpt)

* op_intrpt_strm ()
   active strm (0)

_______state (FR_RCV): enter executives_________

* op_pk_nfd_get (pkptr, fd_name, value_ptr)
   packet ID   (4459)
   field name  (src_addr)
   value    (11)

_______state (FR_RCV): exit executives_________

____state (FR_REPEAT): enter executives_________

* op_pk_nfd_get (pkptr, fd_name, value_ptr)
   packet ID   (4459)
   field name  (dest_addr)
   value    (31)

* op_pk_total_size_get (pkptr)
   packet ID   (4459)
   total size  (25040)

* op_pk_nfd_get (pkptr, fd_name, value_ptr)
   packet ID   (4459)
   field name  (info)
   value    (pk id (4458))

* op_ici_attr_set (iciptr, attr_name, attr_value)
   ICI id      (81)
   attr name   (src_addr)
   value    (11)
```

```
          * op_ici_attr_set (iciptr, attr_name, attr_value)
             ICI id       (81)
             attr name    (dest_addr)
             value     (31)

          * op_ici_install (iciptr)
             ICI ID       (81)

          * op_pk_send_delayed (pkptr, outstrm_index, delay)
             packet ID    (4458)
             stream index    (1)
             delay     (0.0002504)

          * op_pk_destroy (pkptr)
             packet ID    (4459)

          * Kernel Action: Destroying Packet
             packet ID    (4459)

          _____state (FR_REPEAT): exit executives_________

          _________state (IDLE): enter executives_________

          * returning from process ("fddi_mac")

___________ (ODB 2.4.A: Event) ____________________________

    * Time    :  0.995886571808 sec, [00d 00h 00m 00s . 995ms 886us 571ns
   808ps]
    * Event   :  execution ID (461708), schedule ID (#502837), type
   (stream intrpt)
    * Source  :  execution ID (461607), top.ring0.fll.mac (queue)
    * Data    :  instrm (0), packet ID (4457)
    > Module  :  top.ring0.fll.phy_tx (pt-pt transmitter)

odb> next
```

APPENDIX E. MAC "C" CODE:

"fddi_mac_mult.pr.c"

The line numbering in this appendix is used for reference within this thesis only, and does not correspond with that seen in OPNET®'s text editors.

```
1   /* Process model C form file: fddi_mac_mult.pr.c */
2   /* Portions of this file Copyright (C) MIL 3, Inc. 1992 */

3   /* OPNET system definitions */
4   #include <opnet.h>
5   #include "fddi_mac_mult.pr.h"
6   FSM_EXT_DECS

7   /* Header block */
8   /* Define a timer structure used to implement */
9   /* the TRT and THT timers.  The primitives defined to */
10  /* operate on these timers can be found in the */
11  /* function block of this process model. */
12  typedef struct
13     {
14     int            enabled;
15     double            start_time;
16     double            accum;
17     double            target_accum;
18     } FddiT_Timer;

19  /* 08FEB94: define the number of stations here. -Nix */
20  #define NUM_STATIONS 50

21  /* Declare certain primitives dealing with timer.s */
22  double         fddi_timer_remaining ();
23  FddiT_Timer*   fddi_timer_create ();
24  double         fddi_timer_value ();

25  /* Scratch strings for trace statements */
26  char        str0 [512], str1 [512];

27  /* define constants particular to this implementation */
28  #define FDDI_MAX_STATIONS        512
```

```
29  /* define possible values for the frame control field */
30  #define FDDI_FC_FRAME            0
31  #define FDDI_FC_TOKEN            1

32  /* define possible service classes for frames */
33  #define FDDI_SVC_ASYNC           0
34  #define FDDI_SVC_SYNC            1

35  /* define input stream indices */
36  #define FDDI_LLC_STRM_IN         1
37  #define FDDI_PHY_STRM_IN         0

38  /* define output stream indices */
39  #define FDDI_LLC_STRM_OUT        1
40  #define FDDI_PHY_STRM_OUT        0

41  /* define token classes */
42  #define FDDI_TK_NONRESTRICTED    0
43  #define FDDI_TK_RESTRICTED       1

44  /* Ring Constants */
45  #define FDDI_TX_RATE          1.0e+08
46  #define FDDI_SA_SCAN_TIME        28.0e-08

47  /* Token transmission time: based on 6 symbols plus 16 symbols of
48  preamble */
49  #define FDDIC_TOKEN_TX_TIME         88.0e-08

50  /* Codes used to differentiate remote interrupts */
51  #define FDDIC_TRT_EXPIRE         0
52  #define FDDIC_TK_INJECT          1

53  /* Define symbolic expressions used on transition */
54  /* conditions and in executive statements. */
55  #define TRT_EXPIRE (op_intrpt_type () == OPC_INTRPT_REMOTE &&
56     op_intrpt_code () == FDDIC_TRT_EXPIRE)

57  #define TK_RECEIVED phy_arrival && frame_control == FDDI_FC_TOKEN

58  #define RC_FRAME phy_arrival && frame_control == FDDI_FC_FRAME

59  #define FRAME_ARRIVAL                                    \
60     op_intrpt_type () == OPC_INTRPT_STRM &&               \
61     op_intrpt_strm () == FDDI_LLC_STRM_IN

62  #define STRIP  my_address == src_addr

63  /* Define the maximum value for ring_id.  This is the */
64  /* maximum number of FDDI rings that can exist in a */
65  /* simulation.  Note that if this number is changed, */
```

```
66  /* the initialization for fddi_claim_start below must */
67  /* also be modified accordingly.        */
68  #define FDDI_MAX_RING_ID            8

69  /* Declare the operative TTRT value 'T_Opr' which is the final */
70  /* negotiated value of TTRT.  This value is shared by all stations */
71  /* on a ring so that all agree on its value.        */
72  double      fddi_t_opr [FDDI_MAX_RING_ID];
73  #define     Fddi_T_Opr (fddi_t_opr [ring_id])

74  /* This flag indicates that the negotiation for the final TTRT */
75  /* has not yet begun.  It is statically initialized here, and */
76  /* is reset by the first station which modifies T_Opr. */
77  /* Initialize to 1 for all rings.                */
78  int         fddi_claim_start [FDDI_MAX_RING_ID] = {1,1,1,1,1,1,1,1};
79  #define  Fddi_Claim_Start  (fddi_claim_start [ring_id])

80  /* Declare station latency parameters.        */
81  /* These are true globals, so they do not need to be arrays. */
82  double      Fddi_St_Latency;
83  double      Fddi_Prop_Delay;

84  /* Declare globals for Token Acceleration Mechanism. */
85  /* Hop delay and token acceleration are true globals. */
86  double      Fddi_Tk_Hop_Delay;
87  int         Fddi_Tk_Accelerate = 1;

88  /* These are actually values shared by all nodes on a ring, */
89  /* so they must be defined as arrays.         */
90  double      fddi_tk_block_base_time [FDDI_MAX_RING_ID];
91  #define  Fddi_Tk_Block_Base_Time (fddi_tk_block_base_time [ring_id])

92  int           ddi_tk_block_base_station [FDDI_MAX_RING_ID];
93  #define  Fddi_Tk_Block_Base_Station (fddi_tk_block_base_station
94     [ring_id])

95  int         fddi_tk_blocked [FDDI_MAX_RING_ID];
96  #define  Fddi_Tk_Blocked (fddi_tk_blocked [ring_id])

97  int         fddi_num_stations [FDDI_MAX_RING_ID];
98  #define  Fddi_Num_Stations (fddi_num_stations [ring_id])

99  int         fddi_num_registered [FDDI_MAX_RING_ID];
100 #define  Fddi_Num_Registered (fddi_num_registered [ring_id])

101 Objid    fddi_address_table [FDDI_MAX_RING_ID][FDDI_MAX_STATIONS];
102 #define  Fddi_Address_Table (fddi_address_table [ring_id])

103 /* Below is part of the OPBUG 2081 patch; FB ended here, before. -Nix */
```

```
104  /* Event handles for the TRT are maintained at a global level to  */
105  /* allow token acceleration mechanism to adjust these as necessary  */
106  /* when blocking and reinjecting the token. TRT_handle simply  */
107  /* represents the TRT for the local MAC                  */
108  Evhandle fddi_trt_handle [FDDI_MAX_RING_ID][FDDI_MAX_STATIONS];
109  #define  Fddi_Trt_Handle   (fddi_trt_handle [ring_id])
110  #define  TRT_handle     Fddi_Trt_Handle [my_address]

111  /* Similarly, the TRT data structure is maintained on a global level. */
112  FddiT_Timer*   fddi_trt [FDDI_MAX_RING_ID] [FDDI_MAX_STATIONS];
113  #define  Fddi_Trt (fddi_trt [ring_id])
114  #define  TRT         Fddi_Trt [my_address]

115  /* Registers to record the expiration time of each TRT when token is
116     blocked. */
117  double      fddi_trt_exp_time [FDDI_MAX_RING_ID] [FDDI_MAX_STATIONS];
118  #define  Fddi_Trt_Exp_Time (fddi_trt_exp_time [ring_id])

119  /* the 'Late_Ct' flag is declared on a global level so that it can be
120     */
121  /* set at the tim ewhere the token is injected back into the ring.  */
122  int         fddi_late_ct [FDDI_MAX_RING_ID] [FDDI_MAX_STATIONS];
123  #define  Fddi_Late_Ct   (fddi_late_ct [ring_id])
124  #define  Late_Ct        Fddi_Late_Ct [my_address]

125  /* Convenient macro for setting TRT for a given station and absolute
126     time. */
127  #define TRT_SET(station_id,abs_time) fddi_timer_set (Fddi_Trt
128     [station_id], abs_time - op_sim_time()); Fddi_Trt_Handle [station_id]
129     = op_intrpt_schedule_remote (abs_time, FDDIC_TRT_EXPIRE,
130     Fddi_Address_Table [station_id]);

131  /* State variable definitions */
132  typedef struct
133     {
134     FSM_SYS_STATE
135     int                  sv_ring_id;
136     FddiT_Timer*      sv_THT;
137     double               sv_T_Req;
138     double               sv_T_Pri [8];
139     Objid             sv_my_objid;
140     int                  sv_spawn_token;
141     int                  sv_my_address;
142     Packet*           sv_tk_pkptr;
143     double               sv_sync_bandwidth;
144     double               sv_sync_pc;
145     int                  sv_restricted;
146     int                  sv_res_peer;
147     int                  sv_tk_registered;
148     Ici*              sv_to_llc_ici_ptr;
```

```
149     int                     sv_tk_trace_on;
150     } fddi_mac_mult_state;

151  #define pr_state_ptr ((fddi_mac_mult_state*) SimI_Mod_State_Ptr)
152  #define ring_id             pr_state_ptr->sv_ring_id
153  #define THT                 pr_state_ptr->sv_THT
154  #define T_Req               pr_state_ptr->sv_T_Req
155  #define T_Pri               pr_state_ptr->sv_T_Pri
156  #define my_objid          pr_state_ptr->sv_my_objid
157  #define spawn_token       pr_state_ptr->sv_spawn_token
158  #define my_address          pr_state_ptr->sv_my_address
159  #define tk_pkptr          pr_state_ptr->sv_tk_pkptr
160  #define sync_bandwidth    pr_state_ptr->sv_sync_bandwidth
161  #define sync_pc             pr_state_ptr->sv_sync_pc
162  #define restricted          pr_state_ptr->sv_restricted
163  #define res_peer          pr_state_ptr->sv_res_peer
164  #define tk_registered     pr_state_ptr->sv_tk_registered
165  #define to_llc_ici_ptr    pr_state_ptr->sv_to_llc_ici_ptr
166  #define tk_trace_on       pr_state_ptr->sv_tk_trace_on

167  /* Process model interrupt handling procedure */

168  void
169  fddi_mac_mult ()
170     {
171     /* Packets and ICI's */
172     Packet*         mac_frame_ptr;
173     Packet*         pdu_ptr;
174     Packet*         pkptr;
175     Packet*         data_pkptr;
176     Ici*            ici_ptr;

177     /* Packet Fields and Attributes */
178     int                req_pri, svc_class, req_tk_class;
179     int                frame_control, src_addr;
180     int                pk_len, pri_level;
181     static
182     int                *da_ptr, dest_addr[];

183     /* Token - Related */
184     int                tk_usable, res_station, tk_class;
185     int                current_tk_class;
186     double             accum_sync;

187     /* Timer - Related */
188     double             tx_time, timer_remaining, accum_bandwidth;
189     double             tht_value;

190     /* Miscellaneous */
191     int             i;
```

```
192      int            spawn_station, phy_arrival;
193      char        error_string [512];
194      int            num_frames_sent, num_bits_sent;

195   /* 26DEC93: loop management variables, used in RCV_TK  */
196      /* and ENCAP states. -Nix         */
197      int            NUM_PRIOS;
198      int            punt;
199      int            q_check;

200   /* 08FEB94: case management variables, used in FR_REPEAT. -Nix */
201      int            for_me;
202      int            count_addees;

203   /* 08MAR94: "field holding" variables, used in FR_REPEAT. -Nix */
204      Packet*     info_ptr;

205      FSM_ENTER (fddi_mac_mult)

206      FSM_BLOCK_SWITCH
207         {
208         /*-------------------------------------------------------*/
209         /** state (INIT) enter executives **/
210         FSM_STATE_ENTER_FORCED (0, state0_enter_exec, "INIT")
211            {
212   /* Obtain the station's address .  This is an attribute */
213   /* of this process. Addressing is simplified by */
214   /* simply using integers, and only one mode. */
215   /* This mode is 16 bit addressing unless the */
216   /* packet format 'fddi_mac_fr' is modified. */
217            my_objid = op_id_self();   /* 29DEC93 */
218            op_ima_obj_attr_get (my_objid, "station_address", &my_address);

219   /* Register the station's object id in a global table. */
220   /* This table is used by the mechanism which improves */
221   /* simulation efficiency by 'jumping over' idle periods */
222   /* rather than circulating an unusable token. */
223            fddi_station_register (my_address, my_objid);

224   /* Obtain the station latency for tokens and frames. */
225   /* Default value is set at 100 nanoseconds. */
226            Fddi_St_Latency = 100.0e-09;
227            op_ima_sim_attr_get (OPC_IMA_DOUBLE, "station_latency",
228               &Fddi_St_Latency);

229   /* Obtain the propagation delay separating stations. */
230   /* This value is given in seconds with default value 3.3 microseconds.
231      */
232            Fddi_Prop_Delay = 3.3e-06;
```

```
233             op_ima_sim_attr_get (OPC_IMA_DOUBLE, "prop_delay",
234                &Fddi_Prop_Delay);

235     /* Derive the Delay for a 'hop' of a freely circulating packet. */
236             Fddi_Tk_Hop_Delay = Fddi_Prop_Delay + Fddi_St_Latency;


237     /* The T_Pri [] state variable array supports priority */
238     /* assignments on a station by station basis by */
239     /* establishing a correspondence between integer priority */
240     /* levels assigned to frames and the maximum values of the*/
241     /* Token holding timer (THT) which would allow packets to be*/
242     /* sent.  Eight levels are supported here, but this can easily */
243     /* be changed by redimensioning the priority array. */
244     /* By default all levels are identical here, allowing */
245     /* any frame to make use of the token, so that in fact */
246     /* priority levels are not used in the default case. */
247     /* 01JAN94: (8-i) is a quick attempt to impart different weighting */
248     /* scales on each priority level, and is not necessarily realistic.-Nix
249        */
250             for (i = 0; i < 8; i++)
251                {
252                T_Pri[i] = (double) Fddi_T_Opr/(8.0 - i);  /* 01JAN94 */
253     /* printf("INIT: T_Pri[%d] = %d; Fddi_T_Opr = %d\n",    */
254     /* i, T_Pri[i], Fddi_T_Opr);      */
255                }


256     /*Create the token holding timer (THT) used to restrict the */
257     /* asynchronous bandwidth consumption of the station */
258             THT = fddi_timer_create ();

259     /* Create the token rotation timer (TRT) used to measure the */
260     /* rotations of the token, detect late tokens and initialize */
261     /* the THT timer before asynchronous tranmsmissions. */
262             TRT = fddi_timer_create ();

263     /* Set the TRT timer to expire in one TTRT */
264             TRT_SET (my_address, op_sim_time () + Fddi_T_Opr);

265     /* Initialize the Late_Ct variable which keeps track. */
266     /* of the number of TRT expirations. */
267             Late_Ct = 0;

268     /* initially the ring operates in nonrestricted mode */
269             restricted = 0;


270     /* Create an Interface Control Information structure */
271     /* to use when delivering received frames to the LLC. */
```

```
272             to_llc_ici_ptr = op_ici_create ("fddi_mac_ind");

273     /* The 'tk_registered' variable indicates if the station */
274     /* has registered its intent to use the token. */
275             tk_registered = 0;


276     /* Determine if the model is to make use of the token */
277     /* 'acceleration' mechanism.  If not, every passing of the */
278     /* token will be explicityly modeled, leading to large  */
279     /* number of events being scheduled when the ring is idle */
280     /* (i.e, no stations have data to send). */
281             op_ima_sim_attr_get (OPC_IMA_INTEGER, "accelerate_token",
282                &Fddi_Tk_Accelerate);

283     /* Obtain the synchronous bandwidth assigned */
284     /* to this station.  It is expressed as a */
285     /* percentage of TTRT, and then converted to seconds */
286             op_ima_obj_attr_get (my_objid, "sync bandwidth", &sync_pc);
287             sync_bandwidth = sync_pc * Fddi_T_Opr;


288     /* Only one station in the ring is selected to */
289     /* introduce the first token.  Test if this station is it. */
290     /* If so, set the 'spawn_token' flag. */
291             op_ima_sim_attr_get (OPC_IMA_INTEGER, "spawn station",
292                &spawn_station);
293             spawn_token = (spawn_station == my_address);


294     /* If the station is to spawn the token, create */
295     /* the packet which represents the token. */
296             if (spawn_token)
297                {
298                tk_pkptr = op_pk_create_fmt ("fddi_mac_tk");

299     /* assign its frame control field */
300                op_pk_nfd_set (tk_pkptr, "fc", FDDI_FC_TOKEN);

301     /* the first token issued is non-restricted */
302                op_pk_nfd_set (tk_pkptr, "class", FDDI_TK_NONRESTRICTED);

303     /* The transition will be made into the ISSU_TK */
304     /* state where the tk_usable variable is used. */
305     /* In case any data has been generated, prset */
306     /* this variable to one. */
307                tk_usable = 1;

308     /* When sending packets the variable accum_bandwidth is */
309     /* used as a scheduling base.  Init this value to zero. */
```

```
310   /* This statement is required in case this is the spawning */
311   /* station, and the next state entered is ISSUE_TK */
312            accum_bandwidth = 0.0;
313            }

314   /** state (INIT) exit executives **/
315         FSM_STATE_EXIT_FORCED (0, state0_exit_exec, "INIT")
316            {
317            }

318   /** state (INIT) transition processing **/
319         FSM_INIT_COND (spawn_token)
320         FSM_DFLT_COND
321         FSM_TEST_LOGIC ("INIT")

322         FSM_TRANSIT_SWITCH
323            {
324            FSM_CASE_TRANSIT (0, 2, state2_enter_exec, ;)
325            FSM_CASE_TRANSIT (1, 1, state1_enter_exec, ;)
326            }
327         /*---------------------------------------------------------*/

328   /** state (IDLE) enter executives **/
329         FSM_STATE_ENTER_UNFORCED (1, state1_enter_exec, "IDLE")
330            {
331            }

332   /** blocking after enter executives of unforced state. **/
333         FSM_EXIT (3,fddi_mac_mult)

334   /** state (IDLE) exit executives **/
335         FSM_STATE_EXIT_UNFORCED (1, state1_exit_exec, "IDLE")
336            {
337   /* Determine if a trace is activated for the FDDI model */
338            tk_trace_on = op_prg_odb_ltrace_active ("fddi_tk");

339   /* Trap packets arriving from physical layer so that their */
340   /* FC field can be extracted before evaluating conditions  */
341            if (op_intrpt_type () == OPC_INTRPT_STRM && op_intrpt_strm ()
342               != FDDI_LLC_STRM_IN)
343               {
344   /* Acquire the arriving packet. */
345               pkptr = op_pk_get (FDDI_PHY_STRM_IN);

346   /* Determine the type of packet by extracting */
347   /* the frame control field. */
348               op_pk_nfd_get (pkptr, "fc", &frame_control);

349   /* Physical layer arrival flag is set. */
350               phy_arrival = 1;
```

```
351                }
352             else{
353    /* The interrupt is not due to a physical layer arrival. */
354                phy_arrival = 0;

355    /* If the interrupt is a remote interrupt with specified code, it
356       signifies  */
357    /* the reinsertion of the token into the ring after an idle period. This
358       only */
359    /* occurs if the token acceleration mechanism is active.          */
360                if (op_intrpt_type () == OPC_INTRPT_REMOTE && op_intrpt_code
361                   () == FDDIC_TK_INJECT)
362                   {
363    /* create a new token */
364                   tk_pkptr = op_pk_create_fmt ("fddi_mac_tk");

365    /* assign its frame control field */
366                   op_pk_nfd_set (tk_pkptr, "fc", FDDI_FC_TOKEN);

367    /* the token is non-restricted */
368                   op_pk_nfd_set (tk_pkptr, "class", FDDI_TK_NONRESTRICTED);

369    /* insert it into the ring */
370                   op_pk_send (tk_pkptr, FDDI_PHY_STRM_OUT);
371                   }
372                }
373             }

374    /** state (IDLE) transition processing **/
375          FSM_INIT_COND (TK_RECEIVED)
376          FSM_TEST_COND (RC_FRAME)
377          FSM_TEST_COND (TRT_EXPIRE)
378          FSM_TEST_COND (FRAME_ARRIVAL)
379          FSM_DFLT_COND
380          FSM_TEST_LOGIC ("IDLE")

381          FSM_TRANSIT_SWITCH
382             {
383             FSM_CASE_TRANSIT (0, 3, state3_enter_exec, ;)
384             FSM_CASE_TRANSIT (1, 4, state4_enter_exec, ;)
385             FSM_CASE_TRANSIT (2, 7, state7_enter_exec, ;)
386             FSM_CASE_TRANSIT (3, 8, state8_enter_exec, ;)
387             FSM_CASE_TRANSIT (4, 1, state1_enter_exec, ;)
388             }
389    /*--------------------------------------------------------*/

390    /** state (ISSUE_TK) enter executives **/
391          FSM_STATE_ENTER_FORCED (2, state2_enter_exec, "ISSUE_TK")
392             {
393    /* If the token is sent without having been used, and the station */
```

```
394   /* has no data to send, then indicate this fact to the */
395   /* token acceleration mechanism which may have an */
396   /* oppurtunity to block the token. */
397            if (!tk_usable && op_q_stat (OPC_QSTAT_PKSIZE) == 0.0)
398               {
399   /* Note that if the token cannot be blocked, */
400   /* this procedure will forward the token physically. */
401               fddi_tk_indicate_no_data (tk_pkptr, my_address,
402                  accum_bandwidth);
403               }
404            else{
405               if (tk_trace_on == OPC_TRUE)
406                  {
407                  sprintf (str0, "Issuing token. accum_bw (%.9f), prop_del
408                     (%.9f)", accum_bandwidth, Fddi_Prop_Delay);
409                  op_prg_odb_print_major (str0, OPC_NIL);
410                  }

411   /* Send out the token packet using the accumulated */
412   /* consumed bandwidth as a scheduling base. */
413   /* In the case of the initial spawning of the token */
414   /* this will be zero; otherwise this variable will */
415   /* reflect the bandwidth consumed since the last capture */
416   /* of the usable token. Propagation delay is also accounted for. */
417               op_pk_send_delayed (tk_pkptr, FDDI_PHY_STRM_OUT,
418                  accum_bandwidth + Fddi_Prop_Delay);
419               }
420            }

421   /** state (ISSUE_TK) exit executives **/
422         FSM_STATE_EXIT_FORCED (2, state2_exit_exec, "ISSUE_TK")
423            {
424            }

425   /** state (ISSUE_TK) transition processing **/
426         FSM_TRANSIT_FORCE (1, state1_enter_exec, ;)
427   /*---------------------------------------------------------*/

428   /** state (RCV_TK) enter executives **/
429         FSM_STATE_ENTER_FORCED (3, state3_enter_exec, "RCV_TK")
430            {
431   /* The arriving packet, when received in the IDLE state */
432   /* is placed in the variable 'pkptr'.  Since it is now  */
433   /* known that it is a token, it can be placed in 'tk_pkptr. */
434            tk_pkptr = pkptr;

435   /* Load the token's class into the temporary variable 'tk_class.' */
436            op_pk_nfd_get (pkptr, "class", &tk_class);

437   /* If the token is restricted, determine for which station. */
```

```
438             if (tk_class == FDDI_TK_RESTRICTED)
439                {
440   /* Place the station address in the variable 'res_station' */
441   /* which may factor in to the determination of token usability. */
442                op_pk_nfd_get (tk_pkptr, "res_station", &res_station);
443                }

444   /* Determine if the token is usable: */
445   /* assume by default that it is not */
446   /* Subsequent conditions may override this. */
447             tk_usable = 0;
448   /* The token can only be usable if there are frames enqueued */
449   /* 27DEC93: the entire bank of subqueues must be checked,    */
450   /* starting at the highest priority (corresponding to */
451   /* synchronous traffic), and stopping when a packet is */
452   /* found. Then the loop is broken. -Nix    */
453             NUM_PRIOS = 9;
454             for (i = NUM_PRIOS - 1; i > -1; i--)
455             {
456             if (op_subq_stat (i, OPC_QSTAT_PKSIZE) > 0.0)
457                {
458   /* examine the attributes of the packet at the */
459   /* head of the queue. */
460   /* fddi_load_frame_attrs (&dest_addr, &svc_class, &pri_level); */
461                fddi_load_frame_attrs (dest_addr, &svc_class, &pri_level);

462   /* If synchronous data is queued, the token is */
463   /* necessarily usable, regardless of timing conditions. */
464                if (svc_class == FDDI_SVC_SYNC)
465                   {
466                   tk_usable = 1;
467                   break;
468                   }
469                else{
470   /* Otherwise, if asynchronous data is queued, it must */
471   /* meet several criteria for the token to be usable.  */

472   /* The token is only usable only if it is early. */
473                   if (Late_Ct == 0)
474                      {
475   /* The token's class must be nonrestricted, unless */
476   /* this station is involved in the restricted transfer. */
477                      if (tk_class == FDDI_TK_NONRESTRICTED || res_station
478                         == my_address || restricted)
479                         {
480   /* Test the frame's priority assignment against the current TRT */
481   /* This test uses the priority indirection table T_Pri */
482   /* so that only packets whose T_Pri [pri_level] exceeds */
483   /* the TRT can be transmitted.  In other words, by */
484   /* assigning lower values to T_Pri for a given priority */
```

```
485   /* level, packets of that level will be further restricted */
486   /* from using the ring bandwidth. */
487                         if (T_Pri [pri_level] >= fddi_timer_value (TRT))
488                            {
489                            tk_usable = 1;
490                               break;
491                            }
492                         }
493                      }
494                   }
495                } /* closes the "if (op_subq_stat (OPC_QSTAT_PKSIZE) > 0.0"
496                   statment */
497                }/* closes the "for" loop */

498   /* If the token is usable, timers must be readjusted. */
499             if (tk_usable)
500                {
501   /* The timer adjustment depends on whether the token is early or late.
502      */
503                if (Late_Ct == 0)
504                   {
505   /* Transfer the contents of TRT into THT. */
506                   fddi_timer_copy (TRT, THT);

507   /* Disable the THT timer. */
508                   fddi_timer_disable (THT);

509   /* Reset TRT to time the next rotation. */
510                   op_ev_cancel (TRT_handle);
511                   TRT_SET (my_address, op_sim_time () + Fddi_T_Opr);
512                   }
513                else{
514   /* If the token is late, set the THT to its expired */
515   /* value, and disable it.  This will prevent any */
516   /* asynchronous transmissions from occuring. */
517                   fddi_timer_set_value (THT, Fddi_T_Opr);
518                      fddi_timer_disable (THT);

519   /* clear the Late token counter (note that TRT is not modified, */
520   /* so that less than a full TTRT remains before TRT expires again. */
521                   Late_Ct = 0;
522                   }
523                }


524   /* If the token is not usable, different adjustments are made. */
525             else{
526   /* Again, the adjustments depend on the lateness of the token */
527                if (Late_Ct == 0)
528                   {
```

```
529   /* If the token is not late, the TRT is reset to time the next rotation.
530      */
531                   op_ev_cancel (TRT_handle);
532                   TRT_SET (my_address, op_sim_time () + Fddi_T_Opr);
533                   }
534               else{
535   /* clear the Late token counter (note that TRT is not modified, */
536   /* so that less than a full TTRT remains before TRT expires again. */
537                   Late_Ct = 0;
538                   }

539   /* also, account for the time needed by the token  */
540   /* to traverse the station, since it is about to be sent.  */
541   /* Note: station latency is not inclusive of token  */
542   /* transmission time, but only of the time required to  */
543   /* process and repeat the token's symbols.       */
544               accum_bandwidth = Fddi_St_Latency;
545               }
546           }

547   /** state (RCV_TK) exit executives **/
548         FSM_STATE_EXIT_FORCED (3, state3_exit_exec, "RCV_TK")
549            {
550            }

551   /** state (RCV_TK) transition processing **/
552         FSM_INIT_COND (tk_usable)
553         FSM_DFLT_COND
554         FSM_TEST_LOGIC ("RCV_TK")

555         FSM_TRANSIT_SWITCH
556            {
557            FSM_CASE_TRANSIT (0, 9, state9_enter_exec, ;)
558            FSM_CASE_TRANSIT (1, 2, state2_enter_exec, ;)
559            }
560   /*--------------------------------------------------------*/

561   /** state (FR_RCV) enter executives **/
562         FSM_STATE_ENTER_FORCED (4, state4_enter_exec, "FR_RCV")
563            {
564   /* A frame has been received from the physical layer. Note that */
565   /* at this time, only the leading edge of the frame has arrived. */

566   /* Extract the frame's source address (this will be used to */
567   /* determine whether or not to strip the frame from the ring). */
568            op_pk_nfd_get (pkptr, "src_addr", &src_addr);
569            }

570   /** state (FR_RCV) exit executives **/
571         FSM_STATE_EXIT_FORCED (4, state4_exit_exec, "FR_RCV")
```

```
572             {
573             }

574     /** state (FR_RCV) transition processing **/
575           FSM_INIT_COND (STRIP)
576           FSM_DFLT_COND
577           FSM_TEST_LOGIC ("FR_RCV")

578           FSM_TRANSIT_SWITCH
579              {
580              FSM_CASE_TRANSIT (0, 5, state5_enter_exec, ;)
581              FSM_CASE_TRANSIT (1, 6, state6_enter_exec, ;)
582              }
583        /*---------------------------------------------------------*/

584     /** state (FR_STRIP) enter executives **/
585           FSM_STATE_ENTER_FORCED (5, state5_enter_exec, "FR_STRIP")
586              {
587     /* Destroy the frame which has now circulated the entire ring. */
588              op_pk_destroy (pkptr);
589              }

590     /** state (FR_STRIP) exit executives **/
591           FSM_STATE_EXIT_FORCED (5, state5_exit_exec, "FR_STRIP")
592              {
593              }

594     /** state (FR_STRIP) transition processing **/
595           FSM_TRANSIT_FORCE (1, state1_enter_exec, ;)
596        /*---------------------------------------------------------*/

597     /** state (FR_REPEAT) enter executives **/
598           FSM_STATE_ENTER_FORCED (6, state6_enter_exec, "FR_REPEAT")
599              {
600     /* Extract the destination address of the frame. */
601     /* 20FEB94: use a pointer to the array dest_addr, */
602     /* since referring to dest_addr directly produces */
603     /* unexpected results. -Nix    */
604              op_pk_nfd_get (pkptr, "dest_addr", &da_ptr);

605     /*********************************************************/
606     /* printf("*da_ptr: %d; da_ptr: %d; &da_ptr: %d\n", *da_ptr, da_ptr,
607        &da_ptr); */
608     /*********************************************************/
609
610              for (i = 0; i < NUM_STATIONS+1; i++);
611                 dest_addr[i] = da_ptr[i];
612
613     /*********************************************************/
614     /* 02MAR94: print out the address, and the contents.   */
```

```
615   /* for (i = 0; i < NUM_STATIONS+1; i++)          */
616   /* {                                  */
617   /* printf("1.FR_REPEAT:element: %d, address: %X/%d, content: %d\n",
618      */
619   /* i, &(dest_addr[i]), &(dest_addr[i]), dest_addr[i]);*/
620   /* }                                  */
621   /*****************************************************/

622   /* 08FEB94: re-initialize counters. -Nix  */
623            for_me = 0;
624            count_addees = 0;

625   /* 08FEB94: inspect the address field; interested in */
626   /* whether this packet is sent here only, or here and    */
627   /* to others, or to others only.Note that a real packet  */
628   /* would carry all the addresses; the simulation refers  */
629   /* to memory locations. -Nix      */
630            for (i = 1; i < NUM_STATIONS+1; i++)
631            {
632            if (dest_addr[i] == 1)
633               count_addees += 1;
634            }

635   /* If the frame is for this station, make a copy */
636   /* of the frame's data field and forward it to */
637   /* the higher layer. */
638   /* if (dest_addr == my_address) */
639   /* 08FEB94: if this packet is addressed only to this   */
640   /* station, make a copy of the frame's data field and  */
641   /* and forward it to the higher layer. -Nix */
642   /* (a) If the packet is addressed to me only... */
643   /* (note offset applied)          */
644            if (dest_addr[my_address+1] == 1 && count_addees == 1)
645               {

646   /* ****************************** */
647   /* printf("Here is Case 1.\n");   */
648   /* ****************************** */

649   /* record total size of the frame (including data) */
650               pk_len = op_pk_total_size_get (pkptr);

651   /* decapsulate the data contents of the frame */
652   /* 29JAN94: a new field, "pri", has been added to  */
653   /* the fddi_llc_fr packet format in the Parameters */
654   /* Editor, so that output statistics can be */
655   /* generated by class and priority. -Nix */
656               op_pk_nfd_get (pkptr, "info", &data_pkptr);
657               op_pk_nfd_get (pkptr, "pri", &pri_level);
```

```
658   /* The source and destination address are placed in the */
659   /* LLC's ICI before delivering the frame's contents. */
660               op_ici_attr_set (to_llc_ici_ptr, "src_addr", src_addr);
661               op_ici_attr_set (to_llc_ici_ptr, "dest_addr", da_ptr);
662               op_ici_install (to_llc_ici_ptr);

663   /* 18FEB94: print out the address, and the contents. */
664            for (i = 0; i < NUM_STATIONS+1; i++)
665            {
666               dest_addr[i] = da_ptr[i];

667   /*********************************************************************/
668   /* printf("2.FR_REPEAT:element: %d, address: %X/%d, content: %d\n",  */
669   /* i, &(dest_addr[i]), &(dest_addr[i]), dest_addr[i]);   */
670   /*********************************************************************/
671            }

672   /* Because, as noted in the FR_RCV state, only the */
673   /* frame's leading edge has arrived at this time, the */
674   /* complete frame can only be delivered to the higher */
675   /* layer after the frame's transmission delay has elapsed. */
676   /* (since decapsulation of the frame data contents has occured, */
677   /* the original MAC frame length is used to calculate delay) */
678               tx_time = (double) pk_len / FDDI_TX_RATE;
679               op_pk_send_delayed (data_pkptr, FDDI_LLC_STRM_OUT, tx_time);

680   /* Note that the standard specifies that the original */
681   /* frame should be passed along until the originating station */
682   /* receives it, at which point it is stripped from the ring. */
683   /* However, in the simulation model, there is no interest */
684   /* in letting the frame continue past its destination unless */
685   /* group addresses are used, so that the same frame could be */
686   /* destined for several stations.  Here the frame is stripped */
687   /* for efficiency as it reaches the destination; if the model */
688   /* is modified to include group addresses, this should be change. */
689   /* so that the frame is copied and the original repeated. */
690   /* Logic is already present for stripping the frame at the origin. */
691               op_pk_destroy (pkptr);
692               }

693   /* 08FEB94: (b)...or if this packet is not for me at all... -Nix  */
694            else if (dest_addr[my_address+1] == 0)
695               {

696   /* ****************************** */
697   /*  printf("Here is Case 2.\n");  */
698   /* ****************************** */

699   /* Repeat the original frame on the ring and account for */
700   /* the latency through the station and the propagation delay */
```

```
701   /* for a single hop. */
702   /* (Only the originating station can strip the frame). */
703                op_pk_send_delayed (pkptr, FDDI_PHY_STRM_OUT,
704                   Fddi_St_Latency + Fddi_Prop_Delay);
705                }

706   /* 08FEB94: (c)...or if this packet is for me and for others, will  */
707   /* need to send the contents to the SINK, then re-encapsulate the   */
708   /* packet for further transmission. Much of this code is */
709   /* duplicated from the above.  -Nix        */

710             else if (dest_addr[my_address+1] == 1 && count_addees > 1)
711                {

712   /* ************************** */
713   /* printf("Here is Case 3.\n"); */
714   /* ************************** */

715   /* record total size of the frame (including data) */
716                pk_len = op_pk_total_size_get (pkptr);

717   /* decapsulate the data contents of the frame */
718                op_pk_nfd_get (pkptr, "info", &data_pkptr);
719                op_pk_nfd_get (pkptr, "pri", &pri_level);

720   /* ********************************************************************* */
721   /* Print out the address of the "info" field information */
722   /* printf("Case 3: 'info' is located at address %X\n", &data_pkptr); */
723   /* ********************************************************************* */

724   /* 08MAR94: copy the "info" address into a local variable, so that   */
725   /* it may be held for re-installation. -Nix      */
726             info_ptr = op_pk_copy(data_pkptr);


727   /* The source and destination address are placed in the */
728   /* LLC's ICI before delivering the frame's contents. */
729                op_ici_attr_set (to_llc_ici_ptr, "src_addr", src_addr);
730                op_ici_attr_set (to_llc_ici_ptr, "dest_addr", dest_addr);
731                op_ici_install (to_llc_ici_ptr);

732   /* Because, as noted in the FR_RCV state, only the */
733   /* frame's leading edge has arrived at this time, the */
734   /* complete frame can only be delivered to the higher */
735   /* layer after the frame's transmission delay has elapsed. */
736   /* (since decapsulation of the frame data contents has occured, */
737   /* the original MAC frame length is used to calculate delay) */
738                tx_time = (double) pk_len / FDDI_TX_RATE;
739                op_pk_send_delayed (data_pkptr, FDDI_LLC_STRM_OUT, tx_time);
```

```
740   /* 08FEB94: remove this station from the dest_addr array, reassemble */
741   /* the packet, and send the packet on its way.  -Nix */
742            dest_addr[my_address+1] = 0;
743            op_pk_nfd_set(pkptr,"src_addr", src_addr);
744            op_pk_nfd_set(pkptr,"dest_addr", dest_addr);
745            op_pk_nfd_set(pkptr,"pri", pri_level);
746            op_pk_nfd_set(pkptr,"info", info_ptr);
747            op_pk_send_delayed (pkptr, FDDI_PHY_STRM_OUT,
748                  Fddi_St_Latency + Fddi_Prop_Delay);

749               }

750            }

751   /** state (FR_REPEAT) exit executives **/
752         FSM_STATE_EXIT_FORCED (6, state6_exit_exec, "FR_REPEAT")
753            {
754            }

755   /** state (FR_REPEAT) transition processing **/
756         FSM_TRANSIT_FORCE (1, state1_enter_exec, ;)
757   /*---------------------------------------------------------*/

758   /** state (TRT_EXP) enter executives **/
759         FSM_STATE_ENTER_FORCED (7, state7_enter_exec, "TRT_EXP")
760            {
761   /* The timer is reset and allowed to continue running. */
762            TRT_SET (my_address, op_sim_time () + Fddi_T_Opr);

763   /* The late token counter is incremented.  This will  */
764   /* prevent this station from making any asynchronous  */
765   /* transmissions when it next captures the token.  */
766            Late_Ct++;

767            }

768   /** state (TRT_EXP) exit executives **/
769         FSM_STATE_EXIT_FORCED (7, state7_exit_exec, "TRT_EXP")
770            {
771            }

772   /** state (TRT_EXP) transition processing **/
773         FSM_TRANSIT_FORCE (1, state1_enter_exec, ;)
774   /*---------------------------------------------------------*/

775   /** state (ENCAP) enter executives **/
776         FSM_STATE_ENTER_FORCED (8, state8_enter_exec, "ENCAP")
777            {
778   /* A frame has arrived from a higher layer; place it in 'pdu_ptr'. */
779            pdu_ptr = op_pk_get (op_intrpt_strm ());
```

```
780   /* Also get the interface control information */
781   /* associated with the new frame. */
782            ici_ptr = op_intrpt_ici ();
783            if (ici_ptr == OPC_NIL)
784               {
785               sprintf (error_string, "Simulation aborted; error in object
786                  (%d)", op_id_self ());
787               op_sim_end (error_string, "fddi_mac: required ICI not
788                  received", " ", " ");
789               }

790   /* Extract the requested service class */
791   /* (e.g, synchronous or asynchronous). */
792            if (op_ici_attr_exists (ici_ptr, "svc_class"))
793               op_ici_attr_get (ici_ptr, "svc_class", &svc_class);
794            else svc_class = FDDI_SVC_ASYNC;

795   /*****************************************************************/
796   /* for (i=0; i<NUM_STATIONS+1; i++)           */
797   /*  printf("ENCAP a.Field:%d, Address(dec/hex):%d/%X, Contents:%d\n",
798      */
799   /* i, &(dest_addr[i]), &(dest_addr[i]), dest_addr[i]); */
800   /*****************************************************************/



801   /* Extract the destination address. */
802   /* 20FEB94: use a pointer to the array, since the   */
803   /* use of dest_array as its own pointer causes */
804   /* unexpected results. -Nix */
805            op_ici_attr_get (ici_ptr, "dest_addr", &da_ptr);

806   /*****************************************************************/
807   /* printf("&da_ptr: %d/%X; da_ptr: %d\n\n", &da_ptr, &da_ptr, da_ptr);
808      */
809   /* for (i=0; i<NUM_STATIONS+1; i++)                */
810   /* printf("%d:  &da_ptr: %d/%X; da_ptr: %d\n\n", i, &da_ptr, &da_ptr,
811      da_ptr);  */
812   /*****************************************************************/

813            for (i=0; i<NUM_STATIONS+1; i++)
814            {
815            dest_addr[i] = da_ptr[i];
816            }

817   /*****************************************************************/
818   /*  for (i=0; i<NUM_STATIONS+1; i++)             */
819   /*  printf("ENCAP b.Field:%d, Address(dec/hex):%d/%X, Contents:%d\n",
820      */
821   /*  i, &(dest_addr[i]), &(dest_addr[i]), dest_addr[i]); */
```

```
822  /*****************************************************************/

823  /* If the frame is asynchronous, the priority and */
824  /* requested token class parameter may be specified. */
825          if (svc_class == FDDI_SVC_ASYNC)
826             {
827  /* Extract the requested priority level. */
828             if (op_ici_attr_exists (ici_ptr, "pri"))
829                op_ici_attr_get (ici_ptr, "pri", &req_pri);
830             else req_pri = 0;

831  /* Extract the token class (restrictred or non-restricted). */
832             if (op_ici_attr_exists (ici_ptr, "tk_class"))
833                op_ici_attr_get (ici_ptr, "tk_class", &req_tk_class);
834             else req_tk_class = FDDI_TK_NONRESTRICTED;
835             }


836  /* Compose a mac frame from all these elements. */
837          mac_frame_ptr = op_pk_create_fmt ("fddi_mac_fr");
838          op_pk_nfd_set (mac_frame_ptr, "svc_class", svc_class);
839          op_pk_nfd_set (mac_frame_ptr, "dest_addr", dest_addr);
840          op_pk_nfd_set (mac_frame_ptr, "src_addr", my_address);
841          op_pk_nfd_set (mac_frame_ptr, "info", pdu_ptr);

842          if (svc_class == FDDI_SVC_ASYNC)
843             {
844             op_pk_nfd_set (mac_frame_ptr, "tk_class", req_tk_class);
845             op_pk_nfd_set (mac_frame_ptr, "pri", req_pri);
846             }

847  /* 04JAN94: if the frame is synchronous, assign it a separate */
848  /* priority so that it may be assigned its own subqueue, and  */
849  /* thereby be assigned its own probe for monitoring. -Nix */
850          if (svc_class == FDDI_SVC_SYNC)
851          {
852          op_pk_nfd_set (mac_frame_ptr, "pri", 8);
853          }

854  /* Assign the frame control field, which in the model */
855  /* is used to distinguish between tokens and ordinary */
856  /* frames on the ring. */
857          op_pk_nfd_set (mac_frame_ptr, "fc", FDDI_FC_FRAME);

858  /* Enqueue the frame at the tail of the queue. */
859  /* 27DEC93: at the tail of the prioritized queue.    */
860  /* 04JAN94: must distinguish between synch & asynch. */
861          if (svc_class == FDDI_SVC_ASYNC)
862          {
863          op_subq_pk_insert (req_pri, mac_frame_ptr, OPC_QPOS_TAIL);
```

```
            }
            if (svc_class == FDDI_SVC_SYNC)
            {
            op_subq_pk_insert (8, mac_frame_ptr, OPC_QPOS_TAIL);
            }

/* if this station has not yet registered its intent to */
/* use the token, it may do so now since it has data to send */
            if (!tk_registered)
               {
               fddi_tk_register ();
               tk_registered = 1;
               }
            }

/** state (ENCAP) exit executives **/
      FSM_STATE_EXIT_FORCED (8, state8_exit_exec, "ENCAP")
            {
            }

/** state (ENCAP) transition processing **/
      FSM_TRANSIT_FORCE (1, state1_enter_exec, ;)
/*---------------------------------------------------------*/

/** state (TX_DATA) enter executives **/
      FSM_STATE_ENTER_FORCED (9, state9_enter_exec, "TX_DATA")
            {
/* In this state, frames are transmitted until the */
/* token is no longer usable.  Frames are taken from */
/* the single input queue in FIFO order. */

/* Reset the accumulator used to keep track of bandwidth */
/* consumed by the transmissions.  Because all the transmissions */
/* are scheduled to happen at the appropriate times, but */
/* these schedulings occur instantly, this accumulator serves */
/* as the scheduling base for the transmissions. */
/* In other words, each successively transmitted frame */
/* is delayed relative to the previous one by the time which */
/* the latter took to send.  At the end of transmission (e.g, */
/* when the token is no longer usable), this accumulator */
/* serves to delay the forwarding of the token. */
            accum_bandwidth = 0.0;

/* Note that, because all tranmsmissions are */
/* scheduled, the value of the THT timer will not progress */
/* between shcedulings (these all happen in zero time), and so  */
/* the variable 'tht_value' is used to emulate the timer's progress. */
            tht_value = fddi_timer_value (THT);

/* Reset an accumulator which reflects the consumed */
```

```
907   /* synchronous bandwidth. */
908           accum_sync = 0.0;

909   /* Reset counters for transmitted frames and bits. */
910           num_frames_sent = 0;
911           num_bits_sent = 0;

912   /* The transmission sequence must end if the input queue */
913   /* becomes exhausted.  Other termination conditions are */
914   /* embedded in the loop. */
915   /* 27DEC93: modify the loop to accomodate subqueue structure. */
916   /* A "for" loop is imposed over the original "while" loop. */
917   /* First, reset the break marker, "punt". -Nix */
918           punt = 0;
919           for (i = NUM_PRIOS - 1; i > -1; i--)
920           {
921           while (op_subq_stat (i,OPC_QSTAT_PKSIZE) > 0.0)
922              {
923   /* Remove the next frame for transmission. */
924              pkptr = op_subq_pk_remove (i, OPC_QPOS_HEAD);

925   /* Obtain the frame's service class. */
926              op_pk_nfd_get (pkptr, "svc_class", &svc_class);

927   /* Synchronous and asynchronous frames are treated differently. */
928              if (svc_class == FDDI_SVC_SYNC)
929                 {
930   /* Obtain the frame's length, and compute */
931   /* the time required to transmit it. */
932                 pk_len = op_pk_total_size_get (pkptr);
933                 tx_time = (double) pk_len / FDDI_TX_RATE;

934   /* Check if synchronous bandwidth allocation for this */
935   /* station would be exceeded if the transmission were to occur. */
936                 if (accum_sync + tx_time > sync_bandwidth)
937                    {
938   /* The frame could not be sent without exceeding */
939   /* the allocated synchronous bandwidth, */
940   /* so it is replaced on the queue. */
941   /* 27DEC93: in this case, i is the highest priority, */
942   /* which is reserved for synchronous traffic. -Nix   */
943                    op_subq_pk_insert (i, pkptr, OPC_QPOS_HEAD);

944   /* Exit the transmission loop since the frame  */
945   /* transmission request cannot be honored. */
946                    punt = 1;
947                    break;
948                    }
949                 else{
950   /* Send the frame into the ring after other frames have completed. */
```

```
951   /* Also, account for its proagation delay; because the token propagation
952      */
953   /* delay and the frame propagation delay must be consistent, and the */
954   /* token propagation delay is specified as a ring parameter (i.e,
955      stations */
956   /* are assumed to be equally spaced), the ring is intended to run with
957      */
958   /* the "delay" attributes of point-to-point links set at zero. */
959                     op_pk_send_delayed (pkptr, FDDI_PHY_STRM_OUT,
960                        accum_bandwidth + Fddi_Prop_Delay);

961   /* increase the consumed bandwidth to reflect this */
962   /* transmission.  Also increase synchronous consumption. */
963                     accum_bandwidth += tx_time;
964                     accum_sync += tx_time;

965   /* Increase counters for transmitted bits and frames. */
966                     num_frames_sent++;
967                     num_bits_sent += pk_len;
968                     }
969                  }
970               else{
971   /* The request enqueued at the head of the queue is */
972   /* asynchronous. It may only be honored if the THT timer */
973   /* has not expired. */
974                  if (tht_value >= Fddi_T_Opr)
975                     {
976   /* replace the packet on the queue and exit the transmission loop. */
977                     op_subq_pk_insert (i, pkptr, OPC_QPOS_HEAD);
978                     punt = 1;
979                     break;
980                     }
981                  else{
982   /* Obtain the priority assignment of the frame. */
983                     op_pk_nfd_get (pkptr, "pri", &pri_level);

984   /* If the packet's assigned priority level */
985   /* is too low for it to be serviced, then exit the loop */
986   /* after replacing the packet in the queue. */

987   /* ************************************************ */
988   /* 08MAR94: print the values to be compared. -Nix */
989   /* printf("1. TX_DATA: T_Pri[%d] < tht_value ?\n", i); */
990   /* printf("%d < %d ?\n", T_Pri [i], tht_value); */
991   /* ************************************************ */

992                     if (T_Pri [pri_level] < tht_value)
993                        {
994                        op_subq_pk_insert (i, pkptr, OPC_QPOS_HEAD);
995                        punt = 1;
```

```
                        break;
                        }

/* Obtain the frame's length, and compute the time */
/* which would be required to transmit it. */
                     pk_len = op_pk_total_size_get (pkptr);
                     tx_time = (double) pk_len / FDDI_TX_RATE;

/* Determine the requested token class to be */
/* released after this frame is transmitted. */
                     op_pk_nfd_get (pkptr, "tk_class", &tk_class);

/* If the station is in restricted mode, then it may */
/* exit this mode if the class is now nonrestricted */
/* or if the restricted peer is not the addressee. */
                     if (restricted)
                        {
/* Determine the destination address for the new packet. */
                        op_pk_nfd_get (pkptr, "dest_addr", dest_addr);

/* if (tk_class == FDDI_TK_NONRESTRICTED ||   */
/*    res_peer != dest_addr) */

                     if (tk_class == FDDI_TK_NONRESTRICTED ||
                        dest_addr[res_peer] != 1)
                           {
/* Exit restricted mode */
                           restricted = 0;

/* Modify the token to reflect the mode change. */
                           op_pk_nfd_set (tk_pkptr, "class",
                              FDDI_TK_NONRESTRICTED);
                           }
                        }
                     else{
/* Determine the class of the current captured token. */
                        op_pk_nfd_get (tk_pkptr, "class",
&current_tk_class);

/* When not in restricted mode, this mode may be entered */
/* if the passed packet has the appropriate token class requested, */
/* and the token is not already restricted. */
                        if (tk_class == FDDI_TK_RESTRICTED &&
                           current_tk_class != FDDI_TK_RESTRICTED)
                           {

/* Enter restricted mode. */
                           restricted = 1;

/* Store the address of the resticted peer station. */
```

```
1038    /* op_pk_nfd_get (pkptr, "dest_addr", &res_peer); */
1039                            op_pk_nfd_get (pkptr, "dest_addr",
1040                               &dest_addr[res_peer]);

1041    /* Modify the token to reflect the mode change. */
1042                            op_pk_nfd_set (tk_pkptr, "class",
1043                               FDDI_TK_RESTRICTED);
1044                            op_pk_nfd_set (tk_pkptr, "res_station",
1045    res_peer);
1046                            }
1047                         }

1048    /* Send the frame once previous transmissions have completed. */
1049    /* Account for propagation delay as well. */
1050                      op_pk_send_delayed (pkptr, FDDI_PHY_STRM_OUT,
1051                         accum_bandwidth + Fddi_Prop_Delay);


1052    /* Increment THT emulation variable, and consumed bandwidth accumulator.
1053       */
1054    /* 08MAR94: note that tht_value is incrementing, not decrementing. -Nix
1055       */
1056                      tht_value += tx_time;
1057                      accum_bandwidth += tx_time;

1058    /* ************************************************** */
1059    /* 08MAR94: print the Token Holding Time value. -Nix */
1060    /* printf("2. TX_DATA: tht_value is %d\n", tht_value);   */
1061    /* ************************************************** */

1062    /* Increase counters for transmitted bits and frames. */
1063                      num_frames_sent++;
1064                      num_bits_sent += pk_len;
1065                      }
1066                   }
1067                }/* closes the 'while' loop */
1068              if (punt == 1)    /* If the 'while' loop was broken, */
1069              {
1070              punt = 0;/* then reset the 'break' marker, */
1071              break;/* and break out of the 'for' loop too. */
1072              }
1073              }/* closes the 'for' loop.  */

1074    /* Since the token is about to be sent, its transmission time  */
1075    /* must be reflected in the accumulated bandwidth. This is not */
1076    /* done in the ISSUE_TK state because when the token is merely */
1077    /* repeated, full transmission delay is not required, only  */
1078    /* a small delay for repeating.                    */
1079              accum_bandwidth += FDDIC_TOKEN_TX_TIME;
```

```
1080   /* If the station has no more data to send (synchronous or */
1081   /* asynchronous), it should indicate this to the token acceleration */
1082   /* mechanism by deregistering its interest in the token. */
1083   /* 27DEC94: the original code must be modified to include a check   */
1084   /* of subqueues. -Nix */
1085            q_check = 1;
1086            for (i = NUM_PRIOS - 1; i < -1; i--)
1087            {
1088            if (op_subq_stat (i, OPC_QSTAT_PKSIZE) == 0.0)
1089               {
1090               q_check = 0;
1091               }
1092            else {
1093               q_check = 1;
1094               break;
1095               }
1096            }

1097            if (tk_registered && q_check == 0)
1098               {
1099               tk_registered = 0;
1100               fddi_tk_deregister ();
1101               }
1102            }

1103   /** state (TX_DATA) exit executives **/
1104         FSM_STATE_EXIT_FORCED (9, state9_exit_exec, "TX_DATA")
1105            {
1106            }

1107   /** state (TX_DATA) transition processing **/
1108         FSM_TRANSIT_FORCE (2, state2_enter_exec, ;)
1109   /*---------------------------------------------------------*/

1110   /** state (CLAIM) enter executives **/
1111         FSM_STATE_ENTER_UNFORCED (10, state10_enter_exec, "CLAIM")
1112            {
1113   /* Obtain this station's object id which is used */
1114   /* to access the station's attribute assignments. */
1115            my_objid = op_id_self ();

1116   /* Using the object id, obtain the ring id. */
1117   /* The ring id is used by macros defined in the */
1118   /* header block to obtain "ring-global" values, */
1119   /* values shared by all stations on a ring. */
1120            op_ima_obj_attr_get (my_objid, "ring_id", &ring_id);

1121   /* Initialize global variable values. */
1122            Fddi_Tk_Blocked = 0;
1123            Fddi_Num_Stations = 0;
```

```
1124            Fddi_Num_Registered = 0;

1125    /* Using the object id, obtain the value of 'T_Req', */
1126    /* the value of TTRT requested by this station. */
1127            op_ima_obj_attr_get (my_objid, "T_Req", &T_Req);

1128    /* The lowest value of T_Req becomes T_Opr for the ring as a whole. */
1129            if (T_Req < Fddi_T_Opr || Fddi_Claim_Start)
1130               {
1131    /* The T_Req for this station is lower than any other to date */
1132    /* so it is installed in the T_Opr variable. */
1133               Fddi_T_Opr = T_Req;

1134    /* The flag indicating that the claim process is just */
1135    /* beginning may now be cleared. */
1136               Fddi_Claim_Start = 0;
1137               }

1138    /* Request a self interrupt from the Simulation Kernel at the current */
1139    /* time so that after all stations have executed their claim states,  */
1140    /* they can proceed with initializations.  This is necessasary */
1141    /* because some initializations are based in the value of T_Opr */
1142    /* and it must therefore be known that all stations have settled */
1143    /* on a final value. */
1144            op_intrpt_schedule_self (op_sim_time (), 0);
1145            }

1146    /** blocking after enter executives of unforced state. **/
1147          FSM_EXIT (21,fddi_mac_mult)

1148    /** state (CLAIM) exit executives **/
1149          FSM_STATE_EXIT_UNFORCED (10, state10_exit_exec, "CLAIM")
1150             {
1151             }

1152    /** state (CLAIM) transition processing **/
1153          FSM_TRANSIT_FORCE (0, state0_enter_exec, ;)
1154    /*---------------------------------------------------------*/
1155          }
1156       FSM_EXIT (10,fddi_mac_mult)
1157       }

1158    void
1159    fddi_mac_mult_svar (prs_ptr,var_name,var_p_ptr)
1160       fddi_mac_mult_state     *prs_ptr;
1161       char                       *var_name, **var_p_ptr;
1162       {

1163       FIN (fddi_mac_mult_svar (prs_ptr))
```

```
1164      *var_p_ptr = VOS_NIL;
1165      if (Vos_String_Equal ("ring_id" , var_name))
1166         *var_p_ptr = (char *) (&prs_ptr->sv_ring_id);
1167      if (Vos_String_Equal ("THT" , var_name))
1168         *var_p_ptr = (char *) (&prs_ptr->sv_THT);
1169      if (Vos_String_Equal ("T_Req" , var_name))
1170         *var_p_ptr = (char *) (&prs_ptr->sv_T_Req);
1171      if (Vos_String_Equal ("T_Pri" , var_name))
1172         *var_p_ptr = (char *) (prs_ptr->sv_T_Pri);
1173      if (Vos_String_Equal ("my_objid" , var_name))
1174         *var_p_ptr = (char *) (&prs_ptr->sv_my_objid);
1175      if (Vos_String_Equal ("spawn_token" , var_name))
1176         *var_p_ptr = (char *) (&prs_ptr->sv_spawn_token);
1177      if (Vos_String_Equal ("my_address" , var_name))
1178         *var_p_ptr = (char *) (&prs_ptr->sv_my_address);
1179      if (Vos_String_Equal ("tk_pkptr" , var_name))
1180         *var_p_ptr = (char *) (&prs_ptr->sv_tk_pkptr);
1181      if (Vos_String_Equal ("sync_bandwidth" , var_name))
1182         *var_p_ptr = (char *) (&prs_ptr->sv_sync_bandwidth);
1183      if (Vos_String_Equal ("sync_pc" , var_name))
1184         *var_p_ptr = (char *) (&prs_ptr->sv_sync_pc);
1185      if (Vos_String_Equal ("restricted" , var_name))
1186         *var_p_ptr = (char *) (&prs_ptr->sv_restricted);
1187      if (Vos_String_Equal ("res_peer" , var_name))
1188         *var_p_ptr = (char *) (&prs_ptr->sv_res_peer);
1189      if (Vos_String_Equal ("tk_registered" , var_name))
1190         *var_p_ptr = (char *) (&prs_ptr->sv_tk_registered);
1191      if (Vos_String_Equal ("to_llc_ici_ptr" , var_name))
1192         *var_p_ptr = (char *) (&prs_ptr->sv_to_llc_ici_ptr);
1193      if (Vos_String_Equal ("tk_trace_on" , var_name))
1194         *var_p_ptr = (char *) (&prs_ptr->sv_tk_trace_on);

1195      FOUT;
1196      }

1197   void
1198   fddi_mac_mult_diag ()
1199      {
1200   /* Packets and ICI's */
1201      Packet*        mac_frame_ptr;
1202      Packet*        pdu_ptr;
1203      Packet*        pkptr;
1204      Packet*        data_pkptr;
1205      Ici*           ici_ptr;

1206   /* Packet Fields and Attributes */
1207      int            req_pri, svc_class, req_tk_class;
1208      int            frame_control, src_addr;
1209      int            pk_len, pri_level;
1210      static
```

```
1211      int             *da_ptr, dest_addr[];

1212    /* Token - Related */
1213      int             tk_usable, res_station, tk_class;
1214      int             current_tk_class;
1215      double          accum_sync;

1216    /* Timer - Related */
1217      double          tx_time, timer_remaining, accum_bandwidth;
1218      double          tht_value;

1219    /* Miscellaneous */
1220      int             i;
1221      int             spawn_station, phy_arrival;
1222      char         error_string [512];
1223      int             num_frames_sent, num_bits_sent;

1224    /* 26DEC93: loop management variables, used in RCV_TK  */
1225    /* and ENCAP states. -Nix */
1226      int   NUM_PRIOS;
1227      int   punt;
1228      int   q_check;

1229    /* 08FEB94: case management variables, used in FR_REPEAT. -Nix */
1230      int   for_me;
1231      int   count_addees;

1232    /* 08MAR94: "field holding" variables, used in FR_REPEAT. -Nix */
1233      Packet*  info_ptr;

1234      FIN (fddi_mac_mult_diag ())

1235    /* Print out values of timers, and late token counter.  */
1236    /* Also print out data about restricted mode.    */
1237    /* (This code may be executed by the simulation debugger  */
1238    /* by invoking the command 'modprint').       */

1239      sprintf (str0, "Timers (count upwards):  TRT (%.9g), THT (%.9g)",
1240         fddi_timer_value (TRT), fddi_timer_value (THT));
1241      sprintf (str1, "Late_ct (%d)", Late_Ct);
1242      op_prg_odb_print_major (str0, str1, OPC_NIL);
1243      if (restricted)
1244         sprintf (str0, "token is in restricted dialog with (%d)\n",
1245            res_peer);
1246      else sprintf (str0, "token is unrestricted\n");
1247      op_prg_odb_print_major (str0, OPC_NIL);
1248      FOUT;
1249      }

1250    void
```

```
1251  fddi_mac_mult_terminate ()
1252     {
1253  /* Packets and ICI's */
1254     Packet*        mac_frame_ptr;
1255     Packet*        pdu_ptr;
1256     Packet*        pkptr;
1257     Packet*        data_pkptr;
1258     Ici*           ici_ptr;

1259  /* Packet Fields and Attributes */
1260     int            req_pri, svc_class, req_tk_class;
1261     int            frame_control, src_addr;
1262     int            pk_len, pri_level;
1263     static
1264     int            *da_ptr, dest_addr[];

1265  /* Token - Related */
1266     int            tk_usable, res_station, tk_class;
1267     int            current_tk_class;
1268     double         accum_sync;

1269  /* Timer - Related */
1270     double         tx_time, timer_remaining, accum_bandwidth;
1271     double         tht_value;

1272  /* Miscellaneous */
1273     int            i;
1274     int            spawn_station, phy_arrival;
1275     char        error_string [512];
1276     int            num_frames_sent, num_bits_sent;

1277  /* 26DEC93: loop management variables, used in RCV_TK  */
1278  /* and ENCAP states. -Nix         */
1279     int   NUM_PRIOS;
1280     int   punt;
1281     int   q_check;

1282  /* 08FEB94: case management variables, used in FR_REPEAT. -Nix */
1283     int   for_me;
1284     int   count_addees;

1285  /* 08MAR94: "field holding" variables, used in FR_REPEAT. -Nix */
1286     Packet*  info_ptr;

1287     FIN (fddi_mac_mult_terminate ())
1288     FOUT;
1289     }

1290  Compcode
1291  fddi_mac_mult_init (pr_state_pptr)
```

```
1292      fddi_mac_mult_state        **pr_state_pptr;
1293      {
1294      static VosT_Cm_Obtype   obtype = OPC_NIL;

1295      FIN (fddi_mac_mult_init (pr_state_pptr))

1296      if (obtype == OPC_NIL)
1297         {
1298         if (Vos_Catmem_Register ("proc state vars (fddi_mac_mult)", sizeof
1299            (fddi_mac_mult_state), Vos_Nop, &obtype) == VOSC_FAILURE)
1300            FRET (OPC_COMPCODE_FAILURE)
1301         }

1302      if ((*pr_state_pptr = (fddi_mac_mult_state*) Vos_Catmem_Alloc
1303         (obtype, 1)) == OPC_NIL)
1304         FRET (OPC_COMPCODE_FAILURE)
1305      else
1306         {
1307         (*pr_state_pptr)->current_block = 20;
1308         FRET (OPC_COMPCODE_SUCCESS)
1309         }
1310      }

1311   /** The procedures defined in this section serve **/
1312   /** to simplify the code in the main body of the **/
1313   /** process model by providing primitives for timer **/
1314   /** manipulation.. **/

1315      fddi_timer_disable (timer_ptr)
1316         FddiT_Timer*      timer_ptr;
1317         {
1318   /* if the timer is already disabled, do nothing */
1319         if (timer_ptr->enabled)
1320            {
1321   /* disable the timer */
1322            timer_ptr->enabled = 0;

1323   /* reassign the accumulated time so far */
1324            timer_ptr->accum = op_sim_time () - timer_ptr->start_time;
1325            }
1326         }

1327      fddi_timer_enable (timer_ptr)
1328         FddiT_Timer*      timer_ptr;
1329         {
1330   /* if the timer is already enabled, simply return */
1331         if (!timer_ptr->enabled)
1332            {
1333   /* reenable the timer */
1334            timer_ptr->enabled = 1;
```

```
1335    /* set the start time to the current time */
1336    /* less the accumulated time so far */
1337            timer_ptr->start_time = op_sim_time () - timer_ptr->accum;
1338            }
1339        }

1340        fddi_timer_expired (timer_ptr)
1341            FddiT_Timer*        timer_ptr;
1342            {
1343            if (fddi_timer_remaining (timer_ptr) <= 0.0)
1344                return 1;
1345            else return 0;
1346            }


1347        double
1348        fddi_timer_remaining (timer_ptr)
1349            FddiT_Timer*        timer_ptr;
1350            {
1351    /* if the timer is enabled, update the accumulated time */
1352            if (timer_ptr->enabled)
1353                {
1354                timer_ptr->accum = op_sim_time () - timer_ptr->start_time;
1355                }

1356    /* return the timer remaining before expiration */
1357    /* a non-positive value indicates an expired timer */
1358            return (timer_ptr->target_accum - timer_ptr->accum);
1359            }

1360        double
1361        fddi_timer_value (timer_ptr)
1362            FddiT_Timer*        timer_ptr;
1363            {
1364    /* if the timer is enabled, update the accumulated time */
1365            if (timer_ptr->enabled)
1366                {
1367                timer_ptr->accum = op_sim_time () - timer_ptr->start_time;
1368                }
1369            return (timer_ptr->accum);
1370            }

1371        fddi_timer_set_value (timer_ptr, value)
1372            FddiT_Timer*        timer_ptr;
1373            double              value;
1374            {
1375            timer_ptr->accum = value;
1376            }

1377        fddi_timer_copy (from_timer_ptr, to_timer_ptr)
```

```
1378          FddiT_Timer*      from_timer_ptr;
1379          FddiT_Timer*      to_timer_ptr;
1380          {
1381          Vos_Copy_Memory (from_timer_ptr, to_timer_ptr, sizeof
1382             (FddiT_Timer));
1383          }

1384       fddi_timer_set (timer_ptr, duration)
1385          FddiT_Timer*      timer_ptr;
1386          {
1387    /* clear out accumulated time */
1388          timer_ptr->accum = 0.0;

1389    /* assign the timer duration */
1390          timer_ptr->target_accum = duration;

1391    /* assign the current time */
1392          timer_ptr->start_time = op_sim_time ();

1393    /* enable the timer */
1394          timer_ptr->enabled = 1;
1395          }

1396       FddiT_Timer*
1397       fddi_timer_create ()
1398          {
1399          FddiT_Timer*      timer_ptr;

1400    /* allocate memory for a timer structure */
1401          timer_ptr = (FddiT_Timer*) malloc (sizeof (FddiT_Timer));

1402    /* initialize the timer in the disabled mode */
1403          fddi_timer_init (timer_ptr);

1404    /* return the timer's address */
1405          return (timer_ptr);
1406          }

1407       fddi_timer_init (timer_ptr)
1408          FddiT_Timer*      timer_ptr;
1409          {
1410    /* the timer is initially disabled */
1411          timer_ptr->enabled = 0;

1412    /* the accumulated time is zero */
1413          timer_ptr->accum = 0.0;

1414    /* the target accumulated time is infinite */
1415          timer_ptr->target_accum = VOS_DOUBLE_INFINITY;
```

```
1416   /* the start time is now */
1417         timer_ptr->start_time = op_sim_time ();
1418         }

1419      fddi_station_register (address, objid)
1420         Objid          objid;
1421         int               address;
1422         {
1423   /* Fill an entry in the table which maps station */
1424   /* addresses to OPNET object ids */
1425         FIN (fddi_station_register (address, objid))

1426         Fddi_Address_Table [address] = objid;

1427   /* Keep track of total number of stations on the ring */
1428         Fddi_Num_Stations++;

1429         FOUT
1430         }

1431      fddi_tk_register ()
1432         {
1433   /* Register the station's intent to use the token.  */
1434   /* This should be done whenever an unregistered  */
1435   /* station obtains new data to transmit.  */
1436         FIN (fddi_tk_regsister ())

1437   /* increase the number of registered stations */
1438         Fddi_Num_Registered++;

1439   /* if the token is currently blocked, unblock it */
1440         if (Fddi_Tk_Blocked && Fddi_Tk_Accelerate)
1441            {
1442            fddi_tk_unblock ();
1443            }

1444         FOUT
1445         }

1446      fddi_tk_deregister ()
1447         {
1448   /* Cancel the station's intent to use the token.  */
1449   /* This should be done whenever a registered  */
1450   /* station exhausts its transmittable data.   */
1451         FIN (fddi_tk_deregsister ())

1452   /* decrease the number of registered stations */
1453         Fddi_Num_Registered--;

1454         FOUT
```

```
1455        }

1456     fddi_tk_indicate_no_data (token, address, delay)
1457        Packet*        token;
1458        int            address;
1459        double         delay;
1460        {
1461        FIN (fddi_tk_indicate_no_data (token, address, delay))

1462  /* The calling station is indicating that it has captured  */
1463  /* the token, but has no data to send.  If no other stations */
1464  /* have data to send either, the token may be blocked to gain  */
1465  /* simulation efficiency.                          */
1466        if (Fddi_Num_Registered == 0 && Fddi_Tk_Accelerate)
1467           {
1468           fddi_tk_block (token, address);
1469           }
1470        else{
1471  /* If the token cannot be blocked, send it into the ring. */
1472           op_pk_send_delayed (token, FDDI_PHY_STRM_OUT,
1473              delay + Fddi_Prop_Delay);
1474           }
1475        FOUT
1476        }

1477     fddi_tk_block (token, address)
1478        Packet*        token;
1479        int               address;
1480        {
1481        int            i;

1482        FIN (fddi_tk_block (token, address))

1483  /* Record the address of the blocking station and blocking time. */
1484        Fddi_Tk_Block_Base_Time = op_sim_time ();
1485        Fddi_Tk_Block_Base_Station = address;

1486        if (tk_trace_on == OPC_TRUE)
1487           {
1488           sprintf (str0, "Blocking Token: station (%d), time (%.9f)",
1489              Fddi_Tk_Block_Base_Station, Fddi_Tk_Block_Base_Time);
1490           op_prg_odb_print_major (str0, OPC_NIL);
1491           }

1492  /* Indicate that the token is blocked */
1493        Fddi_Tk_Blocked = 1;

1494  /* discard the token packet; another one will be */
1495  /* created when the token is unblocked. */
1496        op_pk_destroy (token);
```

```
1497   /* Cancel TRT timers at all MAC interfaces; otherwise these */
1498   /* timers may continue to expitr during the idle period,  */
1499   /* generating unnecessary events.                          */
1500         if (tk_trace_on == OPC_TRUE)
1501            {
1502            sprintf (str0, "Canceling timers for (%d) stations",
1503               Fddi_Num_Stations);
1504            op_prg_odb_print_major (str0, OPC_NIL);
1505            }

1506         for (i = 0; i < Fddi_Num_Stations; i++)
1507            {
1508   /* Retain the time at which the TRT would have expired; */
1509   /* this is used for calculations when the token is  */
1510   /* reinjected into the ring.                        */
1511            Fddi_Trt_Exp_Time [i] = op_ev_time (Fddi_Trt_Handle [i]);

1512   /* Cancel the TRT expiration event. */
1513            op_ev_cancel (Fddi_Trt_Handle [i]);
1514            }
1515         FOUT
1516         }

1517      fddi_tk_unblock ()
1518         {
1519         double         elapsed_time, first_tk_rx, last_tk_rx;
1520         double         tk_lap_time, next_time, current_time;
1521         double         dbl_num_hops, num_tk_rx, floor (), ceil ();
1522         int            i, num_hops, next_station;

1523         FIN (fddi_tk_unblock ())

1524   /* reset the blocking indicator */
1525         Fddi_Tk_Blocked = 0;

1526   /* Get the current time, used for many calculations below */
1527         current_time = op_sim_time ();

1528         if (tk_trace_on == OPC_TRUE)
1529            {
1530            sprintf (str0, "Unblocking token for ring (%d)", ring_id);
1531            op_prg_odb_print_major (str0, OPC_NIL);
1532            }

1533   /* For all stations on the ring, adjust TRT timer and Late_Ct flag. */
1534         for (i = 0; i < Fddi_Num_Stations; i++)
1535            {
1536            if (tk_trace_on == OPC_TRUE)
1537               {
1538               sprintf (str0, "adjusting state of station (%d)", i);
```

```
1539                op_prg_odb_print_minor ("", str0, OPC_NIL);
1540                }
1541    /* Calculate number of hops separating station i from block base
1542       station.  */
1543    /* In special case where i is the base station, the token must run a
1544       full */
1545    /* lap before returning.                                          */
1546             if (i != Fddi_Tk_Block_Base_Station)
1547                {
1548                num_hops = (i - Fddi_Tk_Block_Base_Station) %
1549                   Fddi_Num_Stations;
1550                if (num_hops < 0)
1551                   num_hops = Fddi_Num_Stations + num_hops;
1552                }
1553             else num_hops = Fddi_Num_Stations;

1554    /* Calculate first time at which token would have been received by
1555       station i.  */
1556    /* Note that initial release of token from base station takes a
1557       different */
1558    /* amount of time than repeating of token by other stations. Thus, the
1559       first */
1560    /* hop is assumed, and the base time is augmented by the time required
1561       to  */
1562    /* complete it.                                                   */
1563             first_tk_rx = Fddi_Tk_Block_Base_Time + FDDIC_TOKEN_TX_TIME +
1564                Fddi_Prop_Delay + (num_hops - 1) * Fddi_Tk_Hop_Delay;

1565             if (tk_trace_on == OPC_TRUE)
1566                {
1567                sprintf (str0, "station is (%d) hops from base", num_hops);
1568                sprintf (str1, "first receipt of token would be at (%.9f)",
1569                   first_tk_rx);
1570                op_prg_odb_print_minor (str0, str1, OPC_NIL);
1571                }

1572    /* Case 1: the token would not yet have been received by station i. */
1573             if (first_tk_rx > current_time)
1574                {
1575    /* Case 1a: the TRT at station i would not yet have expired. */
1576                if (Fddi_Trt_Exp_Time [i] > current_time)
1577                   {
1578    /* Late_Ct remains at its original value; only the TRT needs  */
1579    /* to be started again, with the same expiration time. */
1580                   TRT_SET (i, Fddi_Trt_Exp_Time [i])

1581                   if (tk_trace_on == OPC_TRUE)
1582                      {
1583                      sprintf (str0, "Restoring TRT to previous exp. time
1584                         (%.9f)", Fddi_Trt_Exp_Time [i]);
```

```
1585                    op_prg_odb_print_minor ("Token would not be received
1586                       and TRT not expired", str0, OPC_NIL);
1587                    }
1588                 }
1589    /* Case 1b: the TRT at station i would have expired. */
1590              else
1591                 {
1592    /* Late_Ct would have been set; also the timer would have been
1593       rescheduled  */
1594    /* for an entire TTRT at the time of expiration.              */
1595                 Fddi_Late_Ct [i] = 1;
1596                 TRT_SET (i, (Fddi_T_Opr + Fddi_Trt_Exp_Time [i]))

1597                 if (tk_trace_on == OPC_TRUE)
1598                    {
1599                    sprintf (str0, "Restoring TRT to proper exp. time
1600                       (%.9f)", Fddi_T_Opr + Fddi_Trt_Exp_Time [i]);
1601                    op_prg_odb_print_minor ("Token would not be received
1602                       and TRT would have expired", str0, OPC_NIL);
1603                    }
1604                 }
1605              }

1606    /* Case 2: the token would have been received (perhaps more than once).
1607       */
1608           else
1609              {
1610    /* Calculate the number of times the token would have been received */
1611    /* not including the first receipt.                  */
1612              tk_lap_time = Fddi_Tk_Hop_Delay * Fddi_Num_Stations;
1613              num_tk_rx = floor ((current_time - first_tk_rx) /
1614                 tk_lap_time);

1615    /* Calculate the latest time at which the token would have been
1616       received. */
1617              last_tk_rx = first_tk_rx + (num_tk_rx * tk_lap_time);

1618    /* Clear Late_Ct and schedule timer to expire at last receipt of token
1619       */
1620    /* plus one full TTRT. */
1621              Fddi_Late_Ct [i] = 0;
1622              TRT_SET (i, (last_tk_rx + Fddi_T_Opr))
1623              if (tk_trace_on == OPC_TRUE)
1624                 {
1625                 sprintf (str0, "token received (%g) times, last receipt
1626                    at (%.9f)", num_tk_rx + 1.0, last_tk_rx);
1627                 sprintf (str1, "Restoring TRT to proper exp. time
1628                    (%.9f)", last_tk_rx + Fddi_T_Opr);
1629                 op_prg_odb_print_minor ("Token would have been received;
1630                    Late_Ct is cleared", str1, str0, OPC_NIL);
```

```
1631                    }
1632                 }
1633              }

1634     /* compute the time since the token was blocked */
1635           elapsed_time = current_time - Fddi_Tk_Block_Base_Time;

1636     /* compute the number of hops completed on the ring. For the first hop
1637        */
1638     /* the token is transmitted directly, not repeated. For all remaining
1639        */
1640     /* hops, the delay is the station latency plus the propagation delay.
1641        */
1642     /* Thus, the first hop is assumed, and the remaining time for
1643        additional*/
1644     /* hops is computed beginning at the time where the token enters the  */
1645     /* base station's downstream neighbor */
1646           dbl_num_hops = 1.0 + (elapsed_time - FDDIC_TOKEN_TX_TIME -
1647              Fddi_Prop_Delay) / Fddi_Tk_Hop_Delay;

1648     /* If the token was unblocked in less time than it would have taken to
1649        */
1650     /* be fully transmitted by the base station, dbl_num_hops will be  */
1651     /* negative. However, 1 full hop would still be required before the  */
1652     /* token could be used, since the station had already committed to  */
1653     /* issuing the token. Thus, the actual of number of hops should never */
1654     /* be less than 1. If it is, round it to 1. */
1655           if (dbl_num_hops < 1.0)
1656              dbl_num_hops = 1.0;
1657           else
1658              {
1659     /* In all other cases, round the number of hops up to the nearest */
1660     /* integer value. If already an integer, then leave as is. */
1661              dbl_num_hops = ceil (dbl_num_hops);
1662              }

1663     /* Obtain an integer equivalent of dbl_num_hops. */
1664           num_hops = dbl_num_hops;

1665     /* Based on the number of hops and the base station, compute the */
1666     /* next station where the token will appear. */
1667           next_station = (num_hops + Fddi_Tk_Block_Base_Station) %
1668              Fddi_Num_Stations;

1669     /* Compute the time at which the token will appear there.  */
1670     /* Again, assume the first hop occurred, and measure time */
1671     /* from there forward.                        */
1672           next_time = Fddi_Tk_Block_Base_Time + (FDDIC_TOKEN_TX_TIME +
1673              Fddi_Prop_Delay) + (dbl_num_hops - 1.0) * Fddi_Tk_Hop_Delay;
```

```
1674          if (tk_trace_on == OPC_TRUE)
1675             {
1676             sprintf (str0, "Re-introducing token at station (%d), at time
1677                (%.9f)", next_station, next_time);
1678             op_prg_odb_print_minor (str0, OPC_NIL);
1679             }

1680    /* reinject the token at that station */
1681          fddi_tk_inject (next_station, next_time);

1682          FOUT
1683          }

1684       fddi_tk_inject (address, arv_time)
1685          int             address;
1686          double          arv_time;
1687          {
1688    /* Re-insert the token into the ring after an idle period. */
1689          FIN (fddi_tk_inject (address, arv_time))

1690    /* The token is recreated and reinserted onto the ring */
1691    /* at the specified station which is not necesssarily the */
1692    /* station now requesting the token. */
1693    /* The station which will r  nsert the token is */
1694    /* asked to do so by means  f a remote interrupt. */
1695          op_intrpt_schedule_remote (arv_time, FDDIC_TK_INJECT,
1696                Fddi_Address_Table [address]);

1697          FOUT
1698          }

1699       fddi_load_frame_attrs (dest_addr_ptr, svc_class_ptr, pri_level_ptr)
1700          int             *dest_addr_ptr, *svc_class_ptr, *pri_level_ptr;
1701          {
1702       int                NUM_PRIOS, i;    /* 26JAN94 */
1703          Packet          *pkptr;

1704          FIN (fddi_load_frame_attrs (dest_addr_ptr, svc_class_ptr,
1705             pri_level_ptr))

1706    /* remove next packet in queue */
1707    /* 27DEC94: loop structure superimposed to handle a bank of subqueues.
1708       */
1709    /* Extract the packet with the highest priority, that is, the packet
1710       */
1711    /* at the head of the highest-numbered subqueue containing packets.
1712       */
1713    /* Note that the C language vector numbering convention numbers the
1714       */
1715    /* subqueues from 0 to 7, while FDDI convention is to number the */
```

```
1716  /* corresponding asynchronous priorities from 1 to 8. This is */
1717  /* reconciled in the statistical outputs available in the Analysis */
1718  /* Editor, where labels are assigned accordingly.  Also note that */
1719  /* synchronous traffic is assigned priority 8 as an artifice to allow
1720     */
1721  /* routing through a separate subqueue, by which statistics may be */
1722  /* gathered for traffic by class and by priority. -Nix */
1723     NUM_PRIOS = 9;
1724     for (i = NUM_PRIOS - 1; i > -1; i--)
1725        {
1726           if (op_subq_stat (i, OPC_QSTAT_PKSIZE) > 0.0)
1727           {
1728           pkptr = op_subq_pk_remove (i, OPC_QPOS_HEAD);
1729           break;
1730           }
1731        }
1732  /* extract the fields of interest */
1733        op_pk_nfd_get (pkptr, "dest_addr", dest_addr_ptr);
1734        op_pk_nfd_get (pkptr, "svc_class", svc_class_ptr);

1735  /* only read priority level if frame is asynchronous */
1736        if (*svc_class_ptr == FDDI_SVC_ASYNC)
1737           op_pk_nfd_get (pkptr, "pri", pri_level_ptr);

1738  /* replace the packet on the proper subqueue */
1739        op_subq_pk_insert (i, pkptr, OPC_QPOS_HEAD);
1740        FOUT
1741        }
```

APPENDIX F. SOURCE "C" CODE:

"fddi_gen_mult.pr.c"

The line numbering in this appendix is used for reference within this thesis only, and does not correspond with that seen in OPNET®'s text editors.

```
1   /* Process model C form file: fddi_gen_mult.pr.c */
2   /* Portions of this file Copyright (C) MIL 3, Inc. 1992 */

3   /* OPNET system definitions */
4   #include <opnet.h>
5   #include "fddi_gen_mult.pr.h"
6   FSM_EXT_DECS

7   /* Header block */
8   #define MAC_LAYER_OUT_STREAM        0

9   /* define possible service classes for frames */
10  #define FDDI_SVC_ASYNC          0
11  #define FDDI_SVC_SYNC           1

12  /* define token classes */
13  #define FDDI_TK_NONRESTRICTED   0
14  #define FDDI_TK_RESTRICTED      1

15  /* 07FEB94: define the number of stations    */
16  #define NUM_STATIONS       50

17  /* a global counting variable */
18  /* nt genARRIVAL = 0; */

19  /* State variable definitions */
20  typedef struct
21     {
22     FSM_SYS_STATE
23     Distribution*        sv_inter_dist_ptr;
24     Distribution*        sv_len_dist_ptr;
25     Distribution*        sv_dest_dist_ptr;
26     Distribution*        sv_pkt_priority_ptr;
27     Objid                sv_mac_objid;
```

```
28      Objid                   sv_my_id;
29      int                        sv_low_dest_addr;
30      int                        sv_high_dest_addr;
31      int                        sv_station_addr;
32      int                        sv_low_pkt_priority;
33      int                        sv_high_pkt_priority;
34      double                     sv_arrival_rate;
35      double                     sv_mean_pk_len;
36      double                     sv_async_mix;
37      Ici*                    sv_mac_iciptr;
38      Distribution*           sv_num_addees_dist_ptr;
39      int                        sv_num_addees;
40      int                        sv_min_num_addees;
41      int                        sv_max_num_addees;
42      int                        sv_dest_addr[NUM_STATIONS+1];
43      } fddi_gen_mult_state;

44  #define pr_state_ptr           ((fddi_gen_mult_state*)
45  SimI_Mod_State_Ptr)
46  #define inter_dist_ptr         pr_state_ptr->sv_inter_dist_ptr
47  #define len_dist_ptr           pr_state_ptr->sv_len_dist_ptr
48  #define dest_dist_ptr          pr_state_ptr->sv_dest_dist_ptr
49  #define pkt_priority_ptr       pr_state_ptr->sv_pkt_priority_ptr
50  #define mac_objid              pr_state_ptr->sv_mac_objid
51  #define my_id                     pr_state_ptr->sv_my_id
52  #define low_dest_addr          pr_state_ptr->sv_low_dest_addr
53  #define high_dest_addr         pr_state_ptr->sv_high_dest_addr
54  #define station_addr           pr_state_ptr->sv_station_addr
55  #define low_pkt_priority       pr_state_ptr->sv_low_pkt_priority
56  #define high_pkt_priority      pr_state_ptr->sv_high_pkt_priority
57  #define arrival_rate           pr_state_ptr->sv_arrival_rate
58  #define mean_pk_len            pr_state_ptr->sv_mean_pk_len
59  #define async_mix              pr_state_ptr->sv_async_mix
60  #define mac_iciptr                pr_state_ptr->sv_mac_iciptr
61  #define num_addees_dist_ptr    pr_state_ptr->sv_num_addees_dist_ptr
62  #define num_addees                pr_state_ptr->sv_num_addees
63  #define min_num_addees         pr_state_ptr->sv_min_num_addees
64  #define max_num_addees         pr_state_ptr->sv_max_num_addees
65  #define dest_addr              pr_state_ptr->sv_dest_addr

66  /* Process model interrupt handling procedure */

67  void
68  fddi_gen_mult ()
69     {
70     Packet      *pkptr;
71     int         pklen;
72     int         *da_ptr;

73     int         i, restricted;
```

```
74      int         pkt_prio;
75      int         nix;

76      FSM_ENTER (fddi_gen_mult)

77      FSM_BLOCK_SWITCH
78         {
79   /*---------------------------------------------------------*/
80   /** state (INIT) enter executives **/
81         FSM_STATE_ENTER_UNFORCED (0, state0_enter_exec, "INIT")
82            {
83   /* determine id of own processor to use in finding attrs */
84            my_id = op_id_self ();

85   /* 07FEB94: determine the upper and lower limits for multiple  */
86   /* addressing from this station. -Nix      */
87            op_ima_obj_attr_get (my_id, "min num addees", &min_num_addees);
88            op_ima_obj_attr_get (my_id, "max num addees", &max_num_addees);

89   /* 07FEB94: set up a distribution for the number of stations   */
90   /* receive this packet. -Nix            */
91            num_addees_dist_ptr = op_dist_load ("uniform_int",
92               min_num_addees, max_num_addees);

93   /* determine address range for uniform desination assignment */
94            op_ima_obj_attr_get (my_id, "low dest address",
95   &low_dest_addr);
96            op_ima_obj_attr_get (my_id, "high dest address",
97               &high_dest_addr);

98   /* determine object id of connected 'mac' layer process */
99            mac_objid = op_topo_assoc (my_id, OPC_TOPO_ASSOC_OUT,
100                OPC_OBJMTYPE_MODULE, MAC_LAYER_OUT_STREAM);

101   /* determine the address assigned to it */
102   /* which is also the address of this station */
103           op_ima_obj_attr_get (mac_objid, "station_address",
104              &station_addr);

105   /* set up a distribution for generation of addresses */
106           dest_dist_ptr = op_dist_load ("uniform_int", low_dest_addr,
107              high_dest_addr);

108   /* added 26DEC93 */
109   /* determine priority range for uniform traffic generation */
110           op_ima_obj_attr_get (my_id, "high pkt priority",
111              &high_pkt_priority);
112           op_ima_obj_attr_get (my_id, "low pkt priority",
113              &low_pkt_priority);
```

```
114   /* set up a distribution for generation of priorities */
115            pkt_priority_ptr = op_dist_load ("uniform_int",
116               low_pkt_priority, high_pkt_priority);

117   /* above added 26DEC93 */

118   /* also determine the arrival rate for packet generation */
119            op_ima_obj_attr_get (my_id, "arrival rate", &arrival_rate);

120   /* determine the mix of asynchronous and synchronous */
121   /* traffic. This is expressed as the proportion of */
122   /* asynchronous traffic. i.e a value of 1.0 indicates */
123   /* that all the produced traffic shall be asynchronous. */
124            op_ima_obj_attr_get (my_id, "async_mix", &async_mix);

125   /* set up a distribution for arrival generations */
126            if (arrival_rate != 0.0)
127               {
128   /* arrivals are exponentially distributed, with given mean */
129               inter_dist_ptr = op_dist_load ("constant", 1.0 /
130                  arrival_rate, 0.0);

131   /* determine the distribution for packet size */
132               op_ima_obj_attr_get (my_id, "mean pk length", &mean_pk_len);

133   /* set up corresponding distribution */
134               len_dist_ptr = op_dist_load ("constant", mean_pk_len, 0.0);
135
136   /* designate the time of first arrival */
137               fddi_gen_schedule ();

138   /* set up an interface control information (ICI) structure */
139   /* to communicate parameters to the mac layer process */
140   /* (it is more efficient to set one up now and keep it */
141   /* as a state variable than to allocate one on each packet xfer) */
142               mac_iciptr = op_ici_create ("fddi_mac_req");
143               }
144            }

145   /** blocking after enter executives of unforced state. **/
146         FSM_EXIT (1,fddi_gen_mult)

147   /** state (INIT) exit executives **/
148         FSM_STATE_EXIT_UNFORCED (0, state0_exit_exec, "INIT")
149            {
150            }

151   /** state (INIT) transition processing **/
152         FSM_TRANSIT_FORCE (1, state1_enter_exec, ;)
153   /*---------------------------------------------------------*/
```

```
154   /** state (ARRIVAL) enter executives **/
155         FSM_STATE_ENTER_UNFORCED (1, state1_enter_exec, "ARRIVAL")
156            {
157   /* determine the length of the packet to be generated */
158            pklen = op_dist_outcome (len_dist_ptr);

159   /* 07FEB94: re-initialize the destination address array  */
160   /* to zeros. -Nix           */
161            for (i = 0; i < NUM_STATIONS+1; i++)
162            {
163            dest_addr[i] = 0;
164            }

165   /* determine the destination */
166   /* don't allow this station's address as a possible outcome */
167   /* gen_packet:                  */
168   /* dest_addr = op_dist_outcome (dest_dist_ptr); */
169   /* if (dest_addr != -1 && dest_addr == station_addr) */
170   /* goto gen_packet;               */

171   /* 07FEB94: determine the destinations. -Nix */

172   /* Determine the number of stations to receive this packet  */
173            num_addees = op_dist_outcome (num_addees_dist_ptr);

174   /* Find these stations, using num_addees as a counter. -Nix */
175            for (i = num_addees; i > 0; i--)
176            {
177            gen_packet:
178            nix = op_dist_outcome (dest_dist_ptr);
179            if (dest_addr[nix] == 1 || nix == station_addr)
180               {
181               goto gen_packet;
182               }
183            dest_addr[nix] = 1;
184            }

185   /* 05MAR94: because the op_pk_nfd_get() command in FR_REPEAT  */
186   /* overwrites the first field with the array address, an  */*/
187   /* offset needs to be applied so that the dest_array[0] */
188   /* contents aren't lost; that is, one field more than the    */*/
189   /* number of stations is included to allow a one-step shift   */
190   /* that will preserve the address array.  In fddi_mac, all    */
191   /* references to dest_addr must allow for this shift. -Nix    */
192            for (i=NUM_STATIONS; i>0; i--)
193            dest_addr[i] = dest_addr[i-1];
194

195   /* 26DEC94 & 29JAN94: determine its priority */
196            pkt_prio =  op_dist_outcome (pkt_priority_ptr);
```

```
197   /* create a packet to send to mac */
198            pkptr = op_pk_create_fmt ("fddi_llc_fr");

199   /* assign its overall size. */
200            op_pk_total_size_set (pkptr, pklen);

201   /* assign the time of creation */
202            op_pk_nfd_set (pkptr, "cr_time", op_sim_time ());

203   /* place the destination address into the ICI */
204   /* (the protocol_type field will default) */

205   /* 15MAR94: note that dest_addr now serves as a */
206   /* pointer to an array in memory, as it is the */
207   /* name of an array of what will be 0s and 1s. -Nix */
208            op_ici_attr_set (mac_iciptr, "dest_addr", dest_addr);

209   /* assign the priority, and requested token class */
210   /* also assign the service class */

211   /* 29JAN94: the fddi_llc_fr format is modified */
212   /* to include a "pri" field. -Nix */

213            if (op_dist_uniform (1.0) <= async_mix)
214               {
215               op_pk_nfd_set (pkptr, "pri", pkt_prio);/* 29JAN94 */
216   op_ici_attr_set (mac_iciptr, "svc_class", FDDI_SVC_ASYNC);
217               op_ici_attr_set (mac_iciptr, "pri", pkt_prio);  /* 29JAN94
218   */
219               }
220            else{
221               op_pk_nfd_set (pkptr, "pri", 8);/* 29JAN94 */
222               op_ici_attr_set (mac_iciptr, "svc_class", FDDI_SVC_SYNC);
223               op_ici_attr_set (mac_iciptr, "pri", 8);    /* 29JAN94 */
224               }

225   /* Request only nonrestricted tokens after transmission */
226            op_ici_attr_set (mac_iciptr, "tk_class",
227   FDDI_TK_NONRESTRICTED);

228   /* send the packet coupled with the ICI */
229            op_ici_install (mac_iciptr);
230            op_pk_send (pkptr, MAC_LAYER_OUT_STREAM);

231   /* ******************************************************* */
232   /*  17MAR94: count and report the running total number    */
233   /*  of packets generated. -Nix       */
234   /* ******************************************************* */
235   /* genARRIVAL ++;  */
236   /* printf("Packets generated: %d\n", genARRIVAL); */
```

```
237   /* schedule the next arrival */
238            fddi_gen_schedule ();

239   /* ******************************************************* */
240   /* 18FEB94: print out the address, and the contents. */
241   /* for (i=0; i<NUM_STATIONS+1; i++)  */
242   /*  printf("ARRIVAL: %d. address: %X; contents: %d\n",    */
243   /*  i, &(dest_addr[i]), dest_addr[i]);  */
244   /* ******************************************************* */
245            }

246   /** blocking after enter executives of unforced state. **/
247         FSM_EXIT (3,fddi_gen_mult)

248   /** state (ARRIVAL) exit executives **/
249         FSM_STATE_EXIT_UNFORCED (1, statel_exit_exec, "ARRIVAL")
250            {
251            }

252   /** state (ARRIVAL) transition processing **/
253         FSM_TRANSIT_FORCE (1, statel_enter_exec, ;)
254   /*---------------------------------------------------------*/
255         }

256      FSM_EXIT (0,fddi_gen_mult)
257      }

258   void
259   fddi_gen_mult_svar (prs_ptr,var_name,var_p_ptr)
260      fddi_gen_mult_state     *prs_ptr;
261      char        *var_name, **var_p_ptr;
262      {

263      FIN (fddi_gen_mult_svar (prs_ptr))

264      *var_p_ptr = VOS_NIL;
265      if (Vos_String_Equal ("inter_dist_ptr" , var_name)) *var_p_ptr =
266         (char*) (&prs_ptr->sv_inter_dist_ptr);
267      if (Vos_String_Equal ("len_dist_ptr" , var_name)) *var_p_ptr =
268         (char*) (&prs_ptr->sv_len_dist_ptr);
269      if (Vos_String_Equal ("dest_dist_ptr" , var_name)) *var_p_ptr =
270         (char*) (&prs_ptr->sv_dest_dist_ptr);
271      if (Vos_String_Equal ("pkt_priority_ptr" , var_name)) *var_p_ptr =
272         (char*) (&prs_ptr->sv_pkt_priority_ptr);
273      if (Vos_String_Equal ("mac_objid" , var_name)) *var_p_ptr = (char*)
274         (&prs_ptr->sv_mac_objid);
275      if (Vos_String_Equal ("my_id" , var_name)) *var_p_ptr = (char*)
276         (&prs_ptr->sv_my_id);
277      if (Vos_String_Equal ("low_dest_addr" , var_name)) *var_p_ptr =
278         (char*) (&prs_ptr->sv_low_dest_addr);
```

```
279      if (Vos_String_Equal ("high_dest_addr" , var_name)) *var_p_ptr =
280         (char*) (&prs_ptr->sv_high_dest_addr);
281      if (Vos_String_Equal ("station_addr" , var_name)) *var_p_ptr =
282         (char*) (&prs_ptr->sv_station_addr);
283      if (Vos_String_Equal ("low_pkt_priority" , var_name)) *var_p_ptr =
284         (char*) (&prs_ptr->sv_low_pkt_priority);
285      if (Vos_String_Equal ("high_pkt_priority" , var_name)) *var_p_ptr
286         = (char*) (&prs_ptr->sv_high_pkt_priority);
287      if (Vos_String_Equal ("arrival_rate" , var_name)) *var_p_ptr =
288         (char*) (&prs_ptr->sv_arrival_rate);
289      if (Vos_String_Equal ("mean_pk_len" , var_name)) *var_p_ptr = (char*)
290         (&prs_ptr->sv_mean_pk_len);
291      if (Vos_String_Equal ("async_mix" , var_name)) *var_p_ptr = (char*)
292         (&prs_ptr->sv_async_mix);
293      if (Vos_String_Equal ("mac_iciptr" , var_name)) *var_p_ptr = (char*)
294         (&prs_ptr->sv_mac_iciptr);
295      if (Vos_String_Equal ("num_addees_dist_ptr" , var_name)) *var_p_ptr =
296         (char*) (&prs_ptr->sv_num_addees_dist_ptr);
297      if (Vos_String_Equal ("num_addees" , var_name)) *var_p_ptr = (char*)
298         (&prs_ptr->sv_num_addees);
299      if (Vos_String_Equal ("min_num_addees" , var_name)) *var_p_ptr =
300         (char*) (&prs_ptr->sv_min_num_addees);
301      if (Vos_String_Equal ("max_num_addees" , var_name)) *var_p_ptr =
302         (char*) (&prs_ptr->sv_max_num_addees);
303      if (Vos_String_Equal ("dest_addr" , var_name)) *var_p_ptr = (char*)
304         (prs_ptr->sv_dest_addr);
305      FOUT;
306      }

307   void
308   fddi_gen_mult_diag ()
309      {
310      Packet      *pkptr;
311      int         pklen;
312      int         *da_ptr;

313      int         i, restricted;
314      int   pkt_prio;
315      int   nix;

316      FIN (fddi_gen_mult_diag ())

317      FOUT;
318      }

319   void
320   fddi_gen_mult_terminate ()
321      {
322      Packet      *pkptr;
323      int         pklen;
```

```
324     int         *da_ptr;

325     int         i, restricted;
326     int         pkt_prio;
327     int         nix;

328     FIN (fddi_gen_mult_terminate ())

329     FOUT;
330     }

331  Compcode
332  fddi_gen_mult_init (pr_state_pptr)
333     fddi_gen_mult_state        **pr_state_pptr;
334     {
335     static VosT_Cm_Obtype   obtype = OPC_NIL;

336     FIN (fddi_gen_mult_init (pr_state_pptr))

337     if (obtype == OPC_NIL)
338        {
339        if (Vos_Catmem_Register ("proc state vars (fddi_gen_mult)",
340           sizeof (fddi_gen_mult_state), Vos_Nop, &obtype) ==
341  VOSC_FAILURE)
342           FRET (OPC_COMPCODE_FAILURE)
343        }

344     if ((*pr_state_pptr = (fddi_gen_mult_state*) Vos_Catmem_Alloc
345        (obtype, 1)) == OPC_NIL)
346        FRET (OPC_COMPCODE_FAILURE)
347     else
348        {
349        (*pr_state_pptr)->current_block = 0;
350        FRET (OPC_COMPCODE_SUCCESS)
351        }
352     }

353  /* static added 2DEC93, on advice from MIL 3, Inc. */
354     static
355     fddi_gen_schedule ()
356        {
357        double         inter_time;

358  /* obtain an interarrival period according to the */
359  /* prescribed distribution */
360        inter_time = op_dist_outcome (inter_dist_ptr);

361  /* schedule the arrival of next generated packet */
362        op_intrpt_schedule_self (op_sim_time () + inter_time, 0);
363        }
```

APPENDIX G. SINK "C" CODE:

"fddi_sink_mult.pr.c"

The line numbering in this appendix is used for reference within this thesis only, and does not correspond with that seen in OPNET®'s text editors.

```
1   /* Process model C form file: fddi_sink_mult.pr.c */
2   /* Portions of this file Copyright (C) MIL 3, Inc. 1992 */

3   /* OPNET system definitions */
4   #include <opnet.h>
5   #include "fddi_sink_mult.pr.h"
6   FSM_EXT_DECS

7   /* Header block */
8   /* Globals */
9   /* array format installed 20JAN94; positions 0-7 represent the asynch
10     priority levels, PRIORITIES + 1  */
11  /* represents synch traffic, and grand totals are as given in the
12     original. */
13

14  #define  PRIORITIES  8  /* 20JAN94 */
15  double      fddi_sink_accum_delay = 0.0;
16  double      fddi_sink_accum_delay_a[PRIORITIES + 1] = {0.0, 0.0, 0.0,
17     0.0, 0.0, 0.0, 0.0, 0.0, 0.0};
18  int      fddi_sink_total_pkts = 0;
19  int      fddi_sink_total_pkts_a[PRIORITIES + 1]  = {0, 0, 0, 0, 0, 0, 0,
20     0, 0};
21  double      fddi_sink_total_bits = 0.0;
22  double      fddi_sink_total_bits_a[PRIORITIES + 1]  = {0.0, 0.0, 0.0,
23     0.0, 0.0, 0.0, 0.0, 0.0, 0.0};
24  double      fddi_sink_peak_delay = 0.0;
25  double      fddi_sink_peak_delay_a[PRIORITIES + 2]  = {0.0, 0.0, 0.0,
26     0.0, 0.0, 0.0, 0.0, 0.0, 0.0};
27  int      fddi_sink_scalar_write = 0;
28  int      pri_set = 20;/* 20JAN94 */

29  /* Externally defined globals. */
30  extern double  fddi_t_opr [];
```

```
/* Attributes fromEnvironment file.  */
double   Offered_Load;/* 12JAN94 */
double   Asynch_Offered_Load; /* 12JAN94 */

/* transition expressions */
#define END_OF_SIM   op_intrpt_type() == OPC_INTRPT_ENDSIM

/* State variable definitions */
typedef struct
   {
   FSM_SYS_STATE
   Gshandle             sv_thru_gshandle;
   Gshandle             sv_m_delay_gshandle;
   Gshandle             sv_ete_delay_gshandle;
   Gshandle             sv_thru_gshandle_a[10];
   Gshandle             sv_m_delay_gshandle_a[10];
   Gshandle             sv_ete_delay_gshandle_a[9];
   } fddi_sink_mult_state;

#define pr_state_ptr             ((fddi_sink_mult_state*)
   SimI_Mod_State_Ptr)
#define thru_gshandle            pr_state_ptr->sv_thru_gshandle
#define m_delay_gshandle         pr_state_ptr->sv_m_delay_gshandle
#define ete_delay_gshandle    pr_state_ptr->sv_ete_delay_gshandle
#define thru_gshandle_a          pr_state_ptr->sv_thru_gshandle_a
#define m_delay_gshandle_a    pr_state_ptr->sv_m_delay_gshandle_a
#define ete_delay_gshandle_a  pr_state_ptr->sv_ete_delay_gshandle_a

/* Process model interrupt handling procedure */

void
fddi_sink_mult ()
   {
   double         delay, creat_time;
   Packet*     pkptr;
   int            src_addr, my_addr;
   Ici*        from_mac_ici_ptr;
   double         fddi_sink_ttrt;

   FSM_ENTER (fddi_sink_mult)

   FSM_BLOCK_SWITCH
      {
/*--------------------------------------------------------*/
/** state (DISCARD) enter executives **/
      FSM_STATE_ENTER_UNFORCED (0, state0_enter_exec, "DISCARD")
         {
/* get the packet and the interface control info */
         pkptr = op_pk_get (op_intrpt_strm ());
         from_mac_ici_ptr = op_intrpt_ici ();
```

```
74  /* 20JAN94: get the packet's priority level, which   */
75  /* will be used to index arrays of thruput and delay */
76  /* computations.              */
77  /* pri_set = op_pk_priority_get (pkptr); doesn't work here */
78            op_pk_nfd_get (pkptr, "pri", &pri_set);/* 29JAN94 */

79  /* add in its size */
80            fddi_sink_total_bits += op_pk_total_size_get (pkptr);
81            fddi_sink_total_bits_a[pri_set] += op_pk_total_size_get
82  (pkptr); /* 20JAN94 */

83  /* determine the time of creation of the packet */
84            op_pk_nfd_get (pkptr, "cr_time", &creat_time);

85  /* accumulate delays */
86            delay = op_sim_time () - creat_time;
87            fddi_sink_accum_delay += delay;
88            fddi_sink_accum_delay_a[pri_set] += delay;  /* 20JAN94 */

89  /* keep track of peak delay value */
90            if (delay > fddi_sink_peak_delay)
91               fddi_sink_peak_delay = delay;

92  /* 20JAN94: keep track by priority levels as well 23JAN94 */
93            if (delay > fddi_sink_peak_delay_a[pri_set])
94            fddi_sink_peak_delay_a[pri_set] = delay;

95  /* ******************************************** */
96  /* printf("DISCARD: pri_set is %d\n", pri_set);  */
97  /* ******************************************** */

98  /* destroy the packet */
99            op_pk_destroy (pkptr);

100 /* 03FEB94: To convert this to the "fddi_sink_link"   */
101 /* model, deactivate the 'destroy' code, and activate */
102 /* the following 'enqueue' code.  This is a first */
103 /* step toward developing a LAN bridging structure.   */
104 /* -Nix                    */
105 /* op_subq_pk_insert (pri_set, pkptr, OPC_QPOS_TAIL); */

106 /* increment packet counter; 20JAN94 */
107           fddi_sink_total_pkts++;
108           fddi_sink_total_pkts_a[pri_set]++;

109 /* if a multiple of 25 packets is reached, update stats */
110 /* 03FEB94: [0]->[7] represent asynch priorities 1->8,  */
111 /* respectively; [8] represents synchronous traffic,    */
112 /* and [9] represents overall asynchronous traffic.-Nix */
113           if (fddi_sink_total_pkts % 1 == 0)
```

```
114             {
115             op_stat_global_write (thru_gshandle, fddi_sink_total_bits /
116                op_sim_time ());
117             op_stat_global_write (thru_gshandle_a[0],
118                fddi_sink_total_bits_a[0] / op_sim_time());
119             op_stat_global_write (thru_gshandle_a[1],
120                fddi_sink_total_bits_a[1] / op_sim_time());
121             op_stat_global_write (thru_gshandle_a[2],
122                fddi_sink_total_bits_a[2] / op_sim_time());
123             op_stat_global_write
124                (thru_gshandle_a[3],fddi_sink_total_bits_a[3] /
125                op_sim_time());
126             op_stat_global_write (thru_gshandle_a[4],
127                fddi_sink_total_bits_a[4] / op_sim_time());
128             op_stat_global_write (thru_gshandle_a[5],
129                fddi_sink_total_bits_a[5] / op_sim_time());
130             op_stat_global_write (thru_gshandle_a[6],
131                fddi_sink_total_bits_a[6] / op_sim_time());
132             op_stat_global_write (thru_gshandle_a[7],
133                fddi_sink_total_bits_a[7] / op_sim_time());
134             op_stat_global_write (thru_gshandle_a[8],
135                fddi_sink_total_bits_a[8] / op_sim_time());

136     /* 30JAN94: gather all asynch stats into one overall figure */
137             op_stat_global_write (thru_gshandle_a[9],
138                (fddi_sink_total_bits_a[0] + fddi_sink_total_bits_a[1] +
139                fddi_sink_total_bits_a[2] + fddi_sink_total_bits_a[3] +
140                fddi_sink_total_bits_a[4] + fddi_sink_total_bits_a[5] +
141                fddi_sink_total_bits_a[6] + fddi_sink_total_bits_a[7]) /
142                op_sim_time());
143
144     /* (fddi_sink_total_bits - fddi_sink_total_bits_a[8]) /  */
145     /* op_sim_time()); */

146             op_stat_global_write (m_delay_gshandle,   fddi_sink_accum_delay
147                / fddi_sink_total_pkts);
148             op_stat_global_write (m_delay_gshandle_a[0],
149                fddi_sink_accum_delay_a[0] / fddi_sink_total_pkts_a[0]);
150             op_stat_global_write (m_delay_gshandle_a[1],
151                fddi_sink_accum_delay_a[1] / fddi_sink_total_pkts_a[1]);
152             op_stat_global_write (m_delay_gshandle_a[2],
153                fddi_sink_accum_delay_a[2] / fddi_sink_total_pkts_a[2]);
154             op_stat_global_write (m_delay_gshandle_a[3],
155                fddi_sink_accum_delay_a[3] / fddi_sink_total_pkts_a[3]);
156             op_stat_global_write (m_delay_gshandle_a[4],
157                fddi_sink_accum_delay_a[4] / fddi_sink_total_pkts_a[4]);
158             op_stat_global_write (m_delay_gshandle_a[5],
159                fddi_sink_accum_delay_a[5] / fddi_sink_total_pkts_a[5]);
160             op_stat_global_write (m_delay_gshandle_a[6],
161                fddi_sink_accum_delay_a[6] / fddi_sink_total_pkts_a[6]);
```

```
162             op_stat_global_write (m_delay_gshandle_a[7],
163                fddi_sink_accum_delay_a[7] / fddi_sink_total_pkts_a[7]);
164             op_stat_global_write (m_delay_gshandle_a[8],
165                fddi_sink_accum_delay_a[8] / fddi_sink_total_pkts_a[8]);

166     /* 30JAN94: gather all asynch stats into one figure */
167             op_stat_global_write (m_delay_gshandle_a[9],
168                (fddi_sink_accum_delay_a[0] + fddi_sink_accum_delay_a[1] +
169                fddi_sink_accum_delay_a[2] + fddi_sink_accum_delay_a[3] +
170                fddi_sink_accum_delay_a[4] + fddi_sink_accum_delay_a[5] +
171                fddi_sink_accum_delay_a[6] + fddi_sink_accum_delay_a[7]) /
172                (fddi_sink_total_pkts_a[0] + fddi_sink_total_pkts_a[1] +
173                fddi_sink_total_pkts_a[2] + fddi_sink_total_pkts_a[3] +
174                fddi_sink_total_pkts_a[4] + fddi_sink_total_pkts_a[5] +
175                fddi_sink_total_pkts_a[6] + fddi_sink_total_pkts_a[7]));
176
177     /* (fddi_sink_accum_delay - fddi_sink_accum_delay_a[8]) /  */
178     /* (fddi_sink_total_pkts - fddi_sink_total_pkts_a[8]));    */
179                }

180     /* also record actual delay values */
181             op_stat_global_write (ete_delay_gshandle, delay);
182             op_stat_global_write (ete_delay_gshandle_a[pri_set], delay);


183             }

184     /** blocking after enter executives of unforced state. **/
185           FSM_EXIT (1,fddi_sink_mult)

186     /** state (DISCARD) exit executives **/
187           FSM_STATE_EXIT_UNFORCED (0, state0_exit_exec, "DISCARD")
188              {
189              }

190     /** state (DISCARD) transition processing **/
191           FSM_INIT_COND (END_OF_SIM)
192           FSM_DFLT_COND
193           FSM_TEST_LOGIC ("DISCARD")

194           FSM_TRANSIT_SWITCH
195              {
196              FSM_CASE_TRANSIT (0, 1, state1_enter_exec, ;)
197              FSM_CASE_TRANSIT (1, 0, state0_enter_exec, ;)
198              }
199     /*---------------------------------------------------------*/

200     /** state (STATS) enter executives **/
201           FSM_STATE_ENTER_UNFORCED (1, state1_enter_exec, "STATS")
202              {
```

```
/* At end of simulation, scalar performance statistics */
/* and input parameters are written out. */

/* Only one station needs to do this */
        if (!fddi_sink_scalar_write)
           {
/* set the scalar write flag */
           fddi_sink_scalar_write = 1;
           op_stat_scalar_write ("Mean End-to-End Delay (sec.),
              Priority 1", fddi_sink_accum_delay_a[0] /
              fddi_sink_total_pkts_a[0]);
           op_stat_scalar_write ("Mean End-to-End Delay (sec.),
              Priority 2", fddi_sink_accum_delay_a[1] /
              fddi_sink_total_pkts_a[1]);
           op_stat_scalar_write ("Mean End-to-End Delay (sec.),
              Priority 3", fddi_sink_accum_delay_a[2] /
              fddi_sink_total_pkts_a[2]);
           op_stat_scalar_write ("Mean End-to-End Delay (sec.),
              Priority 4", fddi_sink_accum_delay_a[3] /
              fddi_sink_total_pkts_a[3]);
           op_stat_scalar_write ("Mean End-to-End Delay (sec.),
              Priority 5", fddi_sink_accum_delay_a[4] /
              fddi_sink_total_pkts_a[4]);
           op_stat_scalar_write ("Mean End-to-End Delay (sec.),
              Priority 6", fddi_sink_accum_delay_a[5] /
              fddi_sink_total_pkts_a[5]);
           op_stat_scalar_write ("Mean End-to-End Delay (sec.),
              Priority 7", fddi_sink_accum_delay_a[6] /
              fddi_sink_total_pkts_a[6]);
           op_stat_scalar_write ("Mean End-to-End Delay (sec.),
              Priority 8", fddi_sink_accum_delay_a[7] /
              fddi_sink_total_pkts_a[7]);
           op_stat_scalar_write ("Mean End-to-End Delay (sec.),
              Asynchronous", (fddi_sink_accum_delay -
              fddi_sink_accum_delay_a[8]) / (fddi_sink_total_pkts -
              fddi_sink_total_pkts_a[8]));

/* (fddi_sink_accum_delay_a[0] + fddi_sink_accum_delay_a[1] +  */
/* fddi_sink_accum_delay_a[2] + fddi_sink_accum_delay_a[3] +  */
/* fddi_sink_accum_delay_a[4] + fddi_sink_accum_delay_a[5] +  */
/* fddi_sink_accum_delay_a[6] + fddi_sink_accum_delay_a[7]) / */
/* (fddi_sink_total_pkts_a[0] + fddi_sink_total_pkts_a[1] +    */
/* fddi_sink_total_pkts_a[2] + fddi_sink_total_pkts_a[3] +    */
/* fddi_sink_total_pkts_a[4] + fddi_sink_total_pkts_a[5] +    */
/* fddi_sink_total_pkts_a[6] + fddi_sink_total_pkts_a[7]));   */

           op_stat_scalar_write ("Mean End-to-End Delay (sec.),
              Synchronous", fddi_sink_accum_delay_a[8] /
              fddi_sink_total_pkts_a[8]);
```

```
249             op_stat_scalar_write ("Mean End-to-End Delay (sec.), Total",
250                fddi_sink_accum_delay / fddi_sink_total_pkts);
251             op_stat_scalar_write ("Throughput (bps), Priority 1",
252                fddi_sink_total_bits_a[0] / op_sim_time ());
253             op_stat_scalar_write ("Throughput (bps), Priority 2",
254                fddi_sink_total_bits_a[1] / op_sim_time ());
255             op_stat_scalar_write ("Throughput (bps), Priority 3",
256                fddi_sink_total_bits_a[2] / op_sim_time ());
257             op_stat_scalar_write ("Throughput (bps), Priority 4",
258                fddi_sink_total_bits_a[3] / op_sim_time ());
259             op_stat_scalar_write ("Throughput (bps), Priority 5",
260                fddi_sink_total_bits_a[4] / op_sim_time ());

261             op_stat_scalar_write ("Throughput (bps), Priority 6",
262                fddi_sink_total_bits_a[5] / op_sim_time ());
263             op_stat_scalar_write ("Throughput (bps), Priority 7",
264                fddi_sink_total_bits_a[6] / op_sim_time ());
265             op_stat_scalar_write ("Throughput (bps), Priority 8",
266                fddi_sink_total_bits_a[7] / op_sim_time ());
267             op_stat_scalar_write ("Throughput (bps), Asynchronous",
268                (fddi_sink_total_bits - fddi_sink_total_bits_a[8]) /
269                op_sim_time ());

270 /* (fddi_sink_total_bits_a[0] + fddi_sink_total_bits_a[1] +   */
271 /* fddi_sink_total_bits_a[2] + fddi_sink_total_bits_a[3] +   */
272 /* fddi_sink_total_bits_a[4] + fddi_sink_total_bits_a[5] +   */
273 /* fddi_sink_total_bits_a[6] + fddi_sink_total_bits_a[7]) /  */
274 /* op_sim_time ()); */

275          op_stat_scalar_write ("Throughput (bps), Synchronous",
276             fddi_sink_total_bits_a[8] / op_sim_time ());
277          op_stat_scalar_write ("Throughput (bps), Total",
278             fddi_sink_total_bits / op_sim_time ());
279          op_stat_scalar_write ("Peak End-to-End Delay (sec.), Priority
280             1",    fddi_sink_peak_delay_a[0]);
281          op_stat_scalar_write ("Peak End-to-End Delay (sec.), Priority
282             2",    fddi_sink_peak_delay_a[1]);
283          op_stat_scalar_write ("Peak End-to-End Delay (sec.), Priority
284             3",    fddi_sink_peak_delay_a[2]);
285          op_stat_scalar_write ("Peak End-to-End Delay (sec.), Priority
286             4",    fddi_sink_peak_delay_a[3]);
287          op_stat_scalar_write ("Peak End-to-End Delay (sec.), Priority
288             5",    fddi_sink_peak_delay_a[4]);
289          op_stat_scalar_write ("Peak End-to-End Delay (sec.), Priority
290             6",    fddi_sink_peak_delay_a[5]);
291          op_stat_scalar_write ("Peak End-to-End Delay (sec.), Priority
292             7",    fddi_sink_peak_delay_a[6]);
293          op_stat_scalar_write ("Peak End-to-End Delay (sec.), Priority
294             8",    fddi_sink_peak_delay_a[7]);
```

```
295             op_stat_scalar_write ("Peak End-to-End Delay (sec.),
296                Synchronous",   fddi_sink_peak_delay_a[8]);
297             op_stat_scalar_write ("Peak End-to-End Delay (sec.), Overall",
298                fddi_sink_peak_delay);

299     /* Write the TTRT value for ring 0.  This preserves */
300     /* the old behavior for single-ring simulations. */
301                op_stat_scalar_write ("TTRT (sec.) - Ring 0",fddi_t_opr
302     [0]);

303     /* 12JAN94: obtain offered load information from the Environment */
304     /* file; this will be used to provide abscissa information that  */
305     /* can be plotted in the Analysis Editor (see "fddi_sink" STATS  */
306     /* state. To the user: it's your job to keep these current in    */
307     /* the Environment File. -Nix */
308             op_ima_sim_attr_get (OPC_IMA_DOUBLE, "total_offered_load",
309                &Offered_Load);
310             op_ima_sim_attr_get (OPC_IMA_DOUBLE, "asynch_offered_load",
311                &Asynch_Offered_Load);
312

313             /* 12JAN94: write the total offered load for this run */
314             op_stat_scalar_write ("Total Offered Load
315     (Mbps)",Offered_Load);
316             op_stat_scalar_write ("Asynchronous Offered Load (Mbps)",
317                Asynch_Offered_Load);
318                }
319             }

320     /** blocking after enter executives of unforced state. **/
321           FSM_EXIT (3,fddi_sink_mult)

322     /** state (STATS) exit executives **/
323           FSM_STATE_EXIT_UNFORCED (1, state1_exit_exec, "STATS")
324              {
325              }

326     /** state (STATS) transition processing **/
327           FSM_TRANSIT_MISSING ("STATS")
328     /*---------------------------------------------------------*/

329     /** state (INIT) enter executives **/
330           FSM_STATE_ENTER_FORCED (2, state2_enter_exec, "INIT")
331              {
332     /* get the gshandles of the global statistic to be obtained */
333     /* 20JAN94: set array format */
334
335             thru_gshandle_a[0] = op_stat_global_reg ("pri 1 throughput
336                (bps)");
```

```
337             thru_gshandle_a[1] = op_stat_global_reg ("pri 2 throughput
338                (bps)");
339             thru_gshandle_a[2] = op_stat_global_reg ("pri 3 throughput
340                (bps)");
341             thru_gshandle_a[3] = op_stat_global_reg ("pri 4 throughput
342                (bps)");
343             thru_gshandle_a[4] = op_stat_global_reg ("pri 5 throughput
344                (bps)");
345             thru_gshandle_a[5] = op_stat_global_reg ("pri 6 throughput
346                (bps)");
347             thru_gshandle_a[6] = op_stat_global_reg ("pri 7 throughput
348                (bps)");
349             thru_gshandle_a[7] = op_stat_global_reg ("pri 8 throughput
350                (bps)");
351             thru_gshandle_a[8] = op_stat_global_reg ("synch throughput
352                (bps)");
353             thru_gshandle_a[9] = op_stat_global_reg ("async throughput
354                (bps)");
355             thru_gshandle  = op_stat_global_reg ("total throughput (bps)");
356             m_delay_gshandle_a[0] = op_stat_global_reg ("pri 1 mean delay
357                (sec.)");
358             m_delay_gshandle_a[1] = op_stat_global_reg ("pri 2 mean delay
359                (sec.)");
360             m_delay_gshandle_a[2] = op_stat_global_reg ("pri 3 mean delay
361                (sec.)");
362             m_delay_gshandle_a[3] = op_stat_global_reg ("pri 4 mean delay
363                (sec.)");
364             m_delay_gshandle_a[4] = op_stat_global_reg ("pri 5 mean delay
365                (sec.)");
366             m_delay_gshandle_a[5] = op_stat_global_reg ("pri 6 mean delay
367                (sec.)");
368             m_delay_gshandle_a[6] = op_stat_global_reg ("pri 7 mean delay
369                (sec.)");
370             m_delay_gshandle_a[7] = op_stat_global_reg ("pri 8 mean delay
371                (sec.)");
372             m_delay_gshandle_a[8] = op_stat_global_reg ("synch mean delay
373                (sec.)");
374             m_delay_gshandle_a[9] = op_stat_global_reg ("async mean delay
375                (sec.)");
376             m_delay_gshandle  = op_stat_global_reg ("total mean delay
377                (sec.)");
378             ete_delay_gshandle_a[0] = op_stat_global_reg ("pri 1 end-to-end
379                delay (sec.)");
380             ete_delay_gshandle_a[1] = op_stat_global_reg ("pri 2 end-to-end
381                delay (sec.)");
382             ete_delay_gshandle_a[2] = op_stat_global_reg ("pri 3 end-to-end
383                delay (sec.)");
384             ete_delay_gshandle_a[3] = op_stat_global_reg ("pri 4 end-to-end
385                delay (sec.)");
```

```
386            ete_delay_gshandle_a[4] = op_stat_global_reg ("pri 5 end-to-end
387               delay (sec.)");
388            ete_delay_gshandle_a[5] = op_stat_global_reg ("pri 6 end-to-end
389               delay (sec.)");
390            ete_delay_gshandle_a[6] = op_stat_global_reg ("pri 7 end-to-end
391               delay (sec.)");
392            ete_delay_gshandle_a[7] = op_stat_global_reg ("pri 8 end-to-end
393               delay (sec.)");
394            ete_delay_gshandle_a[8] = op_stat_global_reg ("synch end-to-end
395               delay (sec.)");
396            ete_delay_gshandle   = op_stat_global_reg ("total end-to-end
397               delay (sec.)");
398            }

399   /** state (INIT) exit executives **/
400         FSM_STATE_EXIT_FORCED (2, state2_exit_exec, "INIT")
401            {
402            }

403   /** state (INIT) transition processing **/
404         FSM_INIT_COND (END_OF_SIM)
405         FSM_DFLT_COND
406         FSM_TEST_LOGIC ("INIT")

407         FSM_TRANSIT_SWITCH
408            {
409            FSM_CASE_TRANSIT (0, 1, state1_enter_exec, ;)
410            FSM_CASE_TRANSIT (1, 0, state0_enter_exec, ;)
411            }
412   /*---------------------------------------------------------*/
413         }

414      FSM_EXIT (2,fddi_sink_mult)
415      }

416   void
417   fddi_sink_mult_svar (prs_ptr,var_name,var_p_ptr)
418      fddi_sink_mult_state    *prs_ptr;
419      char        *var_name, **var_p_ptr;
420      {

421      FIN (fddi_sink_mult_svar (prs_ptr))

422      *var_p_ptr = VOS_NIL;
423      if (Vos_String_Equal ("thru_gshandle" , var_name))
424         *var_p_ptr = (char *) (&prs_ptr->sv_thru_gshandle);
425      if (Vos_String_Equal ("m_delay_gshandle" , var_name))
426         *var_p_ptr = (char *) (&prs_ptr->sv_m_delay_gshandle);
427      if (Vos_String_Equal ("ete_delay_gshandle" , var_name))
428         *var_p_ptr = (char *) (&prs_ptr->sv_ete_delay_gshandle);
```

```
429      if (Vos_String_Equal ("thru_gshandle_a" , var_name))
430         *var_p_ptr = (char *) (prs_ptr->sv_thru_gshandle_a);
431      if (Vos_String_Equal ("m_delay_gshandle_a" , var_name))
432         *var_p_ptr = (char *) (prs_ptr->sv_m_delay_gshandle_a);
433      if (Vos_String_Equal ("ete_delay_gshandle_a" , var_name))
434         *var_p_ptr = (char *) (prs_ptr->sv_ete_delay_gshandle_a);
435      FOUT;
436      }

437   void
438   fddi_sink_mult_diag ()
439      {
440      double          delay, creat_time;
441      Packet*      pkptr;
442      int             src_addr, my_addr;
443      Ici*         from_mac_ici_ptr;
444      double          fddi_sink_ttrt;

445      FIN (fddi_sink_mult_diag ())
446      FOUT;
447      }

448   void
449   fddi_sink_mult_terminate ()
450      {
451      double          delay, creat_time;
452      Packet*      pkptr;
453      int             src_addr, my_addr;
454      Ici*         from_mac_ici_ptr;
455      double          fddi_sink_ttrt;

456      FIN (fddi_sink_mult_terminate ())
457      FOUT;
458      }

459   Compcode
460   fddi_sink_mult_init (pr_state_pptr)
461      fddi_sink_mult_state       **pr_state_pptr;
462      {
463      static VosT_Cm_Obtype   obtype = OPC_NIL;

464      FIN (fddi_sink_mult_init (pr_state_pptr))

465      if (obtype == OPC_NIL)
466         {
467         if (Vos_Catmem_Register ("proc state vars (fddi_sink_mult)",
468            sizeof (fddi_sink_mult_state), Vos_Nop, &obtype) ==
469            VOSC_FAILURE)
470            FRET (OPC_COMPCODE_FAILURE)
471         }
```

```
472       if ((*pr_state_pptr = (fddi_sink_mult_state*) Vos_Catmem_Alloc
473          (obtype, 1)) == OPC_NIL)
474          FRET (OPC_COMPCODE_FAILURE)
475       else
476          {
477          (*pr_state_pptr)->current_block = 4;
478          FRET (OPC_COMPCODE_SUCCESS)
479          }
480       }
```

APPENDIX H. ENVIRONMENT FILE FOR

50-STATION MULTICAST CAPABLE FDDI LAN

```
# fddi50mult.ef
# sample simulation configuration file for fddi example model
# 50 station network with multiple addressing capability

#*** Attributes related to loading used by "fddi_gen" ***

# station addresses
*.*.f0.mac.station_address: 0
*.*.f1.mac.station_address: 1
*.*.f2.mac.station_address: 2
*.*.f3.mac.station_address: 3
*.*.f4.mac.station_address: 4
*.*.f5.mac.station_address: 5
*.*.f6.mac.station_address: 6
*.*.f7.mac.station_address: 7
*.*.f8.mac.station_address: 8
*.*.f9.mac.station_address: 9
*.*.f10.mac.station_address: 10
*.*.f11.mac.station_address: 11
*.*.f12.mac.station_address: 12
*.*.f13.mac.station_address: 13
*.*.f14.mac.station_address: 14
*.*.f15.mac.station_address: 15
*.*.f16.mac.station_address: 16
*.*.f17.mac.station_address: 17
*.*.f18.mac.station_address: 18
*.*.f19.mac.station_address: 19
*.*.f20.mac.station_address: 20
*.*.f21.mac.station_address: 21
*.*.f22.mac.station_address: 22
*.*.f23.mac.station_address: 23
*.*.f24.mac.station_address: 24
*.*.f25.mac.station_address: 25
*.*.f26.mac.station_address: 26
*.*.f27.mac.station_address: 27
*.*.f28.mac.station_address: 28
*.*.f29.mac.station_address: 29
*.*.f30.mac.station_address: 30
*.*.f31.mac.station_address: 31
*.*.f32.mac.station_address: 32
*.*.f33.mac.station_address: 33
```

```
*.*.f34.mac.station_address: 34
*.*.f35.mac.station_address: 35
*.*.f36.mac.station_address: 36
*.*.f37.mac.station_address: 37
*.*.f38.mac.station_address: 38
*.*.f39.mac.station_address: 39
*.*.f40.mac.station_address: 40
*.*.f41.mac.station_address: 41
*.*.f42.mac.station_address: 42
*.*.f43.mac.station_address: 43
*.*.f44.mac.station_address: 44
*.*.f45.mac.station_address: 45
*.*.f46.mac.station_address: 46
*.*.f47.mac.station_address: 47
*.*.f48.mac.station_address: 48
*.*.f49.mac.station_address: 49

*.*.*.mac.ring_id :0

# Range number of stations that may recieve this packet if more
# than one is designated (model defaults are both 1)
# Note that the code does not allow the originating station to
# address a packet to itself, so max_num_addees is less than
# the number of stations.
"*.*.llc_src.min num addees" :      1
"*.*.llc_src.max num addees" :      1

# destination addresses for random message generation
"*.*.llc_src.low dest address" :    0
"*.*.llc_src.high dest address" :   49

# "*.f0.llc_src.low dest address"  : 0
# "*.f0.llc_src.high dest address" : 49

# range of priority values that can be assigned to packets; FDDI
# standards allow for 8 priorities of asynchronous traffic.  MIL3's
# original model is modified to allow each station to generate
# multiple priorities, within a specified range.  (Note that while
# research literature refers to asynchronous priorities ranging
# from 1 to 8, the corresponding numbering here is 0 to 7, in
# keeping with the C language array element numbering convention.)
"*.*.llc_src.high pkt priority" :  7
"*.*.llc_src.low pkt priority" :   0

# arrival rate(frames/sec), and message size (bits) for random
# message generation at each station on the ring.
"*.*.*.arrival rate" :          750
"*.*.*.mean pk length" :        1000

# These are the synchronus transmitters
```

```
"*.f0.*.arrival rate"    :  6000
"*.f0.*.mean pk length"  :  512
"*.f5.*.arrival rate"    :  6000
"*.f5.*.mean pk length"  :  512
"*.f10.*.arrival rate"   :  6000
"*.f10.*.mean pk length" :  512
"*.f15.*.arrival rate"   :  6000
"*.f15.*.mean pk length" :  512
"*.f20.*.arrival rate"   :  60000
"*.f20.*.mean pk length" :  512
"*.f25.*.arrival rate"   :  6000
"*.f25.*.mean pk length" :  512
"*.f30.*.arrival rate"   :  6000
"*.f30.*.mean pk length" :  512
"*.f35.*.arrival rate"   :  6000
"*.f35.*.mean pk length" :  512
"*.f40.*.arrival rate"   :  6000
"*.f40.*.mean pk length" :  512
"*.f45.*.arrival rate"   :  60000
"*.f45.*.mean pk length" :  512

# 12DEC93: total offered load is the sum of all stations'
# loads (Mbps). Compute this by hand; this value is used in
# the sink process model for generating scalar plots where
# offered load is the abscissa.
total_offered_load :                    60.72
asynch_offered_load :                   30.00

# set the proportion of asynchronous traffic
# a value of 1.0 indicates all asynchronous traffic
"*.*.*.async_mix" :        1.0

"*.f0.*.async_mix" :    0.0
"*.f5.*.async_mix" :    0.0
"*.f10.*.async_mix" :   0.0
"*.f15.*.async_mix" :   0.0
"*.f20.*.async_mix" :   0.0
"*.f25.*.async_mix" :   0.0
"*.f30.*.async_mix" :   0.0
"*.f35.*.async_mix" :   0.0
"*.f40.*.async_mix" :   0.0
"*.f45.*.async_mix" :   0.0

#*** Ring configuration attributes used by "fddi_mac" ***

# allocate percentage of synchronous bandwidth to each station
# this value should not exceed 1 for all stations combined; OPNET
# does not
# enforce this; 01FEB94: this must be less than 1; see equation below
"*.*.mac.sync bandwidth" :    0.0
```

```
"*.f0.mac.sync bandwidth"  :  .09358
"*.f5.mac.sync bandwidth"  :  .09358
"*.f10.mac.sync bandwidth" :  .09358
"*.f15.mac.sync bandwidth" :  .09358
"*.f20.mac.sync bandwidth" :  .09358
"*.f25.mac.sync bandwidth" :  .09358
"*.f30.mac.sync bandwidth" :  .09358
"*.f35.mac.sync bandwidth" :  .09358
"*.f40.mac.sync bandwidth" :  .09358
"*.f45.mac.sync bandwidth" :  .09358

# Target Token Rotation Time (one half of maximum
# synchronous response time)
"*.*.mac.T_Req" :             .0107

# Index of the station which initially launches the token
"spawn station":              0

# Delay incurred by packets as they traverse a station's
# ring interface (see Powers, p. 351 for a discussion
# of this (Powers gives 1usec, but 60.0e-08 agrees with
# Dykeman & Bux)
station_latency:             60.0e-08

# Propagation Delay separating stations on the ring.
prop_delay:                5.085e-06

# Simulation related attributes

# Token Acceleration Mechanism enabling flag.
# It is reccomended that this mechanism be enabled for most
# situations
accelerate_token:                     1

seed:                                10

# Run control attributes

duration:                             .5
verbose_sim:                       TRUE
upd_int:                            .1
os_file:                     fddi50mult
ov_file:                     fddi50mult

# Opnet Debugger (odb) enabling attribute
# debug:                     TRUE
```

APPENDIX I. CONVENTIONS

One of the purposes of this report is that it will be used both as a teaching tool and a springboard for future researchers and assessors of fiber optic network simulations implementing OPNET®. Throughout the writing of this report, the author has kept these goals in sight and the resulting narrative contains technical stylistic conventions in keeping with the projected use of this material in a teaching, reference, and research environment. These conventions, implemented in the narrative portion of this report only, are briefly described here.

All excerpted programming code fragments are isolated on their own lines within the text and highlighted by a `standard san-serif font.` Variable names, function names, and names of programming objects referred to within the text of the report are also highlighted in this manner, with a `standard san-serif font.` Messages from the computer or responses to be made to computer queries are set off in double quotes and a `"bold standard san-serif font."` Single keystrokes are highlighted in capitalized italics, (e.g. <*CTRL+S*>), while parameters are also set off in the same manner (e.g. <*number of nodes*>).

APPENDIX J. GLOSSARY

BONeS®	Block Oriented Network Simulator
CDL	Common Data Link
DSPO	Defense Support Project Office
Environment file	A command file containing descriptors and values utilized by a system to define the operating parameters. Sometimes this file is referred to as the "configuration file."
FDDI LAN	Fiber Distributed Data Interface Local Area Network
MAC	Medium Access Control
OPNET®	Optimized Network Engineering Tool
THT	Token Holding Timer
TRT	Token Rotation Timer
TTRT	Target Token Rotation Time

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